

Chapel Background and Motivation





COMPUTE | STORE | ANALYZE

Safe Harbor Statement



This presentation may contain forward-looking statements that are based on our current expectations. Forward looking statements may include statements about our financial guidance and expected operating results, our opportunities and future potential, our product development and new product introduction plans, our ability to expand and penetrate our addressable markets and other statements that are not historical facts. These statements are only predictions and actual results may materially vary from those projected. Please refer to Cray's documents filed with the SEC from time to time concerning factors that could affect the Company and these forward-looking statements.



Chapel's Origins: HPCS



DARPA HPCS: High Productivity Computing Systems

- Goal: improve productivity by a factor of 10x
- Timeframe: summer 2002 fall 2012
- Cray developed a new system architecture, network, software, ...
 - this became the very successful Cray XC30™ Supercomputer Series



...and a new programming language: Chapel





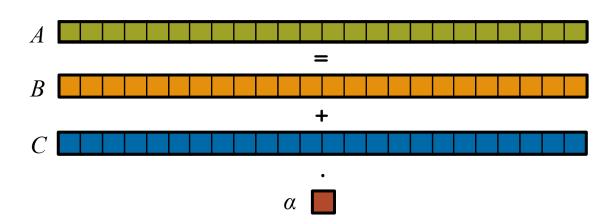
COMPUTE

CRAY

Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m$, $A_i = B_i + \alpha \cdot C_i$

In pictures:



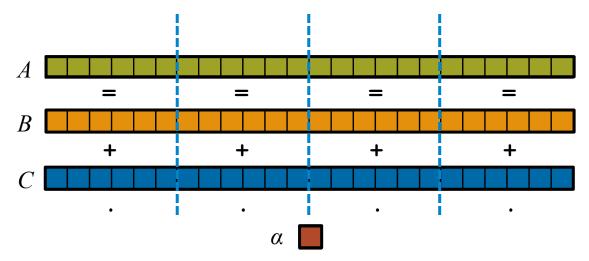


1

Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m$, $A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel:



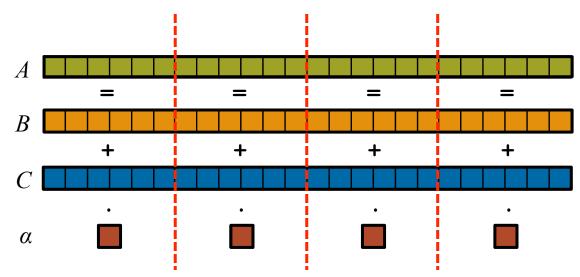


on

Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m$, $A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel (distributed memory):

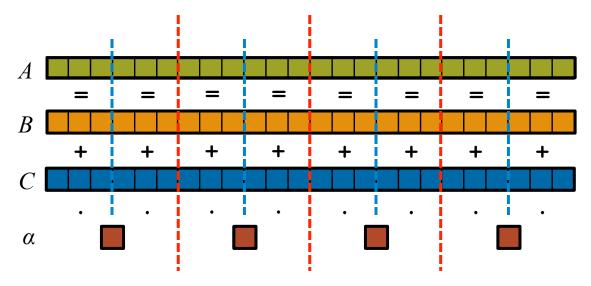




Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel (distributed memory multicore):





COMPUTE

STREAM Triad: MPI

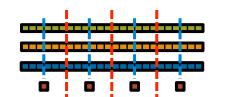


```
#include <hpcc.h>
static int VectorSize;
static double *a, *b, *c;
int HPCC StarStream(HPCC Params *params) {
 int myRank, commSize;
  int rv, errCount;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size( comm, &commSize );
 MPI Comm rank( comm, &myRank );
 rv = HPCC Stream( params, 0 == myRank);
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM,
   0, comm);
 return errCount;
int HPCC Stream(HPCC Params *params, int doIO) {
 register int j;
 double scalar;
 VectorSize = HPCC LocalVectorSize( params, 3,
   sizeof(double), 0);
 a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC( double, VectorSize );
 c = HPCC XMALLOC( double, VectorSize );
```

```
if (!a || !b || !c) {
  if (c) HPCC free(c);
  if (b) HPCC free(b);
  if (a) HPCC free(a);
  if (doIO) {
    fprintf( outFile, "Failed to allocate memory (%d).
 \n", VectorSize );
    fclose( outFile );
  return 1;
for (j=0; j<VectorSize; j++) {</pre>
 b[j] = 2.0;
  c[j] = 0.0;
scalar = 3.0;
for (j=0; j<VectorSize; j++)</pre>
  a[j] = b[j] + scalar*c[j];
HPCC free(c);
HPCC free (b);
HPCC free(a);
```



STREAM Triad: MPI+OpenMP





MPI + OpenMP

```
#include <hpcc.h>
                                                       if (!a || !b || !c) {
#ifdef OPENMP
                                                         if (c) HPCC free(c);
#include <omp.h>
                                                         if (b) HPCC free(b);
#endif
                                                         if (a) HPCC free(a);
static int VectorSize;
                                                         if (doIO) {
static double *a, *b, *c;
                                                           fprintf( outFile, "Failed to allocate memory (%d).
                                                         \n", VectorSize );
int HPCC StarStream(HPCC Params *params) {
                                                           fclose( outFile );
  int myRank, commSize;
  int rv, errCount;
                                                         return 1;
  MPI Comm comm = MPI COMM WORLD;
 MPI Comm size( comm, &commSize );
                                                     #ifdef OPENMP
  MPI Comm rank( comm, &myRank );
                                                     #pragma omp parallel for
                                                     #endif
  rv = HPCC Stream( params, 0 == myRank);
                                                       for (j=0; j<VectorSize; j++) {</pre>
  MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM,
                                                         b[j] = 2.0;
   0, comm);
                                                         c[j] = 0.0;
  return errCount;
                                                       scalar = 3.0;
int HPCC Stream(HPCC Params *params, int doIO) {
                                                     #ifdef OPENMP
  register int j;
                                                     #pragma omp parallel for
  double scalar;
                                                     #endif
                                                       for (j=0; j<VectorSize; j++)</pre>
  VectorSize = HPCC LocalVectorSize( params, 3,
                                                         a[j] = b[j] + scalar*c[j];
   sizeof(double), 0);
                                                       HPCC free(c);
  a = HPCC XMALLOC( double, VectorSize );
                                                       HPCC free(b);
 b = HPCC XMALLOC( double, VectorSize );
                                                       HPCC free(a);
  c = HPCC XMALLOC( double, VectorSize );
```



STREAM Triad: MPI+OpenMP vs. CUDA

MPI + OpenMP

```
#ifdef OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;
int HPCC StarStream(HPCC Params *params) {
   int myRank, commSize;
   int rv, errCount;
   MPI_Comm comm = MPI_COMM_WORLD;

MPI_Comm_size( comm, &commSize );
   MPI_Comm_rank( comm, &myRank );
   rv = HPCC Stream( params, 0 == myRank);
   MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM, 0, comm );
   return errCount;
```

CUDA

```
#define N 2000000
int main() {
  float *d_a, *d_b, *d_c;
  float scalar;

cudaMalloc((void**)&d_a, sizeof(float)*N);
  cudaMalloc((void**)&d_b, sizeof(float)*N);
  cudaMalloc((void**)&d_c, sizeof(float)*N);
```

HPC suffers from too many distinct notations for expressing parallelism and locality

dim3 dimBlock(128)

```
VectorSize = HPCC LocalVectorSize( params, 3, sizeof(double), 0 );
 a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC( double, VectorSize );
 c = HPCC XMALLOC( double, VectorSize );
  if (!a || !b || !c) {
    if (c) HPCC free(c);
    if (b) HPCC free(b);
    if (a) HPCC free(a);
    if (doIO) {
      fprintf( outFile, "Failed to allocate memory (%d).\n", VectorSize );
      fclose( outFile );
    return 1:
#ifdef OPENMP
#pragma omp parallel for
#endif
  for (j=0; j<VectorSize; j++) {
   b[j] = 2.0;
   c[j] = 0.0;
 scalar = 3.0;
#ifdef OPENMP
#pragma omp parallel for
 for (j=0; j<VectorSize; j++)
   a[i] = b[i] + scalar*c[i];
 HPCC free(c);
 HPCC free (b) :
 HPCC free(a);
  return 0;
```



COMPUTE

Why so many programming models?



HPC has traditionally given users...

- ...low-level, control-centric programming models
- ...ones that are closely tied to the underlying hardware
- ...ones that support only a single type of parallelism

Type of HW Parallelism	Programming Model	Unit of Parallelism
Inter-node	MPI	executable
Intra-node/multicore	OpenMP/pthreads	iteration/task
Instruction-level vectors/threads	pragmas	iteration
GPU/accelerator	CUDA/OpenCL/OpenACC	SIMD function/task

benefits: lots of control; decent generality; easy to implement downsides: lots of user-managed detail; brittle to changes



STORE | ANALYZE

Rewinding a few slides...

MPI + OpenMP

```
CUDA
```

```
#define N 2000000
int main() {
  float *d_a, *d_b, *d_c;
  float scalar;

cudaMalloc((void**)&d_a, sizeof(float)*N);
  cudaMalloc((void**)&d_b, sizeof(float)*N);
  cudaMalloc((void**)&d_c, sizeof(float)*N);
```

HPC suffers from too many distinct notations for expressing parallelism and locality

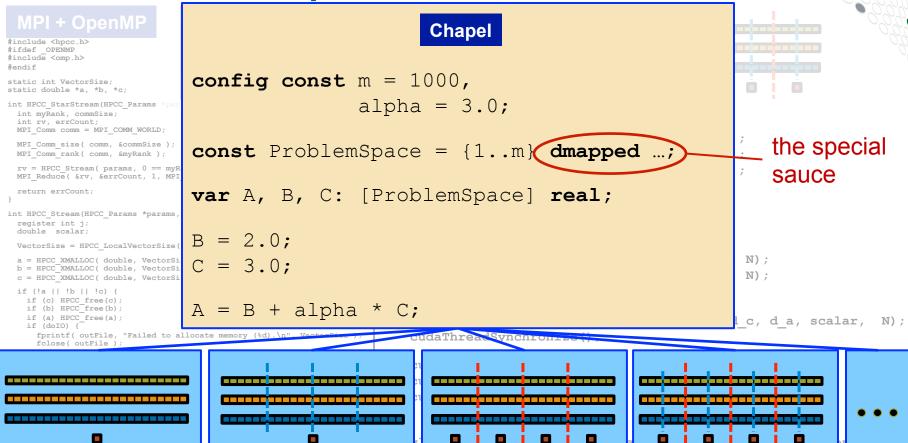
dim3 dimBlock(128)

```
VectorSize = HPCC LocalVectorSize( params, 3, sizeof(double), 0 );
 a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC( double, VectorSize );
 c = HPCC XMALLOC( double, VectorSize );
  if (!a || !b || !c) {
    if (c) HPCC free(c);
    if (b) HPCC free(b);
    if (a) HPCC free(a);
    if (doIO) {
      fprintf( outFile, "Failed to allocate memory (%d).\n", VectorSize );
      fclose( outFile );
    return 1:
#ifdef OPENMP
#pragma omp parallel for
#endif
  for (j=0; j<VectorSize; j++) {
   b[j] = 2.0;
   c[i] = 0.0;
 scalar = 3.0;
#ifdef OPENMP
#pragma omp parallel for
 for (j=0; j<VectorSize; j++)
   a[i] = b[i] + scalar*c[i];
 HPCC free(c);
 HPCC free (b) :
 HPCC free(a);
  return 0;
```



COMPUTE

STREAM Triad: Chapel



<u>Philosophy:</u> Good language design can tease details of locality and parallelism away from an algorithm, permitting the compiler, runtime, applied scientist, and HPC expert to each focus on their strengths.



#endif

HPCC

Motivating Chapel Themes

- 1) General Parallel Programming
- 2) Global-View Abstractions
- 3) Multiresolution Design
- 4) Control over Locality/Affinity
- 5) Reduce HPC ← Mainstream Language Gap



1) General Parallel Programming



With a unified set of concepts...

...express any parallelism desired in a user's program

- Styles: data-parallel, task-parallel, concurrency, nested, ...
- Levels: model, function, loop, statement, expression

...target any parallelism available in the hardware

• Types: machines, nodes, cores, instructions

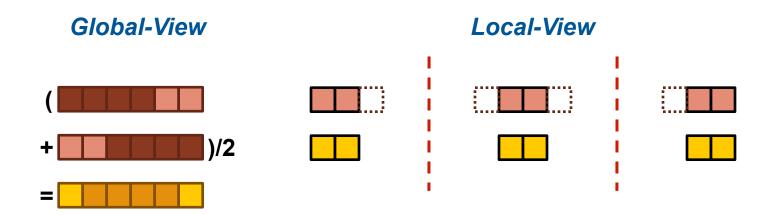
Type of HW Parallelism	Programming Model	Unit of Parallelism
Inter-node	Chapel	executable/task
Intra-node/multicore	Chapel	iteration/task
Instruction-level vectors/threads	Chapel	iteration
GPU/accelerator	Chapel	SIMD function/task



| ANALYZE



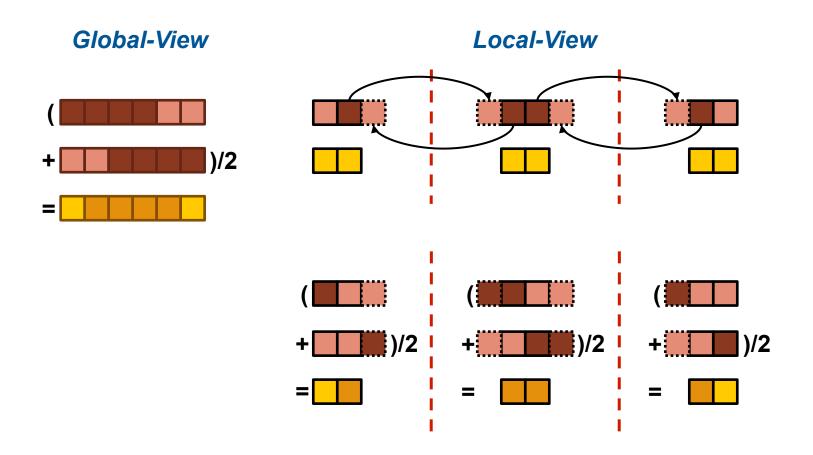
In pictures: "Apply a 3-Point Stencil to a vector"







In pictures: "Apply a 3-Point Stencil to a vector"





In code: "Apply a 3-Point Stencil to a vector"

Global-View

```
proc main() {
    var n = 1000;
    var A, B: [1..n] real;

    forall i in 2..n-1 do
        B[i] = (A[i-1] + A[i+1])/2;
}
```

Local-View (SPMD)

```
proc main() {
  var n = 1000;
  var p = numProcs(),
      me = myProc(),
      myN = n/p,
  var A, B: [0..myN+1] real;
  if (me < p-1) {
    send(me+1, A[myN]);
    recv (me+1, A[myN+1]);
  if (me > 0) {
    send (me-1, A[1]);
    recv (me-1, A[0]);
  forall i in 1..myN do
    B[i] = (A[i-1] + A[i+1])/2;
```

Bug: Refers to uninitialized values at ends of A



C O M P U T E

STORE

In code: "Apply a 3-Point Stencil to a vector"

Global-View

```
proc main() {
   var n = 1000;
   var A, B: [1..n] real;

  forall i in 2..n-1 do
   B[i] = (A[i-1] + A[i+1])/2;
}
```

Communication becomes geometrically more complex for higher-dimensional arrays

```
Local-View (SPMD)
```

```
proc main() {
                    Assumes p divides n
  var n = 1000;
  \mathbf{var} p = \mathbf{numProcs}(),
      me = myProc(),
      myN = n/p,
      myLo = 1,
      myHi = myN;
  var A, B: [0..myN+1] real;
  if (me < p-1) {
    send(me+1, A[myN]);
    recv(me+1, A[myN+1]);
  } else
    myHi = myN-1;
  if (me > 0) {
    send (me-1, A[1]);
    recv (me-1, A[0]);
  } else
    myLo = 2;
  forall i in myLo..myHi do
    B[i] = (A[i-1] + A[i+1])/2;
```



2) Global-View Programming: A Final Note



 A language may support both global- and local-view programming — in particular, Chapel does

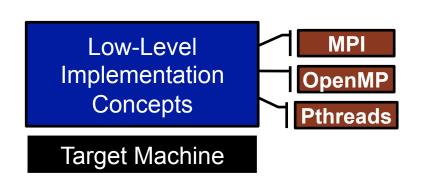
```
proc main() {
  coforall loc in Locales do
    on loc do
       MySPMDProgram(loc.id, Locales.numElements);

proc MySPMDProgram(myImageID, numImages) {
    ...
}
```



3) Multiresolution Design: Motivation





HPF High-Level Abstractions

Target Machine

"Why is everything so tedious/difficult?"

"Why don't my programs port trivially?"

"Why don't I have more control?"



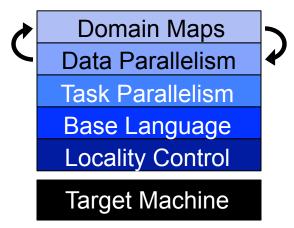
3) Multiresolution Design: Concept



Multiresolution Design: Support multiple tiers of features

- higher levels for programmability, productivity
- lower levels for greater degrees of control

Chapel language concepts



- build the higher-level concepts in terms of the lower
- permit the user to intermix layers arbitrarily



ANALYZE

4) Control over Locality/Affinity



Consider:

- Scalable architectures package memory near processors
- Remote accesses take longer than local accesses

Therefore:

- Placement of data relative to tasks affects scalability
- Give programmers control of data and task placement

Note:

 Increasingly, locality matters more and more within a compute node as well

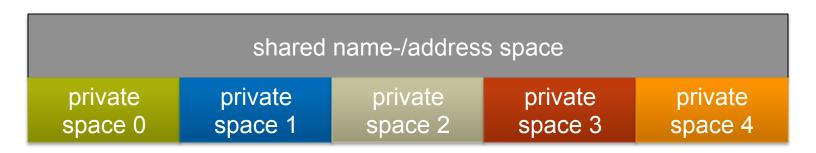


Partitioned Global Address Space Languages

(Or perhaps: partitioned global namespace languages)

abstract concept:

- support a shared namespace on distributed memory
 - permit any parallel task to access any lexically visible variable
 - doesn't matter if it's local or remote





ANALYZE

Partitioned Global Address Space Languages

(Or perhaps: partitioned global namespace languages)

abstract concept:

- support a shared namespace on distributed memory
 - permit any parallel task to access any lexically visible variable
 - doesn't matter if it's local or remote
- establish a strong sense of ownership
 - every variable has a well-defined location
 - local variables are cheaper to access than remote ones

	partitioned sh	nared name-/a	ddress space	
private	private	private	private	private
space 0	space 1	space 2	space 3	space 4



COMPUTE

Chapel and PGAS



Chapel is a PGAS language...

...but unlike most, it's not restricted to SPMD

⇒ never think in terms of "the other copies of the program"



Locales (think: "compute nodes")



C O M P U T E

5) Reduce HPC ↔ Mainstream Language Gap



Consider:

- Students graduate with training in Java, Matlab, Python, etc.
- Yet HPC programming is dominated by Fortran, C/C++, MPI

We'd like to narrow this gulf with Chapel:

- to leverage advances in modern language design
- to better utilize the skills of the entry-level workforce...
- ...while not alienating the traditional HPC programmer
 - e.g., support object-oriented programming, but make it optional





Questions about Context and Motivation?





COMPUTE | STORE | ANALYZE

Legal Disclaimer

Information in this document is provided in connection with Cray Inc. products. No license, express or implied, to any intellectual property rights is granted by this document.

Cray Inc. may make changes to specifications and product descriptions at any time, without notice.

All products, dates and figures specified are preliminary based on current expectations, and are subject to change without notice.

Cray hardware and software products may contain design defects or errors known as errata, which may cause the product to deviate from published specifications. Current characterized errata are available on request.

Cray uses codenames internally to identify products that are in development and not yet publically announced for release. Customers and other third parties are not authorized by Cray Inc. to use codenames in advertising, promotion or marketing and any use of Cray Inc. internal codenames is at the sole risk of the user.

Performance tests and ratings are measured using specific systems and/or components and reflect the approximate performance of Cray Inc. products as measured by those tests. Any difference in system hardware or software design or configuration may affect actual performance.

The following are trademarks of Cray Inc. and are registered in the United States and other countries: CRAY and design, SONEXION, URIKA, and YARCDATA. The following are trademarks of Cray Inc.: ACE, APPRENTICE2, CHAPEL, CLUSTER CONNECT, CRAYPAT, CRAYPORT, ECOPHLEX, LIBSCI, NODEKARE, THREADSTORM. The following system family marks, and associated model number marks, are trademarks of Cray Inc.: CS, CX, XC, XE, XK, XMT, and XT. The registered trademark LINUX is used pursuant to a sublicense from LMI, the exclusive licensee of Linus Torvalds, owner of the mark on a worldwide basis. Other trademarks used in this document are the property of their respective owners.

Copyright 2014 Cray Inc.



COMPUTE | STORE

