

# Chapel: Base Language

# Goals of this Talk

- Help you understand code in subsequent slide decks
- Give you the basic skills to program in Chapel
- Provide a survey of Chapel's base language features
- Impart an appreciation for the base language design

*Note: There is more in this slide deck than we will be able to cover, so consider it a reference and overview*



# "Hello World" in Chapel: Two Versions

- Fast prototyping

```
writeln("Hello, world!");
```

- “Production-grade”

```
module Hello {  
    proc main() {  
        writeln("Hello, world!");  
    }  
}
```



# Characteristics of Chapel

- **Design points**

- Identifying parallelism & locality is user's job, not compiler's
- No unexpected compiler-inserted array temporaries
- No pointers and limited opportunities for aliasing
- Intentionally not an extension of an existing language



# Chapel Influences

**C, Modula:** basic syntax

**ZPL, HPF:** data parallelism, index sets, distributed arrays

**CRAY MTA C/Fortran:** task parallelism, synchronization

**CLU** (see also Ruby, Python, C#): iterators

**Scala** (see also ML, Matlab, Perl, Python, C#): type inference

**Java, C#:** OOP, type safety

**C++:** generic programming/templates  
(but with a different syntax)



# Outline

- Introductory Notes
- Elementary Concepts
  - Lexical structure
  - Types, variables, and constants
  - Operators and Assignments
  - Compound Statements
  - Input and output
- Data Types and Control Flow
- Program Structure



# Lexical Structure

- Comments

```
/* standard  
C style  
multi-line */
```

```
// standard C++ style single-line
```

- Identifiers:

- Composed of A-Z, a-z, \_, \$, 0-9
- Cannot start with 0-9
- Case-sensitive



# Primitive Types

Type	Description	Default Value	Currently-Supported Bit Widths	Default Bit Width
bool	logical value	false	8, 16, 32, 64	impl. dep.
int	signed integer	0	8, 16, 32, 64	64
uint	unsigned integer	0	8, 16, 32, 64	64
real	real floating point	0.0	32, 64	64
imag	imaginary floating point	0.0i	32, 64	64
complex	complex floating points	0.0 + 0.0i	64, 128	128
string	character string	""	N/A	N/A

## • Syntax

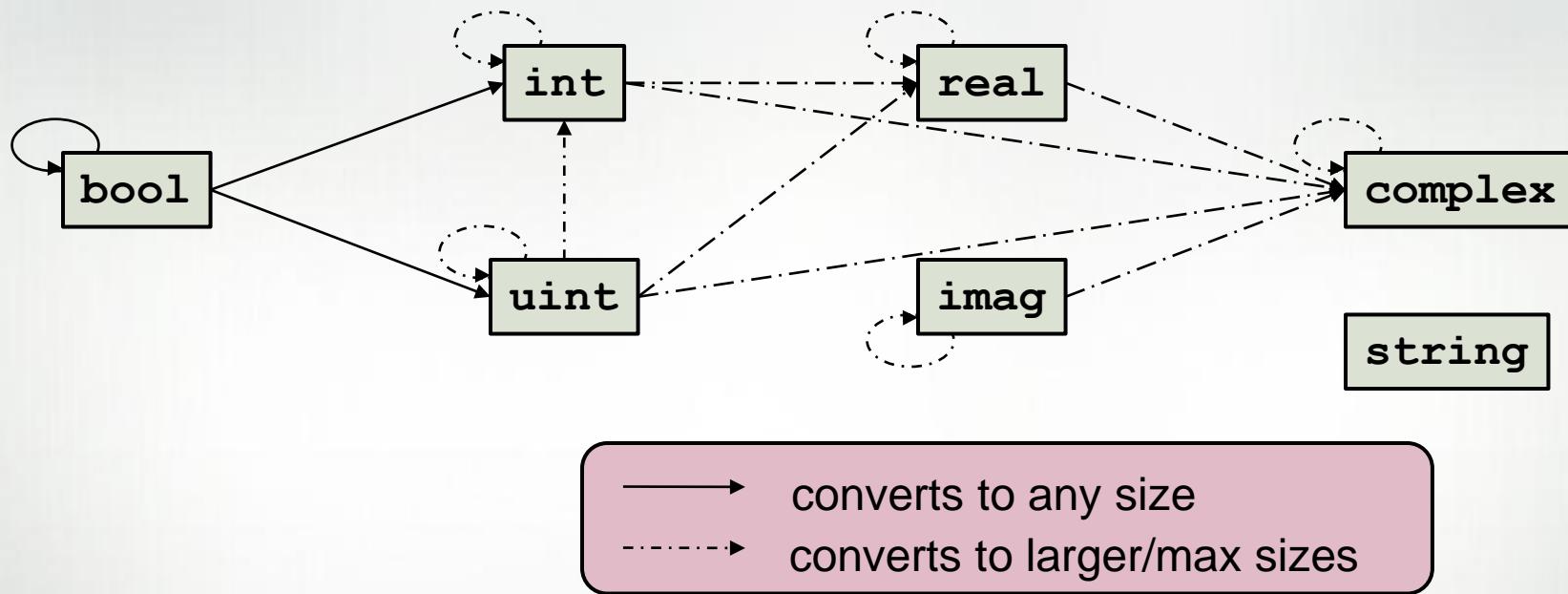
```
primitive-type:  
    type-name [ ( bit-width ) ]
```

## • Examples

```
int(16) // 16-bit int  
real(32) // 32-bit real  
uint      // 64-bit uint
```



# Implicit Type Conversions (Coercions)



- Notes:
  - reals do not implicitly convert to ints as in C
  - ints and uints don't interconvert as handily as in C

# Type Aliases and Casts

- Basic Syntax

```
type-alias-declaration:  
    type identifier = type-expr;  
  
cast-expr:  
    expr : type-expr
```

- Semantics

- type aliases are simply symbolic names for types
- casts are supported between any primitive types

- Examples

```
type elementType = complex(64);  
  
5:int(8) // store value as int(8) rather than int  
"54":int // convert string to an int  
249:elementType // convert int to complex(64)
```



# Variables, Constants, and Parameters

## • Basic syntax

*declaration:*

```
var identifier [: type] [= init-expr];
const identifier [: type] [= init-expr];
param identifier [: type] [= init-expr];
```

## • Semantics

- **var/const**: execution-time variable/constant
- **param**: compile-time constant
- No *init-expr* ⇒ initial value is the type's default
- No *type* ⇒ type is taken from *init-expr*

## • Examples

```
const pi: real = 3.14159;
var count: int;                      // initialized to 0
param debug = true;                  // inferred to be bool
```



# Config Declarations

- ## Syntax

```
config-declaration:  
  config type-alias-declaration  
  config declaration
```

- ## Semantics

- Like normal, but supports command-line overrides
- Must be declared at module/file scope

- ## Examples

```
config param intSize = 32;  
config type elementType = real(32);  
config const epsilon = 0.01:elementType;  
config var start = 1:int(intSize);
```

```
% chpl myProgram.chpl -sintSize=64 -selementType=real  
% a.out --start=2 --epsilon=0.00001
```



# Basic Operators and Precedence

Operator	Description	Associativity	Overloadable
:	cast	<b>left</b>	no
<b>**</b>	exponentiation	<b>right</b>	yes
<b>! ~</b>	logical and bitwise negation	<b>right</b>	yes
<b>* / %</b>	multiplication, division and modulus	<b>left</b>	yes
<i>unary + -</i>	positive identity and negation	<b>right</b>	yes
<b>+ -</b>	addition and subtraction	<b>left</b>	yes
<b>&lt;&lt; &gt;&gt;</b>	shift left and shift right	<b>left</b>	yes
<b>&lt;= &gt;= &lt; &gt;</b>	ordered comparison	<b>left</b>	yes
<b>== !=</b>	equality comparison	<b>left</b>	yes
<b>&amp;</b>	bitwise/logical and	<b>left</b>	yes
<b>^</b>	bitwise/logical xor	<b>left</b>	yes
<b> </b>	bitwise/logical or	<b>left</b>	yes
<b>&amp;&amp;</b>	short-circuiting logical and	<b>left</b>	via <code>isTrue</code>
<b>  </b>	short-circuiting logical or	<b>left</b>	via <code>isTrue</code>



# Assignments

Kind	Description
=	simple assignment
<code>+ = - = * = / = % =</code> <code>* *= &amp; =   = ^ =</code> <code>&amp; &amp; =    = &lt;&lt;= &gt;&gt; =</code>	compound assignment (e.g., <code>x += y;</code> is equivalent to <code>x = x + y;</code> )
<code>&lt;=&gt;</code>	swap assignment

- Note: assignments are only supported at the statement level



# Compound Statements

- **Syntax**

```
compound-stmt:  
{ stmt-list }
```

- **Semantics**

- As in C, permits a series of statements to be used in place of a single statement

- **Example**

```
{  
    writeln("Starting a compound statement");  
    x += 1;  
    writeln("Ending the compound statement");  
}
```



# Console Input/Output

- **Output**

- `write(expr-list)`: writes the argument expressions
- `writeln(...)` variant: writes a linefeed after the arguments

- **Input**

- `read(expr-list)`: reads values into the argument expressions
- `read(type-list)`: reads values of given types, returns as tuple
- `readln(...)` variant: same, but reads through next linefeed

- **Example:**

```
var first, last: string;
write("what is your name? ");
read(first);
last = read(string);
writeln("Hi ", first, " ", last);
```

What is your name?  
**Chapel User**  
Hi Chapel User

- I/O to files and strings also supported



# Outline

- Introductory Notes
- Elementary Concepts
- Data Types and Control Flow
  - Tuples
  - Ranges
  - Arrays
  - For loops
  - Other control flow
- Program Structure



# Tuples

## • Syntax

```
heterogeneous-tuple-type:  
( type, type-list )
```

```
homogenous-tuple-type:  
param-int-expr * type
```

```
tuple-expr:  
( expr, expr-list )
```

## • Purpose

- supports lightweight grouping of values  
(e.g., when passing or returning procedure arguments)
- multidimensional arrays use tuple indices

## • Examples

```
var coord: (int, int, int) = (1, 2, 3);  
var coordCopy: 3*int = coord;  
var (i1, i2, i3) = coord;  
var triple: (int, string, real) = (7, "eight", 9.0);
```



# Range Values

- Syntax

```
range-expr:  
[low] .. [high]
```

- Semantics

- Regular sequence of integers

$low \leq high$ :  $low, low+1, low+2, \dots, high$

$low > high$ : degenerate (an empty range)

$low$  or  $high$  unspecified: unbounded in that direction

- Examples

```
1..6          // 1, 2, 3, 4, 5, 6  
6..1          // empty  
3..          // 3, 4, 5, 6, 7, ...
```



# Range Operators

## • Syntax

*range-op-expr:*

*range-expr* **by** *stride*

*range-expr* **#** *count*

*range-expr* **align** *alignment*

*range-expr*[*range-expr*]

## • Semantics

- **by**: strides range; negative *stride*  $\Rightarrow$  start from *high*
- **#**: selects initial *count* elements of range
- **align**: specifies the alignment of a strided range
- [] or (): intersects the two ranges

## • Examples

```
1..6 by 2    // 1, 3, 5
1..6 by -1   // 6, 5, 4, ..., 1
1..6 #4       // 1, 2, 3, 4
1..6[3..]     // 3, 4, 5, 6
```

```
1.. by 2      // 1, 3, 5, ...
1.. by 2 #3  // 1, 3, 5
1.. by 2 align 2 // 2, 4, ...
1.. #3 by 2  // 1, 3
0..#n        // 0, ..., n-1
```



# Array Types

- ## Syntax

```
array-type:  
  [ index-set-expr ] elt-type
```

- ## Semantics

- Stores an element of *elt-type* for each index
- Array values expressed using square brackets

- ## Examples

```
var A: [1..3] int = [5, 3, 9], // 3-element array of ints
    B: [1..3, 1..5] real,      // 2D array of reals
    C: [1..3][1..5] real;     // array of arrays of reals
```

*Much more on arrays in data parallelism section later...*



# For Loops

- Syntax

```
for-loop:
```

```
  for index-expr in iterable-expr { stmt-list }
```

- Semantics

- Executes loop body serially, once per loop iteration
- Declares new variables for identifiers in *index-expr*
  - type and const-ness determined by *iterable-expr*
  - *iterable-expr* could be a range, array, or iterator

- Examples

```
var A: [1..3] string = [" DO", " RE", " MI"];
```

```
for i in 1..3 { write(A(i)); }           // DO RE MI
for a in A { a += "LA"; } write(A);    // DOLA RELA MILA
```



# Zipper Iteration

- Syntax

```
zipper-for-loop:
```

```
  for index-expr in zip( iteratable-exprs ) { stmt-list }
```

- Semantics

- Zipper iteration is over all yielded indices pair-wise

- Example

```
var A: [0..9] real;  
  
for (a,i,j) in zip(A, 1..10, 2..20 by 2) do  
  a = j + i/10.0;  
  
writeln(A);
```

```
2.1 4.2 6.3 8.4 10.5 12.6 14.7 16.8 18.9 21.0
```



# Other Control Flow Statements

- Conditional statements

```
if cond { computeA(); } else { computeB(); }
```

- While loops

```
while cond {  
    compute();  
}
```

```
do {  
    compute();  
} while cond;
```

- Select statements

```
select key {  
    when value1 { compute1(); }  
    when value2 { compute2(); }  
    otherwise { compute3(); }  
}
```

**Note:** Chapel also has expression-level conditionals and for loops



# Control Flow: Braces vs. Keywords

Most control flow supports keyword-based forms for single-statement versions

- Conditional statements

```
if cond then computeA(); else computeB();
```

- While loops

```
while cond do  
    compute();
```

- For loops

```
for indices in iterable-expr do  
    compute();
```

- Select statements

```
select key {  
    when value1 do compute1();  
    when value2 do compute2();  
    otherwise   do compute3();  
}
```



# Outline

- Introductory Notes
- Elementary Concepts
- Data Types and Control Flow
- Program Structure
  - Procedures and iterators
  - Modules and main()
  - Records and classes
  - Generics
  - Other basic language features



# Procedures, by example

- Example to compute the area of a circle

```
proc area(radius: real): real {
    return 3.14 * radius**2;
}

writeln(area(2.0)); // 12.56
```

```
proc area(radius) {
    return 3.14 * radius**2;
}
```

Argument and return types can be omitted

- Example of argument default values, naming

```
proc writeCoord(x: real = 0.0, y: real = 0.0) {
    writeln((x,y));
}

writeCoord(2.0);           // (2.0, 0.0)
writeCoord(y=2.0);         // (0.0, 2.0)
writeCoord(y=2.0, 3.0);    // (3.0, 2.0)
```



# Iterators

- **Iterator:** a procedure that generates values/variables
  - Used to drive loops or populate data structures
  - Like a procedure, but yields values back to invocation site
  - Control flow logically continues from that point
- Example

```
iter fibonacci(n) {  
    var current = 0,  
        next = 1;  
    for 1..n {  
        yield current;  
        current += next;  
        current <=> next;  
    }  
}
```

```
for f in fibonacci(7) do  
    writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8
```



# Argument and Return Intents

- Arguments can optionally be given intents
  - (blank): varies with type; follows principle of least surprise
    - most types: **const**
    - arrays, domains, sync vars: passed by reference
  - **const**: disallows modification of the formal
  - **in**: copies actual into formal at start; permits modifications
  - **out**: copies formal into actual at procedure return
  - **inout**: does both of the above
  - **ref**: pass by reference
  - **param/type**: formal must be a param/type (evaluated at compile-time)
- Return types can also have intents
  - (blank)/**const**: cannot be modified (without assigning to a variable)
  - **var**: permits modification back at the callsite
  - **type**: returns a type (evaluated at compile-time)
  - **param**: returns a param value (evaluated at compile-time)



# Modules

## • Syntax

```
module-def:  
  module identifier { code }  
  
module-use:  
  use module-identifier;
```

## • Semantics

- all Chapel code is stored in modules
- use-ing a module makes its symbols visible in that scope
- module-level statements are executed at program startup
  - typically used to initialize the module
- for convenience, a file containing code outside of a module declaration creates a module with the file's name



# Program Entry Point: main()

## • Semantics

- Chapel programs start by:
  - initializing all modules
  - executing main(), if it exists

```
M1.chpl:  
use M2;  
writeln("Initializing M1");  
proc main() { writeln("Running M1"); }
```

```
M2.chpl:  
module M2 {  
    writeln("Initializing M2");  
}
```

```
% chpl M1.chpl M2.chpl  
% ./a.out  
  
Initializing M2  
Initializing M1  
Running M1
```



# Revisiting "Hello World"

- Fast prototyping

hello.chpl

```
writeln("Hello, world!");
```

==

```
module hello {  
    writeln("Hello, world!");  
}
```

- “Production-grade”

```
module HelloWorld {  
    proc main() {  
        writeln("Hello, world!");  
    }  
}
```

Module-level code is executed during module initialization

main() executed when program begins running



# Records and Classes

- Chapel's struct/object types
  - Contain variable definitions (fields)
  - Contain procedure & iterator definitions (methods)
  - Records: value-based (*e.g.*, assignment copies fields)
  - Classes: reference-based (*e.g.*, assignment aliases object)
  - Record : Class :: C++ struct : Java class
- Example

```
record circle {  
    var radius: real;  
    proc area() {  
        return pi*radius**2;  
    }  
}
```

```
var c1, c2: circle;  
c1 = new c1(radius=1.0);  
c2 = c1;           // copies c1  
c1.radius = 5.0;  
writeln(c2.radius); // 1.0  
// records deleted by compiler
```



# Records and Classes

- Chapel's struct/object types
  - Contain variable definitions (fields)
  - Contain procedure & iterator definitions (methods)
  - Records: value-based (*e.g.*, assignment copies fields)
  - Classes: reference-based (*e.g.*, assignment aliases object)
  - Record : Class :: C++ struct : Java class
- Example

```
class circle {  
    var radius: real;  
    proc area() {  
        return pi*radius**2;  
    }  
}
```

```
var c1, c2: circle;  
c1 = new c1(radius=1.0);  
c2 = c1; // aliases c1's circle  
c1.radius = 5.0;  
writeln(c2.radius); // 5.0  
delete c1; // users delete classes
```



# Method Examples

Methods without arguments need not use parenthesis

```
proc circle.circumference {  
    return 2* pi * radius;  
}  
  
writeln(c1.area(), " ", c1.circumference);
```

Methods can be defined for any type

```
proc int.square() {  
    return this**2;  
}  
  
writeln(5.square());
```



# Generic Procedures

Generic procedures can be defined using type and param arguments:

```
proc foo(type t, x: t) { ... }  
proc bar(param bitWidth, x: int(bitWidth)) { ... }
```

Or by simply omitting an argument type (or type part):

```
proc goo(x, y) { ... }  
proc sort(A: []) { ... }
```

Generic procedures are instantiated for each unique argument signature:

```
foo(int, 3);           // creates foo(x:int)  
foo(string, "hi");    // creates foo(x:string)  
goo(4, 2.2);          // creates goo(x:int, y:real)
```



# Generic Objects

Generic objects can be defined using type and param fields:

```
class Table { param size: int; var data: size*int; }
class Matrix { type eltType; ... }
```

Or by simply eliding a field type (or type part):

```
record Triple { var x, y, z; }
```

Generic objects are instantiated for each unique type signature:

```
// instantiates Table, storing data as a 10-tuple
var myT: Table(10);
// instantiates Triple as x:int, y:int, z:real
var my3: Triple(int, int, real) = new Triple(1, 2, 3.0);
```



# Other Base Language Features not covered today

- Enumerated types
- Unions
- Type select statements, argument type queries
- Parenthesis-less functions/methods
- Procedure dispatch constraints (“where” clauses)
- Compile-time features for meta-programming
  - type/param procedures
  - folded conditionals
  - unrolled for loops
  - user-defined compile-time warnings and errors



# Status: Base Language Features

- Most features working well
- Performance is currently suboptimal in some cases
- Some semantic checks are incomplete
  - e.g., constness-checking for members, arrays
- Error messages could use improvement at times
- OOP features are limited in certain respects
  - generic classes w/ subclassing, user constructors
- Memory for strings is currently leaked



# Future Directions

- Error handling/Exceptions
- Fixed-length strings
- Interfaces (joint work with CU Boulder)
- Improved namespace control
  - private fields/methods in classes and records
  - module symbol privacy, filtering, renaming
- Interoperability with other languages (joint with LLNL)



# Questions?

- Introductory Notes
  - Characteristics
  - Influences
- Elementary Concepts
  - Lexical structure
  - Types, variables, and constants
  - Operators and assignments
  - Compound Statements
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- Data Types and Control Flow
  - Tuples
  - Ranges
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- Program Structure
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