

Chapel: Sample Codes



Outline



STREAM and RA HPC Challenge Benchmarks

- simple, regular 1D computations
- results from SC '09 competition

AMR Computations

hierarchical, regular computation

• SSCA #2

unstructured graph computation



HPC Challenge



Two classes of competition:

- Class 1: "best performance"
- Class 2: "most productive"
 - Judged on: 50% performance 50% elegance
 - Four recommended benchmarks: STREAM, RA, FFT, HPL
 - Use of library routines: discouraged

• Why you may care:

• provides an alternative to the top-500's focus on peak performance

• Recent Class 2 Winners:

 2008: performance: IBM (UPC/X10) productive: Cray (Chapel), IBM (UPC/X10), Mathworks (Matlab)
 2009: performance: IBM (UPC+X10) elegance: Cray (Chapel)



Chapel Implementation Characteristics







HPC Challenge: Chapel Entries (2008-2009)



Benchmark	2008	2009	Improvement
Global STREAM	1.73 TB/s (512 nodes)	10.8 TB/s (2048 nodes)	6.2x
EP STREAM	1.59 TB/s (256 nodes)	12.2 TB/s (2048 nodes)	7.7x
Global RA	0.00112 GUPs (64 nodes)	0.122 GUPs (2048 nodes)	109x
Global FFT	single-threaded single-node	multi-threaded multi- node	multi-node parallel
Global HPL	single-threaded single-node	multi-threaded single- node	single-node parallel

All timings on ORNL Cray XT4:

- 4 cores/node
- 8 GB/node
- no use of library routines





Global STREAM Triad in Chapel (Excerpts)





EP STREAM Triad in Chapel (Excerpts)





STREAM Triad Performance



Performance of HPCC STREAM Triad (Cray XT4)





Global Random Access in Chapel (Excerpts)

```
const TableDist = new dmap(new Block([0..m-1])),
UpdateDist = new dmap(new Block([0..N U-1]));
```

const TableSpace: domain ... dmapped TableDist = ..., Updates: domain ... dmapped UpdateDist = ...;

```
var T: [TableSpace] uint(64);
```

Chapel: Sample Codes

```
forall ( ,r) in (Updates,RAStream()) do
  on TableDist.idxToLocale(r & indexMask) {
    const myR = r;
    local T(myR & indexMask) ^= myR;
  }
  This body should eventually simply be written:
    on T(r&indexMask) do
    T(r&indexMask) ^= r;
  (and again, can be today, but performance is worse)
```



Random Access Performance





Random Access Efficiency on 32+ Nodes





Portability Results





Performance of HPCC Random Access (Cray CX1) 0.02 0.016 0.012

GUP/s



Performance of HPCC STREAM Triad (IBM pSeries 575)







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Adaptive Mesh Refinement in Chapel

What's so great about domains?

- Ability to reason about unions of rectangular index spaces (as unions of domains)
- Trivial shared-memory parallelism; easy access to distributed parallelism with distributions
- Fewer nested loops, and no bounds to mess up
- Striding allows much better description of grids (vertices, edges, cell centers)
- Dimension-independent code









Chapel (15)



var cell_centers = [1..7 by 2, 1..5 by 2];



Chapel (16)



- var cell_centers = [1..7 by 2, 1..5 by 2];
- var vertical edges = [0..8 by 2, 1..5 by 2];





- var cell centers = [1..7 by 2, 1..5 by 2];
- var vertical edges = [0..8 by 2, 1..5 by 2];

var horizontal edges = [1..7 by 2, 0..6 by 2];



Chapel (18)



- var cell centers = [1..7 by 2, 1..5 by 2];
- **var** vertical edges = [0..8 **by** 2, 1..5 **by** 2];
- **var** horizontal edges = [1..7 **by** 2, 0..6 **by** 2];

var vertices = [0..8 by 2, 0..6 by 2];





Dimension-free stencils

Tasks:

- 1. Create an *N*-dimensional grid.
- 2. Evaluate the function $f(x_1, x_2, \ldots, x_N) = \sin(x_1) \sin(x_2) \cdots \sin(x_N)$ on the grid.
- 3. Approximate the Laplacian,

$$\Delta f = f_{x_1 x_1} + f_{x_2 x_2} + \dots + f_{x_N x_N}$$





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config param N: int = 2; const dimensions = [1..N];







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```
config param N: int = 2;
const dimensions = [1..N];
```

config const num_points = 20; const dx = 1.0 / (num_points-1);

```
var grid_points: domain(N);
```

```
var ranges: N*range;
for d in dimensions do
  ranges(d) = 1..num points;
```

```
grid_points = ranges;
```







2. Evaluate the function

```
f(x_1, x_2, \dots, x_N) = \sin(x_1)\sin(x_2)\cdots\sin(x_N)
```

```
var f: [grid points] real = 1.0;
```

```
forall point in points {
   for d in dimensions do
      f(point) *= sin( (point(d)-1)*dx );
}
```





2. Evaluate the function

```
f(x_1, x_2, \ldots, x_N) = \sin(x_1)\sin(x_2)\cdots\sin(x_N)
```

```
var f: [grid_points] real = 1.0;
forall point in points {
  for d in dimensions do
    f(point) *= sin( (point(d)-1)*dx );
}
Calculates real
```

coordinate \boldsymbol{x}_d







$$\Delta f = f_{x_1 x_1} + f_{x_2 x_2} + \dots + f_{x_N x_N}$$









$$\Delta f = f_{x_1 x_1} + f_{x_2 x_2} + \dots + f_{x_N x_N}$$

var interior_points = grid_points.expand(-1);
var laplacian: [interior_points] real;

Laplacian is only defined on the **interior** of the grid







$$\Delta f = f_{x_1 x_1} + f_{x_2 x_2} + \dots + f_{x_N x_N}$$

var interior_points = grid_points.expand(-1);
var laplacian: [interior_points] real;

```
forall point in interior_points {
    var shift: N*int;
```





}



$$\Delta f = f_{x_1 x_1} + f_{x_2 x_2} + \dots + f_{x_N x_N}$$

```
forall point in interior_points {
    var shift: N*int;
```





$$\Delta f = f_{x_1 x_1} + f_{x_2 x_2} + \dots + f_{x_N x_N}$$

```
forall point in interior points {
       var shift: N*int;
       for d in dimensions {
         shift(d) = 1;
         laplacian(point) += ( f(point+shift)
                               -2*f(point)
Translates in
                               + f(point-shift)
dimension d
                              ) / dx^{*2};
        shift(d) = 0;
```





Chapel (31)

$$\Delta f = f_{x_1 x_1} + f_{x_2 x_2} + \dots + f_{x_N x_N}$$

```
forall point in interior points {
       var shift: N*int;
       for d in dimensions {
         shift(d) = 1;
                                                        point
         laplacian(point) += ( f(point+shift)
                               -2*f(point)
Translates in
                               + f(point-shift)
dimension d
                              ) / dx^{*2};
        shift(d) = 0;
```



$$\Delta f = f_{x_1 x_1} + f_{x_2 x_2} + \dots + f_{x_N x_N}$$

```
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       for d in dimensions {
         shift(d) = 1;
                                                         point
         laplacian(point) += ( f(point+shift)
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$$\Delta f = f_{x_1 x_1} + f_{x_2 x_2} + \dots + f_{x_N x_N}$$



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Given a set of heavy edges *HeavyEdges* in directed graph *G*, find sub-graphs of outgoing paths with *length* ≤ *maxPathLength*







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maxPathLength = 0 maxPathLength = 1 maxPathLength = 2





def rootedHeavySubgraphs(
 G,
 type vertexSet;
 HeavyEdges : domain,
 HeavyEdgeSubG : [],
 in maxPathLength: int) {
 forall (e, subgraph) in
 (HeavyEdges, HeavyEdgeSubG) {
 const (x,y) = e;
 var ActiveLevel: vertexSet;
 }
}

ActiveLevel += y;

subgraph.edges += e; subgraph.nodes += x; subgraph.nodes += y; for pathLength in 1..maxPathLength {
 var NextLevel: vertexSet;
 forall v in ActiveLevel do
 forall w in G.Neighbors(v) do
 atomic {
 if !subgraph.nodes.member(w) {
 NextLevel += w;
 subgraph.nodes += w;
 subgraph.edges += (v, w);
 }
}

if (pathLength < maxPathLength) then
ActiveLevel = NextLevel;</pre>







<pre>def rootedHeavySubgraphs(for pa G, type vertexSet; f</pre>	thLength in 1maxPathLength { ar NextLevel: vertexSet; forall v in ActiveLevel do
Generic with respect to vertex sets vertexSet: A type argument specifying how to represent vertex subsets	<pre>forall w in G.Neighbors(v) do atomic { if !subgraph.nodes.member(w) { NextLevel += w; subgraph.nodes += w; } }</pre>
Requirements:Parallel iterationAbility to add members, test for membership	<pre>subgraph.edges += (v, w); }</pre>
 Options: An associative domain over vertices domain (index (G.vertices)) A sparse subdomain of the vertices sparse subdomain (G.vertices) 	<pre>f (pathLength < maxPathLength) then ActiveLevel = NextLevel;</pre>



def rootedHeavySubgraphs(G,

type vertexSet;

HeavyEdges : domain,

HeavyEdgeSubG : [],

in maxPathLength: int) {

forall (e, subgraph) in
 (HeavyEdges, HeavyEdgeSubG) {

const (x, y) = e;

```
var ActiveLevel: vertexSet;
```

for pathLength in 1..maxPathLength {
 var NextLevel: vertexSet;
 forall v in ActiveLevel do
 forall w in G.Neighbors(v) do
 atomic {

if !subgraph.nodes.member(w) {
 NextLevel += w;

subgraph.nodes += w;

subgraph.edges += (v, w);

ActiveLe

The same genericity applies to subgraphs

subgraph.edges += e; subgraph.nodes += x; subgraph.nodes += y;

if (pathLength < maxPathLength) then
ActiveLevel = NextLevel;</pre>





def rootedHeavySubgraphs(
 G,
 type vertexSet;
 HeavyEdges : domain,
 HeavyEdgeSubG : [],
 in maxPathLength: int) {
 forall (e, subgraph) in
 (HeavyEdges, HeavyEdgeSubG) {
 const (x,y) = e;
 var ActiveLevel: vertexSet;
 }
}

ActiveLevel += y;

subgraph.edges += e; subgraph.nodes += x; subgraph.nodes += y; for pathLength in 1..maxPathLength {
 var NextLevel: vertexSet;
 forall v in ActiveLevel do
 forall w in G.Neighbors(v) do
 atomic {
 if !subgraph.nodes.member(w) {
 NextLevel += w;
 subgraph.nodes += w;
 subgraph.edges += (v, w);
 }
 }
}

if (pathLength < maxPathLength) then
ActiveLevel = NextLevel;</pre>



Questions?



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