Chapel: Base Language

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SC08: Tutorial S07 – 11/16/08

Base Language Themes

- **Style:** Block-structured, imperative
- **Syntax:**
  - Borrow heavily from C family (C, C++, Java, C#, Perl) for familiarity
  - In other cases, use something intuitive and easy to learn
- **Object-oriented programming:**
  - Support it, but don't require it
  - Reference- and value-based objects (Java- and C++-style)
- **Type System:**
  - statically typed for performance, safety
  - permit types to be elided in most contexts for convenience
- **Aliasing:**
  - minimize aliases to help with compiler analysis (*e.g.*, no pointers)
  - main sources: object references, array aliases
- **Compiler-inserted array temporaries:** never require them

Chapel: Base Language (2)

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Chapel Influences

- Intentionally not an extension to an existing language
- Instead, select attractive features from previous work:

  ZPL, HPF: data parallelism, index sets, distributed arrays
  (see also APL, NESL, Fortran90)

  Cray MTA C/Fortran: task parallelism, lightweight synchronization

  CLU: iterators (see also Ruby, Python, C#)

  ML: latent types (see also Scala, Matlab, Perl, Python, C#)

  Java, C#: OOP, type safety

  C++: generic programming/templates (without adopting its syntax)

  C, Modula, Ada: syntax

Outline

- Starting points
- Basics
  - Lexical Structure
  - Scalar Types
  - Variable, Constant, Configuration Declarations
  - Console I/O
  - Conversions
  - Operators
- Middle Ground
- More advanced topics

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Lexical Structure

- Comments: standard C-style comments

```
x = 1;   // single-line comment
x = 2;   /* multi-line
            comment */
```

- Whitespace:
  - spaces, TABs, new-lines
  - ignored, except to separate tokens and end single-line comments
- Identifiers:
  - made up of `A-Z`, `a-z`, `0-9`, `_`, `$`
  - cannot start with `0-9`
- Case-sensitivity: Chapel *is* case-sensitive
- Statement structure:
  - statements terminated by `;`
  - compound statements enclosed by `{ … }`

Scalar Types

	description	default value	default width	currently supported bit-widths
bool	boolean value	false	impl.-dependent	8, 16, 32, 64
int	signed integer	0	32 bits	8, 16, 32, 64
uint	unsigned integer	0	32 bits	8, 16, 32, 64
real	real floating point	0.0	64 bits	32, 64
imag	imaginary floating point	0.0i	64 bits	32, 64
complex	complex value	0.0 + 0.0i	128 bits	64, 128
string	character string	“”	N/A	N/A

- Syntax

```
scalar-type:
  scalar-type-name [(width)]
```

- Examples

```
int(64)    // a 64-bit int
real(32)   // a 32-bit real
uint       // a 32-bit uint
```


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Literals

- Boolean:

```
true    // true bool
false   // false bool
```

- Integer:

```
123      // decimal int
0x1fff   // hexidecimal int
0b1001   // binary int
```

- Floating Point:

```
1.2       // real
3.4e-6    // real
7.8i      // imag
```

```
100i                   // imag
1.2 + 3.4i             // complex
(5.6, 7.8): complex    // complex
```

- String:

```
"hi"    // string
'SC08'  // string
```


Declarations: Variables

- Syntax

```
var-decl-stmt:
  var identifier [: type] [= initializer]
```

- Semantics
  - declares a new variable named *identifier*
  - *type* – if specified, indicates variable's type
    – otherwise, type is inferred from *initializer*
  - *initializer* – if specified, used as the variable's initial value
    – otherwise, the variable's type determines its initial value
- Examples

```
var epsilon: real = 0.01;
var count: int;              // "count" initialized to 0
var name = "Brad";           // "name" inferred to be string
var x, y: int,               // comma-separated forms also
    flag = false;            //   supported
```


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Console Output

- Syntax:

```
write(expr-list)
writeln(expr-list)
```

- Semantics:
  - *write* – print the argument list to the console in order
  - *writeln* – same as write, but also print a new-line at the end
- Examples:

```
var n = 1000;
writeln("n is: ", n);
```

- Output:

```
n is 1000
```


Hello world: simplest version

- Program

```
writeln("Hello, world!");
```

- Output

```
Hello, world!
```


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Console Input

- Syntax (readln versions also supported):

```
read(expr-list)
read(type)
read(type-list)
```

- Semantics:
  - *read(expr-list)* – read values into the argument list expressions
  - *read(type)* – read a value of the specified type and return it
  - *read(type-list)* – read values of the given types and return as a tuple
  - *readln* – same as read, but then read through the next new-line
- Examples:

```
var x, y: real,
    z: int;
read(x, z);                // read a real into x, an int into z
y = read(real);            // read a real into y
(y, z) = read(real, int); // read a real into y, an int into z
```


Declarations: Constants

- Syntax

```
const-decl-stmt:
  const identifier [: type] [= initializer]
```

- Semantics
  - like a variable, but cannot be reassigned after initialization
  - initializer need not be a statically-known value
- Examples

```
const pi = 3.14159;    // pi is a constant 64-bit real
pi = 0.0;              // ILLEGAL: cannot reassign a const
const n = computeN(); // can initialize w/ runtime value
```


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Configuration Variables/Constants

- Syntax

```
config-decl-stmt:
  config const-decl-stmt
| config var-decl-stmt
```

- Semantics
  - like a standard declaration, but supports command-line overrides
  - must be declared at global scope
- Examples

```
config const n = 10,
             epsilon = 0.01,
             verbose = false;
```

- Executable Command-line

```
> ./a.out --n=10000 --epsilon=0.0000001 --verbose=true
```


Hello world: configurable version

- Program

```
config const msg = "Hello, world!";

writeln(msg);
```

- Output

```
> ./a.out
Hello, world!
>
> ./a.out --msg="Hello, SC08!!"
Hello, SC08!!
```


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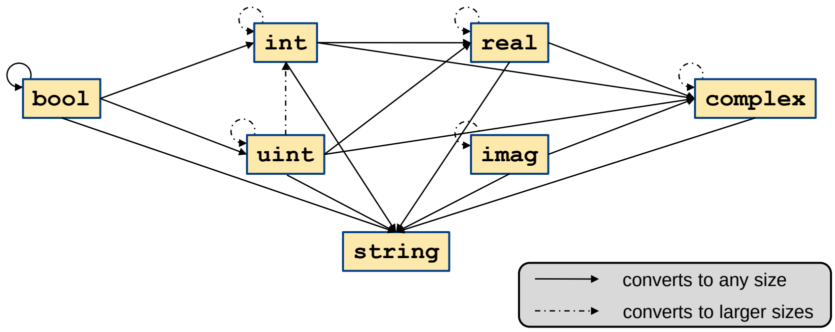
Implicit Conversions



- Notes:
  - reals do not implicitly convert to ints as in C
  - ints and uints don't interconvert as handily as in C
  - C# has served as our model in establishing these rules

Explicit Conversions / Casts

- Syntax

```
cast-expr:
  expr : type
```

- Semantics
  - convert *expr* to the type specified by *type*
- Examples

```
const three = pi: int,
      age   = "3": int;
```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Basic Operators

Operator	Description
`+ - * / %`	arithmetic (plus, minus, multiply, divide, C-style modulus)
`**`	exponentiation
`& | ^ ~ << >>`	bitwise (and, or, xor, not, shift-left, shift-right)
`&& || !`	logical, short-circuiting (and, or, not)
`=`	assignment
`+= -= *= /= %=` `**= &= |= ^=` `<<= >>= &&= ||=`	op-assignment (*e.g.*, `x += y;` ⇒ `x = x + y;`)
`<=>`	swap assignment

Outline

- Starting points
- Basics
- Middle Ground
  - Other Types
    - Ranges
    - Arrays
  - Loops and Control Flow
  - Program Structure
    - Functions and Iterators
    - Modules and `main()`
- More advanced topics

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Other Types

- Covered Today:
  - Ranges: regular integer sequences
  - Domains: index sets
  - Arrays: mappings from indices to variables
- Touched on Today:
  - Tuples: lightweight mechanism for grouping variables/values
  - Records: value-based objects, like C structs or C++ classes
  - Classes: reference-based objects, like Java or C# classes
- Not covered today:
  - Unions: store multiple types in overlapping memory
    - (as in C, but type-safe)
  - Enumerated types: finite list of named values
    - *e.g.*, `enum color {red, green, blue};`

Range Values

- Syntax

```
range-expr:
  [lo]..[hi] [by stride]
```

- Semantics
  - represents a regular sequence of integers
    - if *stride* > 0: *lo*, *lo+stride*, *lo*+2**stride*, … ≤ *hi*
    - if *stride* < 0: *hi*, *hi–stride*, *hi*–2**stride*, … ≥ *lo*
  - *lo* or *hi* can be omitted if *stride* has the appropriate sign
- Examples

```
1..6              // 1, 2, 3, 4, 5, 6
1..6 by -1        // 6, 5, 4, 3, 2, 1
6..1              // an empty sequence
1..6 by 2         // 1, 3, 5
1..6 by -4        // 6, 2
1..               // 1, 2, 3, 4, 5, 6, 7, 8, …
```


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The # operator

- Syntax

```
count-expr:
  range-expr # count-expr
```

- Semantics
  - creates a range from the initial *count-expr* elements of *range-expr*
- Examples

```
0..#6              // 0, 1, 2, 3, 4, 5
0..#6 by 2         // 0, 2, 4
0..by 2 #6         // 0, 2, 4, 6, 8, 10
1..6 #3            // 1, 2, 3
1..6 by -1 # 3     // 6, 5, 4
```


Array Types

- Syntax

```
array-type:
  [index-set] type
```

- Semantics
  - for each index in *index-set*, stores an element of *type*
- Examples

```
var A: [1..3] int,            // a 3-element array of ints
    B: [1..3, 1..5] string,   // a 2D 3x5 array of strings
    C: [1..3] [1..5] real;    // a 3-element array of
                              //   5-element arrays of reals

    D: [1..3] int = (1, 2, 3);  // initialized array
```


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Array Indexing

- Syntax

```
index-expr:
  array-expr[index-expr]
| array-expr(index-expr)
```

- Semantics
  - references the element in *array-expr* corresponding to *index-expr*
- Examples

```
var A: [1..3] int,
    B: [1..3, 1..5] string;

A(1) = 2;
A[2] = 4;
B(1, 2) = "hi";
B[2, 5] = "SC08";
B[0, 0] = "oops";   // error: indexing out-of-bounds
```


For loops

- Syntax

```
for-loop:
  for identifier in iteratable-expr do body-stmt
| for identifier in iteratable-expr { body }
```

- Semantics
  - executes loop body once per value yielded by *iteratable-expr*
  - stores each value in a body-local variable/const named *identifier*
- Examples

```
var A: [1..3] string = ("hi", "SC08", "!!");

for i in 1..3 do write(A(i));    // prints: hiSC08!!
for i in 1..3 do i += 1;    // illegal, ranges yield consts
for a in A {
  a += "-";  write(a);          // prints: hi-SC08-!!
}                          // A is now "hi-" "SC08-" "!!-")
```


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Zippered/Tensor Iteration

- Syntax

```
tensor-for-loop:
  for index-decl in [iter-expr1, iter-expr2, …] loop-body

zippered-for-loop:
  for index-decl in (iter-expr1, iter-expr2, …) loop-body
```

- Semantics
  - *tensor-for-loop* – iterates over all pairs of yielded elements
  - *zippered-for-loop* – iterates over yielded elements pair-wise
- Examples

```
for i in [0..1, 0..1] …  // i = (0,0); (0,1); (1,0); (1,1)

for i in (0..1, 0..1) …      // i = (0,0); (1,1)
for (x,y) in (0..1, 0..1) …  // x=0, y=0;  x=1, y=1
```


Other Control Flow

- While loops

```
while test-expr do body-stmt
while test-expr { body-stmts }
do { body-stmts } while test-expr;
```

- Conditional Statements and Expressions

```
if test-expr then true-stmt [else false-stmt]
if test-expr { true-stmts } [else { false-stmts }]
if test-expr then true-expr [else false-expr]
```

- Also…
  - select: a switch/case statement
  - break: break out of a loop (optionally labeled)
  - continue: skip to next iteration of a loop (optionally labeled)
  - return: return from a function
  - exit: exit the program
  - halt: exit the program due to an exceptional/error condition

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Function Definitions

- Syntax

```
function-decl-stmt:
  def identifier [(formal-list)] [: type] body
```

- Semantics
  - *identifier* – name of function being defined
  - *formal-list* – list of arguments (potentially empty)
  - *type* – if specified, specifies function's return type
    – otherwise, return type inferred from function body
  - *body* – specifies function's definition
- Examples

```
def square(x: real): real {
  return x**2;
}

const pi2 = square(pi);
```


Formal Arguments

- Syntax

```
formal-argument:
  intent identifier [: type] [= init]
```

- Semantics
  - *identifier* – name of formal argument
  - *intent* – how to pass the actual argument
  - *type* – if specified, specifies formal type; otherwise generic (inferred)
  - *init* – if specified, permits argument to be omitted at callsite
- Example

```
def label(x, name: string, end = ".\n") {
  writeln(name, " is ", x, end);        }
label(n, "n", " and ");
label(msg "msg");
```

- Output

```
n is 1000 and msg is Hello, SC08!.
```


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Named Argument Passing

- Arguments may be matched by name rather than position

```
def foo(x: int = 2, y:int = x) { … }

foo();           // equivalent to foo(2, 2);
foo(3);          // equivalent to foo(3, 3);
foo(y=3);        // equivalent to foo(2, 3);
foo(y=3, x=2);   // equivalent to foo(2, 3);
foo(y=3, 2);     // equivalent to foo(2, 3);
```


Argument Intents

- Syntax

```
intent:
  (blank) | const | in | out | inout
```

- Semantics
  - `in` – copy actual into formal at function start and permit modification
  - `out` – copy formal into actual at function return
  - `inout` – combination of "`in`" and "`out`"
  - `const` – varies with type
  - (blank) – varies with type; follows "principle of least surprise"

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Argument Intents and Types

	argument type			
intent	**scalar type**	**domain/array**	**record**	**class**
`in`	"copy in" : copy actual into formal at function start and permit modification			
`out`	"copy out" : copy formal into actual at function return			
`inout`	"copy in and out" : combination of in and out			
`const`	copy in but disallow modification	pass by reference and disallow modification	copy in and disallow modification	copy reference in and disallow modification to reference
(blank)	see `const`	pass-by-reference and permit modification	see `const`	see `const`

Motivation for Iterators

Given a program with a bunch of similar loops…

```
for i in 0..#m do
  for j in 0..#n do
    …A(i,j)…

…

for i in 0..#m do
  for j in 0..#n do
    …A(i,j)…

…
```

Consider the effort to convert them from RMO to CMO…

```
for j in 0..#n do
  for i in 0..#m do
    …A(i,j)…

…

for j in 0..#n do
  for i in 0..#m do
    …A(i,j)…

…
```

Or to tile the loops…

```
for jj in 0..#n by block do
  for ii in 0..#m by block do
    for j in jj.. min(m,jj+block)-1 do
      for i in ii..min(n,ii+block)-1 do
        …A(i,j)…

…

for jj in 0..#n by block do
  for ii in 0..#m by block do
    for j in jj.. min(m,jj+block)-1 do
      for i in ii..min(n,ii+block)-1 do
        …A(i,j)…

…
```


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Motivation for Iterators

Given a program with a bunch of similar loops…

```
for i in 0..#m do
  for j in 0..#n do
    …A(i,j)…
```

Consider the effort to convert them from RMO to CMO…

```
for j in 0..#n do
  for i in 0..#m do
    …A(i,j)…
```

Or to tile the loops…

```
for jj in 0..#n by block do
  for ii in 0..#m by block do
    for j in jj.. min(m,jj+block)-1 do
      for i in ii..min(n,ii+block)-1 do
        …A(i,j)…
```

…

Or to change the iteration order over the tiles…

Or to make them into fragmented loops for an MPI program…

Or to change the distribution of the work/arrays in that MPI program…

Or to label them as parallel for OpenMP or a vectorizing compiler…

Or to do *anything* that we do with loops all the time as a community…

We wouldn't program straight-line code this way, so why are we so tolerant of our lack of loop abstractions?

Iterators

- like functions, but *yield* a number of elements one-by-one:

```
def RMO() {
  for i in 0..#m do
    for j in 0..#n do
      yield (i,j);
```

```
def tiled(block) {
  for jj in 0..#n by block do
    for ii in 0..#m by block do
      for j in jj.. min(m,jj+block)-1 do
        for i in ii..min(n,ii+block)-1 do
          yield (i,j);
}
```

- can be used to drive for loops:

```
for (i,j) in RMO() do
  …A(ij)…
```

```
for (i,j) in tiled(block) do
  …A(ij)…
```

- as with functions…

  …one iterator can be redefined to change the behavior of many loops

  …a single invocation can be altered, or its arguments can be changed
- not necessarily any more expensive than raw, inlined loops

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Modules

- Syntax:

```
module-def:
  module { code }
```

```
module-use:
  use module-name;
```

- Semantics
  - all Chapel code is stored in modules
  - `use`-ing a module causes its symbols to be available from that scope
  - top-level code in a module is executed when the module is first used
- Example

```
module M {
  def foo() {
    writeln("Hi from M!");
  }
  writeln("Someone used M");
}
```

```
use M;
foo();
```

- Output

```
Someone used M
Hi from M!
```


Program Entry Point

- Semantics
  - Each module can define a function "*main*" to serve as an entry point
  - If a module does not define "*main*", its top-level code serves as main
  - If a program defines multiple "*mains*", compiler flags must specify one
- Example

```
module M1 {
  def main() {
    writeln("Running M1");
  }
}
```

```
module M2 {
  def main() {
    writeln("Running M2");
  }
}
```

- Output

```
> chpl --main-module=M1
>
> a.out
Running M1
>
> chpl --main-module=M2
>
> a.out
Running M2
```


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Hello world: structured version

- Program

```
module Hello {
  def main() {
    writeln("Hello, world!");
  }
}
```

- Output

```
Hello, world!
```


Hello world: simplest version

- Program

```
writeln("Hello, world!");
```

- Output

```
Hello, world!
```


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Outline

- Starting points
- Basics
- Middle Ground
- More advanced topics
  - Object-oriented Programming (OOP)
  - Compile-time machinery
  - Generics

Record Types

- Syntax

```
record-type-decl:
  record identifier { decl-list }
```

- Semantics
  - creates a record type named *identifier*
  - *decl-list* – defines member constants/variables, and methods
  - assignment copies members from one record to another
  - similar to C++ classes
- Example

```
record employee { var name: string, id: int; }
var e1: employee = new employee(name="Brad", id=12345),
    e2: employee;    // e2 defaults to name="", id=0)
e1 = e2;             // e2 is a distinct copy of e1
e2.name = "Joe";
writeln(e1.name);    // prints "Brad"
```


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Class Types

- Syntax

```
class-type-decl:
  class identifier { decl-list }
```

- Semantics
  - similar to records, but creates a reference type rather than a “struct”
  - assignment copies object reference, not members
  - similar to Java classes
- Example

```
class employee { var name: string, id: int; }
var e1: employee = new employee(name="Brad", id=12345),
    e2: employee;    // e2 defaults to nil
e1 = e2;             // e2 is an alias of e1
e2.name = "Joe";
writeln(e1.name);    // prints "Joe"
```


OOP Capabilities

- We won’t cover a number of standard OOP features today:
  - inheritance
  - shadowing members/fields
  - dynamic dispatch
  - point of instantiation
  - …

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Standard Methods

- Classes/records support standard user-defined methods:
  - `this()` – permits indexing an instance of the class/record
  - `these()` – permits iteration over an instance of the class/record
  - `writeThis()` – overrides the default way of printing a class
- Example uses:

```
var myC = new C();

myC(i,j);      // legal if C supports a this() function
               // that takes i and j as arguments

for x in myC do …   // legal if C supports a these()
                    // iterator

writeln(myC);   // calls writeThis() if defined,
                // otherwise compiler supplies a default
```


Standard Methods Example

```
class Pair {
  var x, y: real;

  def this(i: int) {
    if (i==0) then
      return x;
    if (i==1) then
      return y;
    halt("out-of-bounds: ", i);
  }

  def these() {
    yield x;
    yield y;
  }

  def writeThis(s: Writer) {
    s.write((x,y));
} }
```

- Use

```
var p = new Pair(x=1.2,
                 y=3.4);
writeln("p(0)=", p(0));
writeln("p(1)=", p(1));
p(0) = 5.6;  p(1) = 7.8;
for x in p {
  writeln(x);
  x = -x;        }
writeln(p);
```

- Output

```
p(0)=1.2
p(1)=3.4
1.2
3.4
(1.2, 3.4)
```


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Standard Methods Using `var` Return Types

```
class Pair {
  var x, y: real;
  def this(i: int) var {
    if (i==0) then
      return x;
    if (i==1) then
      return y;
    halt("out-of-bounds: ", i);
  }
  def these() var {
    yield x;
    yield y;
  }
  def writeThis(s: Writer) {
    s.write((x,y));
} }
```

- Use

```
var p = new Pair(x=1.2,
                 y=3.4);
writeln("p(0)=", p(0));
writeln("p(1)=", p(1));
p(0) = 5.6;   p(1) = 7.8;
for x in p {
  writeln(x);
  x = -x;        }
writeln(p);
```

- Output

```
p(0)=1.2
p(1)=3.4
5.6
7.8
(-5.6, -7.8)
```


Compile-Time Language

- Chapel has rich compile-time capabilities
  - loop unrolling
  - conditional folding
  - user-defined functions that can be evaluated at compile-time
  - …
- Supported via two main language concepts:
  - **`type`** – type variables and expressions
  - **`param`** – compile-time constants
- In order to support static typing and good performance…

  …the compiler must be able to determine the static types of…
    - variables/members
    - function arguments/return types

  …parameter values are required in certain contexts
    - array ranks
    - indexing of heterogeneous tuples

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Compile-time Language Examples

```
param numDims = 2;        // declare a compile-time constant
type elemType = int;      // declare a named type

def sqr(param x) param {  // declare a param function
  return x*x;             // std ops on params create params
}

param nDSq = sqr(numDims);       // use it to create a param

def myInt(param big: bool) type {  // declare a type fun.
  if (big) then return int(64);    // param test =>
          else return int(32);     //   fold conditional
}

var myTuple = (1, "hi", 2.3);      // heterogeneous tuple

for i in 1..3 do                   // illegal: types vary
  writeln(myTuple(i));             //   across iterations

for param i in 1..3 do             // param index =>
  writeln(myTuple(i));             //   unroll loop
```


Generic Functions, Records, Classes

- `type` and `param` are also used for generic programming
  - a copy of the function/class is stamped out for each unique signature
  - generic functions are created by accepting param/type arguments

```
def x2y2(type t, x: t, y: t): t {
  return x**2 + y**2;
}
x2y2(int, 2, 3);
x2y2(real, 1.2, 3.4);
```

Note: *recall that eliding a formal argument's type also results in a function generic in that argument)*

  - generic classes are created by having param/type members

```
class BoundedStack {
  type elemType;
  const bound: int = 10;
  var data: [1..bound] elemType;
}

var myStack = new BoundedStack(string, 100);
```


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Other Base Language Features

- config params -- can be set on the compiler command-line
- function/operator overloading
  - `where` clauses to decide between overloads using type/param exprs
- argument query syntax

```
def x2y2(x: ?t, y: t): t {
  return x**2 + y**2;
}
x2y2(2, 3);
x2y2(1.2, 3.4);
x2y2(1, 2.3);
```

- tuple types, enumerated types, type unions
- file I/O
- nested modules
- standard modules

Questions?

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