

Chapel: Locality



The Locale

Definition

- Abstract unit of target architecture
- Capable of running tasks and storing variables
 - i.e., has processors and memory
- Supports reasoning about locality

Properties

- a locale's tasks have ~uniform access to local vars
- Other locale's vars are accessible, but at a price

Locale Examples

- A multi-core processor
- An SMP node



"Hello World" in Chapel: a Multi-Locale Version

Multi-locale Hello World



Locales and Program Startup

Specify # of locales when running Chapel programs

% a.out --numLocales=8

% a.out -nl 8

• Chapel provides built-in locale variables

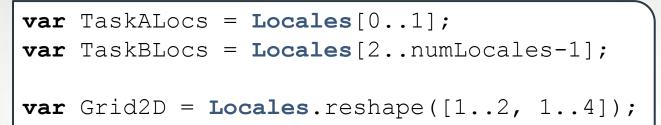
config	g const n	umLocales:	<pre>int;</pre>	
const	LocaleSp	ace: domai	n (1)	= [0numLocales-1];
const	Locales:	[LocaleSp	ace]	locale;

numLocales:	8
LocaleSpace:	
Locales:	L0 L1 L2 L3 L4 L5 L6 L7

• main() begins as a single task on locale #0 (Locales [0])



Create locale views with standard array operations:





Locale Methods



	proc	<pre>locale.id:</pre>	int	{
--	------	-----------------------	-----	---

Returns locale's index in LocaleSpace

proc locale.name: string { ... }

Returns name of locale, if available (like uname -a)

```
proc locale.numCores: int { ... }
```

Returns number of processor cores available to locale

```
proc locale.physicalMemory(...) { ... }
```

Returns physical memory available to user programs on locale

Example

const totalPhysicalMemory =

+ **reduce** Locales.physicalMemory();

The On Statement



• Syntax

on-stmt:
 on expr { stmt }

- Semantics
 - Executes stmt on the locale that stores expr

• Example

writeln("start on locale 0");
on Locales(1) do
 writeln("now on locale 1");
writeln("on locale 0 again");



Locality and Parallelism are Orthogonal

On-clauses do not introduce any parallelism

```
writeln("start on locale 0");
on Locales(1) do
    writeln("now on locale 1");
writeln("on locale 0 again");
```

But can be combined with constructs that do:

```
writeln("start on locale 0");
begin on Locales(1) do
  writeln("now on locale 1");
on Locales(2) do begin
  writeln("now on locale 2");
writeln("on locale 0 again");
```

• (the final three statements could appear in any order)



SPMD Programming in Chapel Revisited

 A language may support both global- and local-view programming — in particular, Chapel does

```
proc main() {
  coforall loc in Locales do
    on loc do
        MySPMDProgram(loc.id, Locales.numElements);
}
proc MySPMDProgram(me, p) {
   ...
}
```

Querying a Variable's Locale



• Syntax

```
locale-query-expr:
    expr . locale
```

- Semantics
 - Returns the locale on which expr is stored

• Example

```
var i: int;
on Locales(1) {
  var j: int;
  writeln(i.locale.id, j.locale.id); // outputs 01
}
```







• Built-in locale value

const here: locale;

- Semantics
 - Refers to the locale on which the task is executing

• Example

writeln(here.id); // outputs 0
on Locales(1) do
writeln(here.id); // outputs 1

Serial Example with Implicit Communication



var x, y: real;	// x and y allocated on locale 0
	// migrate task to locale 1 // z allocated on locale 1
z = x + y;	// remote reads of x and y
	// migrate back to locale 0 // remote write to z // migrate back to locale 1
	<pre>// data-driven migration to locale 0 // remote write to z // migrate back to locale 1 // migrate back to locale 0</pre>

LO	X	L1	z	
	У			



Status: Locales

- Everything should be functioning perfectly
- The compiler is currently conservative about assuming variables may be remote
 - Impact: scalar performance overhead
- The compiler is currently lacking several important communication optimizations
 - Impact: performance impact for programs that would benefit by aggregated communication

Future Directions



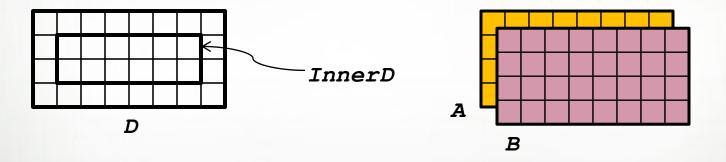
• Hierarchical Locales (joint work with UIUC)

- Support ability to expose hierarchy, heterogeneity within locales
- Particularly important in next-generation nodes
 - CPU+GPU hybrids
 - tiled processors
 - manycore processors



Review: Data Parallelism

- Domains are first-class index sets
 - Specify the size and shape of arrays
 - Support iteration, array operations, etc.



May optionally be distributed over multiple locales

Can be stored in local memories in arbitrary ways





Domain maps are "recipes" that instruct the compiler how to map the global view of a computation...

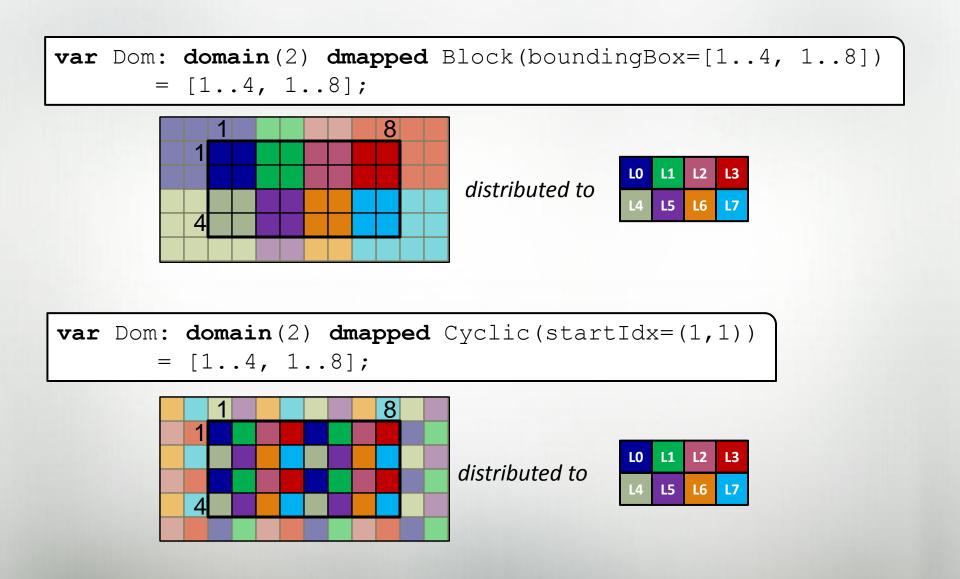


...to a locale's memory and processors:





Sample Distributions: Block and Cyclic





Chapel's Domain Map Strategy

- 1. Chapel provides a library of standard domain maps
 - to support common array implementations effortlessly
- 2. Advanced users can write their own domain maps in Chapel
 - to cope with shortcomings in our standard library
- 3. Chapel's standard layouts and distributions will be written using the same user-defined domain map framework
 - to avoid a performance cliff between "built-in" and user-defined domain maps
- 4. Domain maps should only affect implementation and performance, not semantics
 - to support switching between domain maps effortlessly



 I will be talking at length about domain maps on Tuesday morning, so thought I'd save some time here by asking you to attend that talk



- Full-featured Block, Cyclic, and Replicated distributions
- Single-locale COO and CSR Sparse layouts supported
- Serial quadratic probing Associative layout supported
- Block-Cyclic, Dimensional, and Associative distributions underway
- Adding documentation for defining domain maps
- Memory currently leaked for distributed arrays

Future Directions



Advanced uses of domain maps:

- GPU programming
- Dynamic load balancing
- Resilient computation
- *in situ* interoperability
- Out-of-core computations

Questions?



