

Portability and Platform-Specific Work

Chapel Team, Cray Inc.
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Outline



- Chapel on Intel® Xeon Phi™ "Knights Landing" (KNL)
 - Chapel Support for KNL High-Bandwidth Memory (HBM)
- Unification of Cray (CCE) Configuration
- Python 2/3 Compatibility
- Other Portability Work

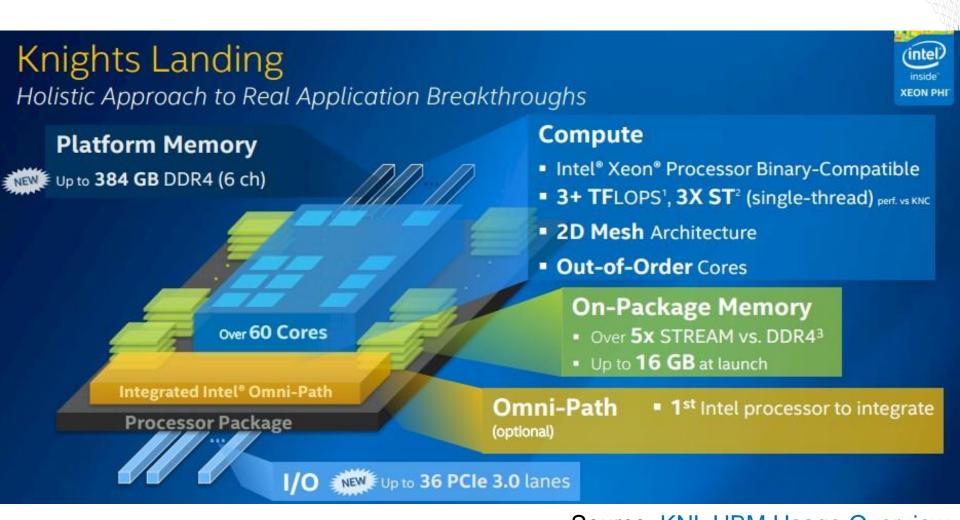




Chapel on Intel® Xeon Phi™ "Knights Landing" (KNL)











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Previous releases included Knights Corner (KNC) support

Targeting KNC with Chapel has several limitations

- requires a special build of the module
- only supports Intel as the target compiler
- all communication has to be bounced through the host processor
 - cannot use "native" communication layers (e.g., ugni/gasnet-aries)
- limited amount of memory

KNC was a good starting point for becoming "KNL-ready"

- allowed us to explore running on a many-core architecture
 - 60+ cores (4 threads per core)
 - each core is relatively slow compared to a traditional Xeon
 - wide (512 bit) vector units





Targeting KNL addresses several KNC-specific limitations:

- binary-compatible with Xeon Processors
- support from Intel, GNU, and Cray target compilers
- can run as a self-hosted processor
- can support native communication layers
- substantially higher memory capacity

That said, KNL remains a unique many-core architecture

- 60+ cores (with 4 threads per core)
- support for AVX-512 (wide vector units)
- high-performance on-package memory (MCDRAM)
 - a.k.a. high-bandwidth memory (HBM)
- several "clustering" modes available
 - all-to-all, quadrant, sub-numa
- for more information, see: https://software.intel.com/en-us/articles/what-disclosures-has-intel-made-about-knights-landing





- Several KNL-based systems will be coming online soon
 - e.g., Cori and Trinity: two large systems being manufactured by Cray
 - each system will have over 9,000 self-hosted KNL nodes
 - will be running a new version of Cray's OS and system management stack
- Main KNL goal for 1.13 was to simply get up and running
 - intentionally modest, due to limited access to hardware/systems
- Secondary goal: start looking at performance
 - improve multi-threaded performance using Qthreads tasking
 - add a locale model that supports targeting HBM
 - ongoing work, reported on in following section
 - start exploring vectorization performance



Chapel on KNL: This Effort



Default to qthreads tasking on KNC

- previously defaulted to fifo
 - got poor performance with qthreads in early days of KNC work
- subsequent work with qthreads has resulted in better performance
 - using different scheduler, affinity options, etc.
- had to disable guard pages for a few tests
 - concurrent calls to mprotect() appear to be expensive on Xeon Phi
- stream performance is now on par with reference for KNC
 - early results on KNL also look promising

Numerous bug fixes for CLE 6.X and KNL

- primarily related to differences with new Cray OS
 - updates to ugni comm layer memory registration
 - fixes to hwloc for /proc/mounts overflow bug
 - updates to module build process
 - several miscellaneous bug fixes



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Chapel on KNL: Impact



- Achieving good multi-threaded performance on KNC
 - and on KNL for the limited testing we've had the chance to do
- Correctness test suite passing for all KNL configurations
 - matrix of nightly testing over comm/tasking layers vs. target compilers:

Configuration Matrix	cray	intel	gnu
none			
gasnet-aries	•	•	•
gasnet-mpi	•	•	•
ugni-muxed	•	•	•
ugni-qthreads	•	•	•

(blue means "clean run" in our test system)



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Chapel on KNL: Next Steps



- Retire Chapel's support for KNC
 - was a useful stepping-stone towards KNL
 - we believe it is no longer important to Chapel's end-users
 - please let us know if we're mistaken!
- Add KNL-specific locale model to support targeting HBM
- Verify that Qthreads provides good performance on KNL
 - and see if guard page overhead can be lowered
- Work on vectorization improvements/optimizations
 - and explore any potential KNL-specific optimizations





Chapel Support for KNL High-Bandwidth Memory (HBM) (an ongoing effort)



KNL HBM: Background



KNL supports two memory systems:

- off-package DDR memory
 - 90+ GB/s and up to 384 GB
- high-performance on-package memory (a.k.a. MCDRAM or HBM)
 - 400+ GB/s and up to 16 GB

MCDRAM can operate in different modes:

- partly or entirely as a cache to system memory
- allocating all data in the MCDRAM
- as a memory that is independent from system memory
 - programs use it by allocating explicitly from the MCDRAM

KNL will support dividing into 4 distinct NUMA domains

Sources: Intel KNL Disclosures KNL HBM Usage Overview



KNL HBM: Representing Using Locales



Represent MCDRAM as part of a Chapel locale model

- our first example of a *memory-only* sub-locale
- such locales can be used like any other:
 - you can target them using on-statements
 - to store data
 - to execute tasks
 - you can store them in an array and target them with distributions

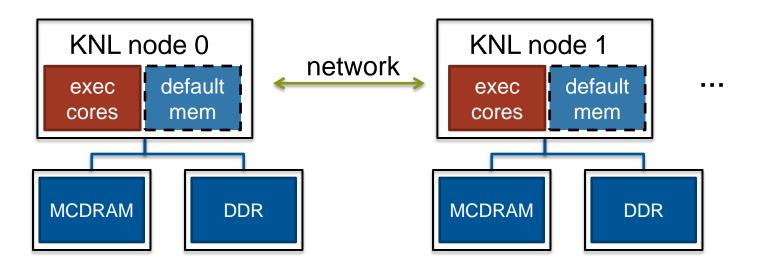
The following slides propose our default policy for KNL

- since it's part of a locale model, none of this is baked into the compiler
 - i.e., if you want a different KNL policy, you can create one
 - (or you can advocate for changes to ours)
- This is all work-in-progress



KNL HBM: Proposed Policy

- Tasks on memory locales execute on the parent
 - but with a constraint: any allocations are from the designated memory
- Tasks on top-level KNL locales use a default child memory



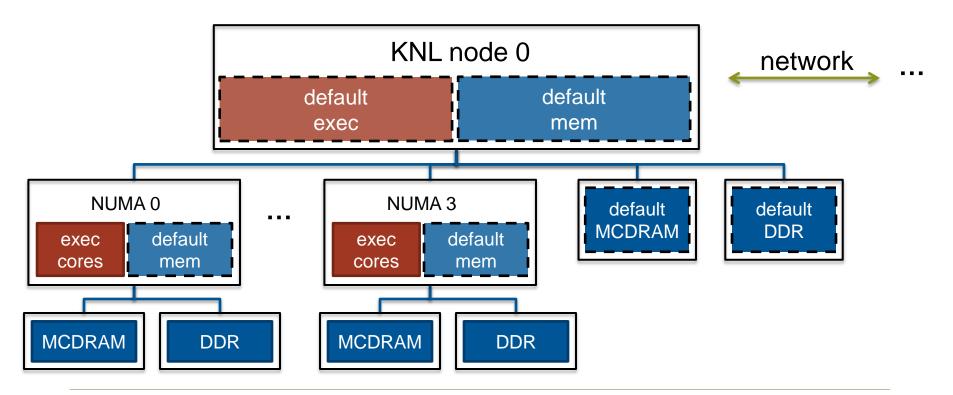


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KNL HBM: Representing NUMA Locales



- Challenge: How to represent distinct NUMA domains?
 - give memory sub-locales to each top-level KNL locale, as before
 - also give each KNL locale multiple NUMA sub-locales
 - give each NUMA sub-locale its own memory sub-locales



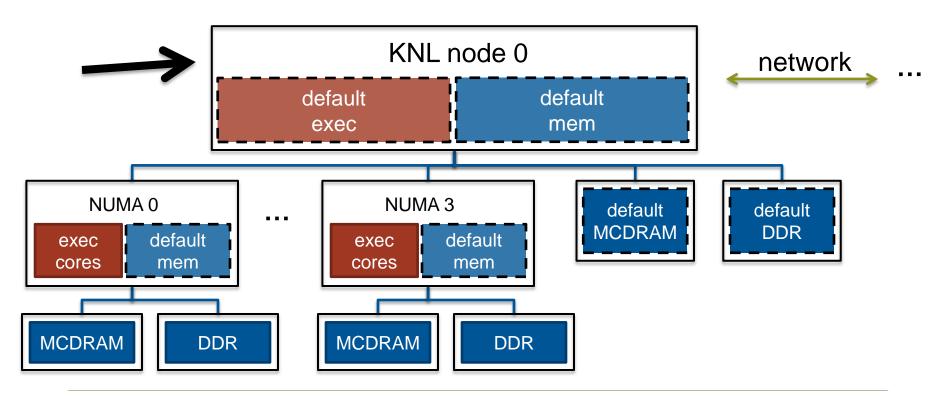


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• Executing on a top-level KNL locale will:

- choose among the NUMA sub-locales for where to run the task (as in the NUMA locale model)
- use a default memory type for its allocations

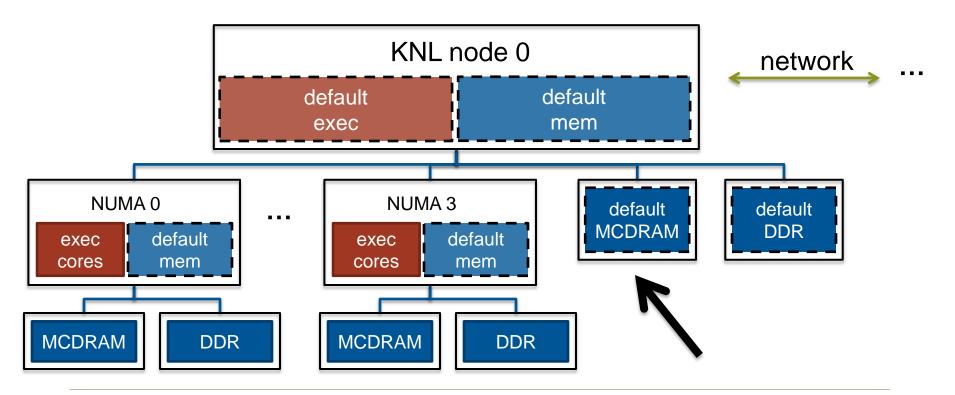




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- Executing on a KNL locale's memory sub-locale will:
 - choose among the NUMA sub-locales for where to run the task (as in the previous case)
 - use the specified memory type for its allocations

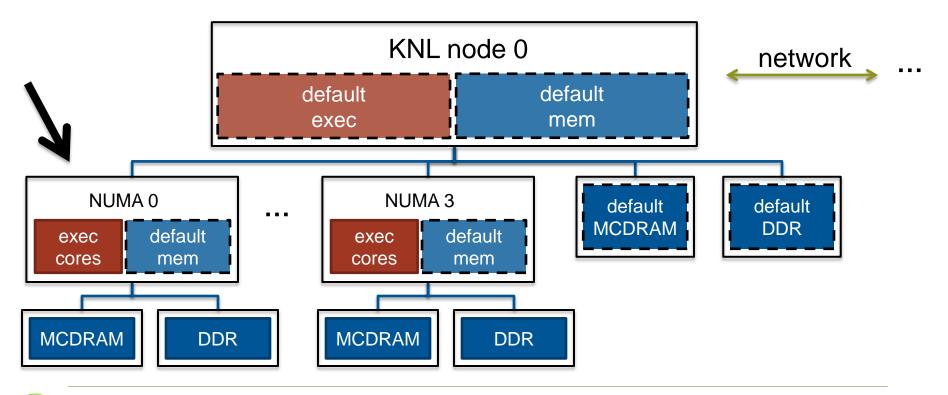




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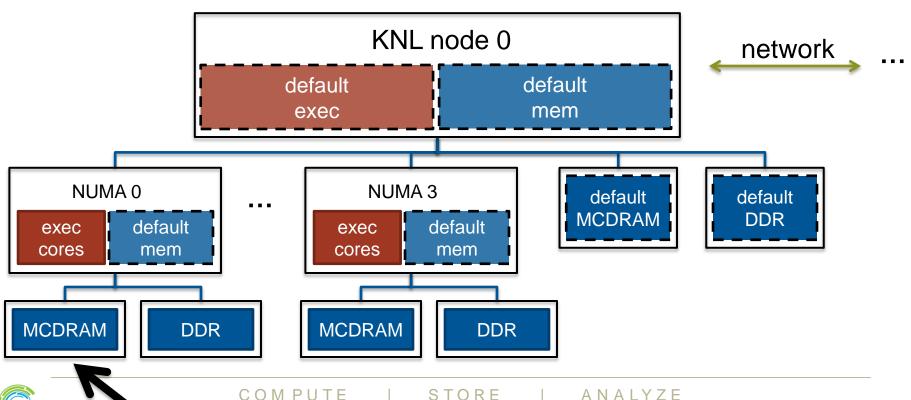
- Executing on a KNL locale's NUMA sub-locale will:
 - run the task on one of the cores within the NUMA domain
 - and limit child tasks to those same cores, unless an on-clause is used
 - use a default memory type for its allocations





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- Executing on a NUMA domain's memory sub-locale will:
 - run the task on one of the cores within the NUMA domain
 - and limit child tasks to those same cores, unless an on-clause is used
 - use the specified memory type for allocations





KNL HBM: Mapping forall loops to KNL



- Memory locales are locales like any other
 - Yet, we don't want forall loops to run tasks on memory sub-locales
- Default foralls currently use simple locale model methods:

proc getChildCount(): int

proc getChild(idx:int) : locale

We'll need to update this interface to avoid the memories



KNL HBM: Allocation Experiments



- Developed prototype KNL locale model
 - targeted hbw_malloc functions provided by memkind library
- Demonstrated allocating explicitly as follows:

```
// get HBM with a temporary interface: LocaleModel.hbm
var hbm = (here:LocaleModel).hbm;

on hbm {
    // demonstrated this array is allocated with hbw_malloc
    var A: [1..1000] int;
    A = 1;
}
```

In-progress: Support for Block distribution over HBM



KNL HBM: Impact, Status, and Next Steps



Impact:

- Well on our way to MCDRAM support for KNL
 - Chapel programs will be able to explicitly program the memory
 - Arrays will be able to be distributed over many MCDRAMs in a system

Status:

early prototype demonstrated (not included in v1.13)

Next Steps:

- evolve Locale interface to support HBM
- support capability-oriented interfaces for users for portability, e.g.:
 - "get locale ancestor running the operating system"
 - "get a fast-memory sub-locale"
- complete implementation of KNL Locale Model
- investigate NUMA support on KNL
- investigate performance





Unification of Cray (CCE) Configuration



CCE Target Compiler: Background



- On Crays we support Cray, GNU, and Intel target compilers
- Cray (CCE) hasn't supported all runtime configurations
 - used fifo tasking
 - used gasnet comm
 - used dlmalloc memory
 - lacked regexp support

couldn't build qthreads or muxed (inline ASM)

couldn't build ugni (missing atomics)

couldn't build tcmalloc (C++, compiler flags, etc)

couldn't build re2 (Makefiles used GNU flags)

- As a result, using the CCE back-end limited performance
 - and did not support regular expressions
- CCE 8.4.0 allows us to improve on this by adding...
 - ... inline ASM support
 - ...additional GNU compatability
 - ...missing atomics



CCE Target Compiler: This Effort



Build qthreads+hwloc with CCE

- required working around broken ffs() for hwloc
- required CCE 8.4.0 for Qthreads (inline ASM)

Build muxed with CCE

required CCE 8.4.0 (inline ASM)

Build jemalloc with CCE

- required building with –hipa2 to work around inlining bug
- had to avoid throwing GNU-specific flags

Build ugni with CCE

- required building with —hipa0 to work around inlining bug
- required jemalloc, since getting tcmalloc to build proved impractical

Build re2 with CCE

edited Makefiles to avoid GNU-specific flags

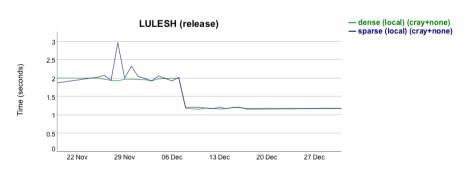


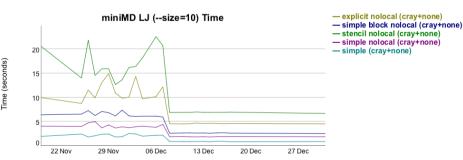
CCE Target Compiler: Impact

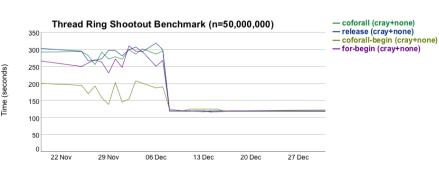


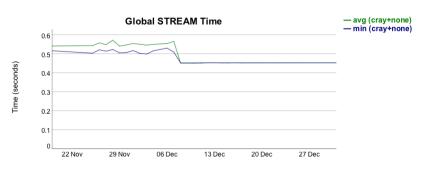
CCE back-end now supports the full runtime

- significantly better performance due to qthreads+hwloc, jemalloc, ugni
- regular expression support











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Python 2/3 Compatibility



Python 2/3 Compatibility: Background



Background:

- Building Chapel required Python 2.x, but Python 3 is used widely
- Had become our most frequent barrier to successful end-user installs

This Effort:

Modified Chapel configuration scripts to support Python 3.x as well

Impact:

Users can now build Chapel with Python 2 or 3

Status:

Nightly testing checks compatibility with Python 2.6, 2.7, and 3.5

Next Steps:

- Make the testing infrastructure Python 3-compatible
 - e.g., start_test





Other Portability Work



Other Portability Work



- Switched PGI C++ compiler from pgCC to pgc++
- Worked around code base incompatibilities with gcc 5.1
 - seemingly a gcc 5.1 compiler bug, as unlikely as that seems...
- Fixed portability issues to newer Cygwin installs
- Removed stale Xcode support



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