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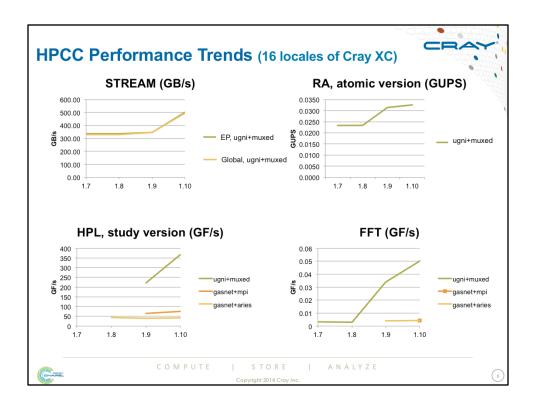
# **Executive Summary**

- Generally speaking, performance has improved with 1.10
- Previous slide decks have shown performance changes:
  - ....in the shootout benchmarks
  - ...due to threading mechanism and policy changes
  - ...due to LICM improvements
  - ...due to the addition of CHPL TARGET ARCH/--specialize
  - ...due to generated code improvements
- These slides contain additional v1.10 performance results
  - not tied to any specific effort, just comparisons across releases



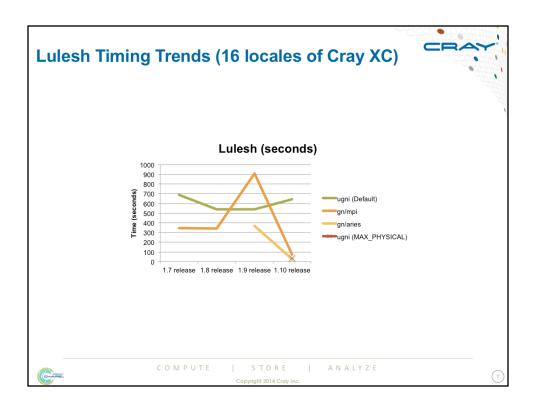






Higher is better in all these graphs.

Note that GASNet over MPI completes for the first time as of 1.10 for FFT.



This is our "study" version of LULESH, running with the sedov15oct input deck gn = GASNet

Lower is better here.

Note that LULESH strongly prefers using a number of threads equal to the number of physical cores for ugni/muxed (shown with the single 'X' point gathered with 1.10, which competes with the gasnet/aries results).

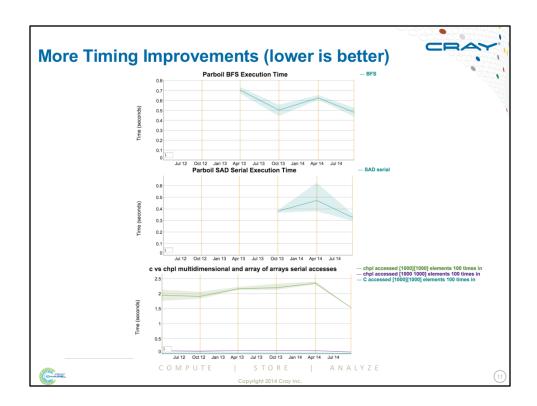


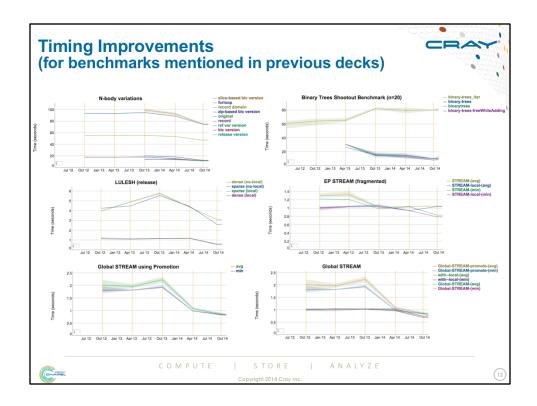
# Release-over-Release Improvements

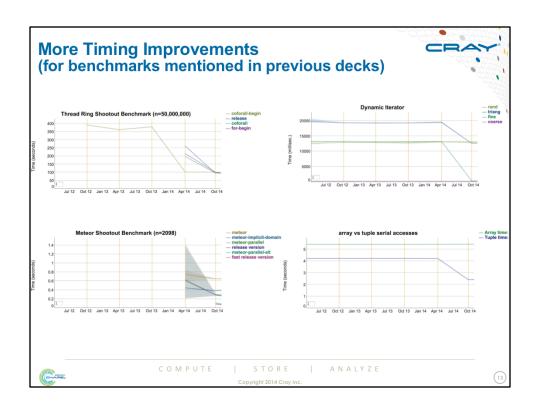
- The following slides summarize release-over-release improvements across several releases
  - today's benchmark codes
  - today's OS, gcc, etc.
  - retroactively run on past releases
- These were collected on chap04
  - 64-bit Linux
  - 2 quad-core Intel Xeon processors with hyper-threading
    - 16 logical cores
  - 48 GB RAM
  - Can be viewed interactively at: <a href="http://chapel.sourceforge.net/perf/">http://chapel.sourceforge.net/perf/</a>

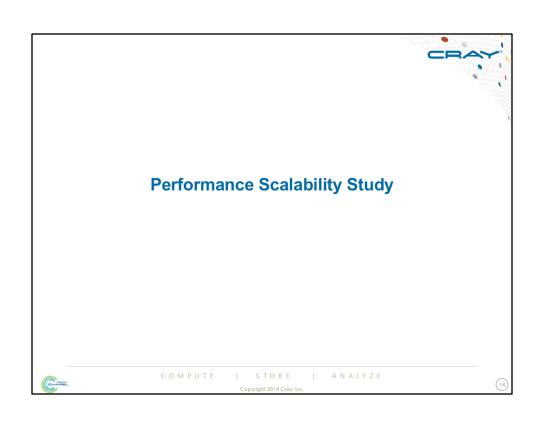










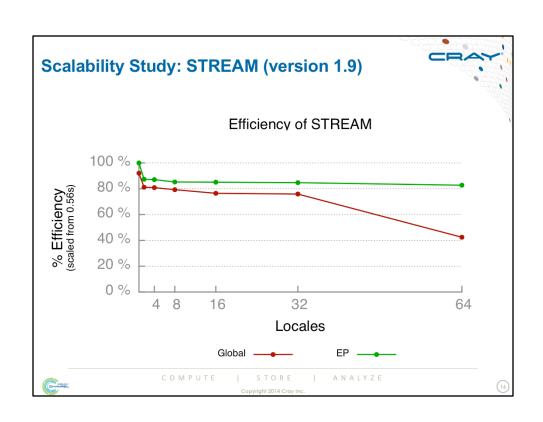


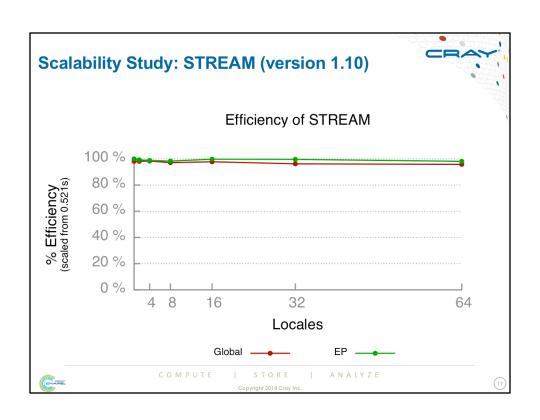
## **Scalability Study: Background**

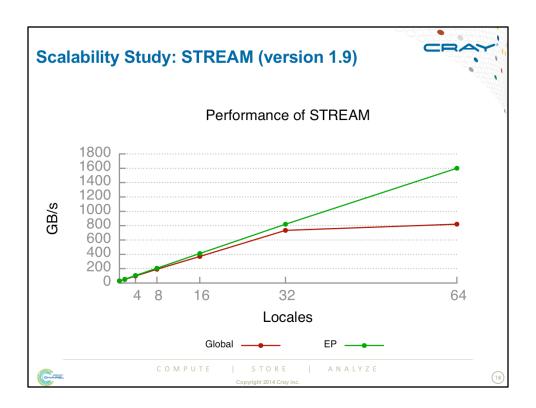
- For this release, we performed a scalability study for some basic benchmarks
  - HPCC Stream: EP and Global
  - HPCC RA: atomic, on-based, and remote memory operations (rmo)
    - these test network atomics, active messages, and puts/gets, respectively
  - Reduction of an array
- All experiments shown here were performed on a Cray XC

  - ugni+muxed runtime
- Each of the following pairs of slides shows 1.9 vs. 1.10
  - Note that some graphs are efficiency, some are performance
  - Also note that the scales of the graph change between versions
    - version 1.10 performs strictly better for these benchmarks on 1 locale

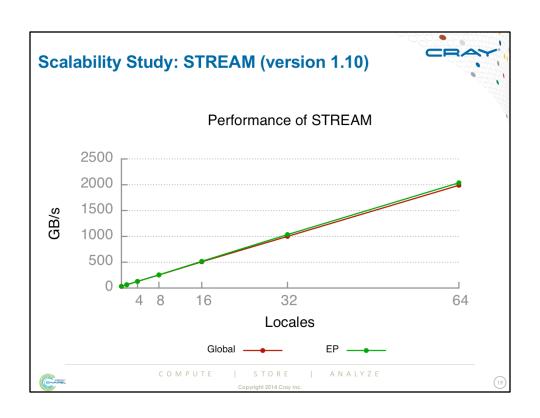


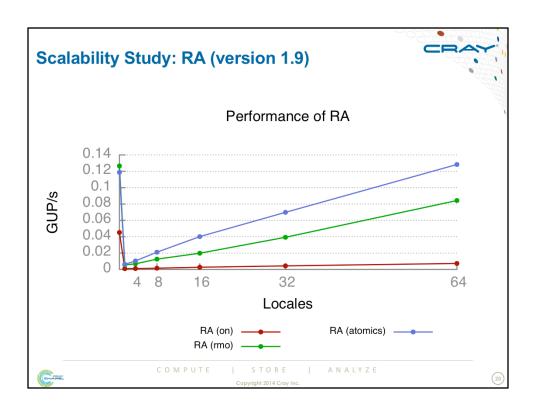




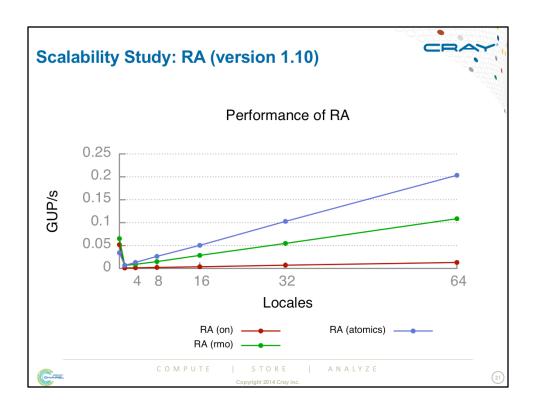


This pair of graphs is just a different way of looking at the data in the previous ones

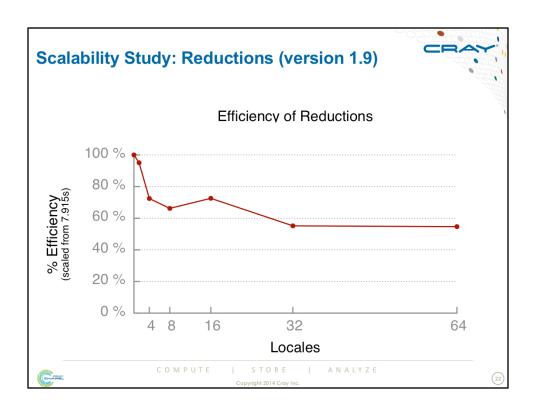




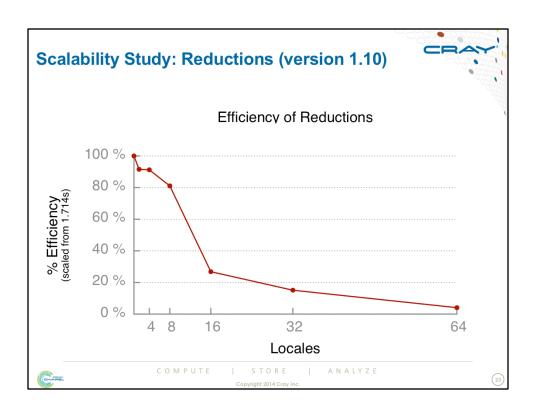
For RA, the drop in performance when moving from 1 to 2 locales is expected because you move from a world in which all updates are local to one in which half are remote. Note that our version does not do bucketing, so tends to do worse on 2 locales than the reference version.



Note that as of 1.10, the parallel performance crosses over the 1-locale performance at a much lower locale count than in 1.9



This test creates an array requiring ¼ of the available memory across all locales and measures the time required to compute a sum-reduce over it.



It's not clear to use yet what is making reductions scale worse in version 1.10. Interestingly, setting CHPL\_RT\_NUM\_THREADS\_PER\_LOCALE='MAX\_PHYSICAL' causes scalability to be much more similar to the 1.9 version.

## **Scalability Study: Next Steps**



- Scale studies to higher locale counts
  - This study limited by availability/size of systems we could access
  - Have since gotten accounts on larger/more available external systems
- Improve NUMA mapping for stream
- Investigate scalability limiters for on-based RA for ugni
- Tune reductions
  - Investigate scalability challenges
  - Re-implement for performance
    - Ideally by moving implementation from compiler to modules
    - Approach would benefit from (ongoing) data parallel task intents work
- Compare across more communication layers, benchmarks





#### Slurm/uGNI conflict



### **Background:** Slurm limits CHPL\_COMM=ugni throughput

- comm=ugni uses multiple uGNI communication domains (NIC cores)
  - default: # comm domains = # pthreads
  - minimizes thread contention, maximizes NIC throughput
  - program memory must be registered separately in each comm domain
- but slurm limits per-process Aries resources, to allow for node sharing
  - effectively limits total memory registration per process to about 240 GiB
- for example, cannot have 32 GiB heap and 8 comm domains

### **This Effort:** Under slurm, throttle to fit within its limits

- reduce default heap size to 16 GiB (v. ALPS: nearly all free memory)
- given heap size, limit # comm domains so total registration < 240 GiB
  - 16 GiB: 15 comm domains; 32 GiB: 7 comm domains; etc.

Impact: May not reach full Aries performance with slurm

**Next Steps:** Requires network-exclusive mode from Slurm





## **Performance Priorities and Next Steps**



- understand/improve cases like RA using ons
- improve coverage, reporting, and graphing of automated testing
- investigate support for ugni with qthreads and/or gasnet with muxed

### NUMA-aware performance

- more focus on NUMA locale model
  - particularly execution-time address representation
- improve array initialization (parallel, appropriate first-touch)
  - currently gated by constructor/default init/noinit capabilities
- consider using NUMA by default (?)

### SIMD-ization of generated code

#### Continue scalability studies

- Reduce unnecessary communication
- Improve implementation of reductions



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