









































































What is Chapel?



Chapel: A productive parallel programming language

- portable & scalable
- open-source & collaborative

Goals:

- Support general parallel programming
 - "any parallel algorithm on any parallel hardware"
- Make parallel programming at scale far more productive





Chapel and Productivity



Chapel aims to be as...

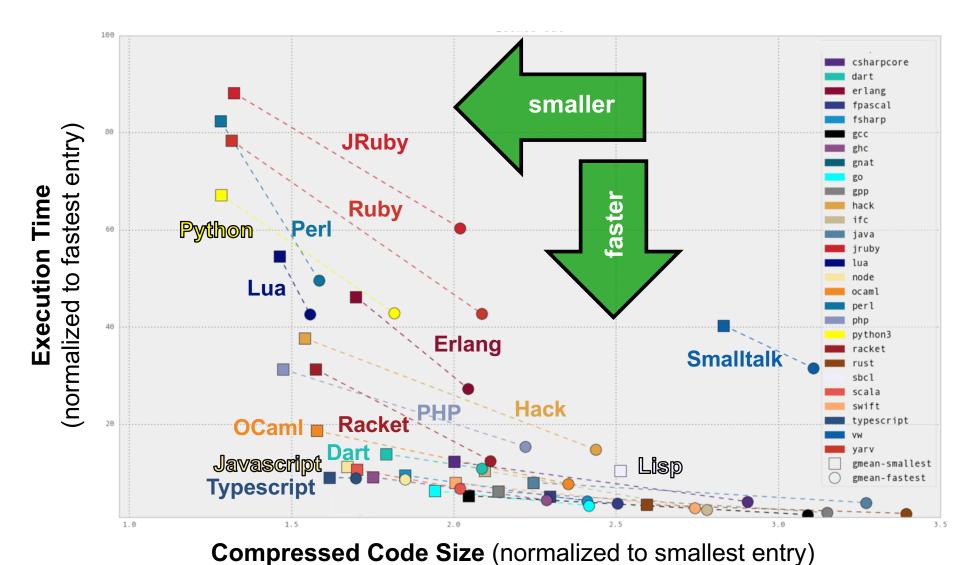
- ...programmable as Python
- ...fast as Fortran
- ...scalable as MPI, SHMEM, or UPC

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- ...portable as C
- ...flexible as C++
- ...fun as [your favorite programming language]



(Oct 2017 standings)

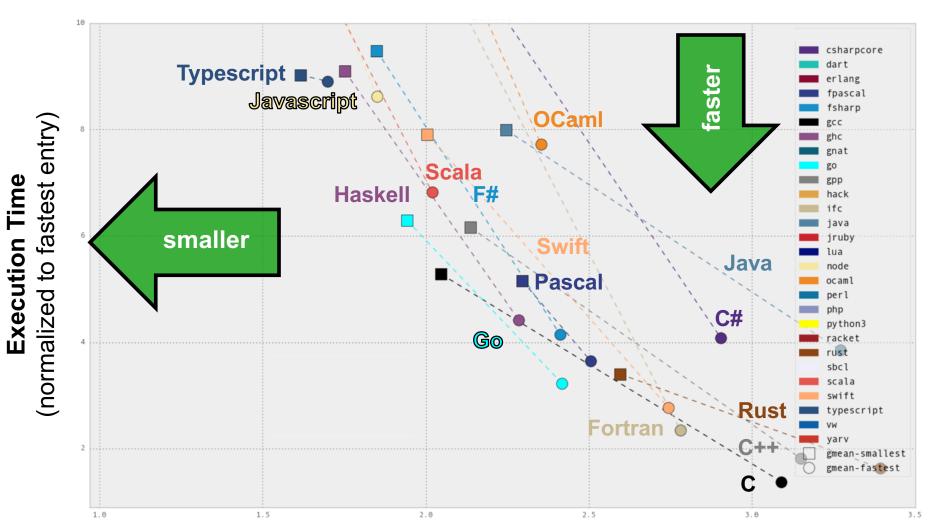


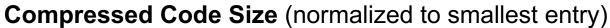


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(Oct 2017 standings, zoomed in)



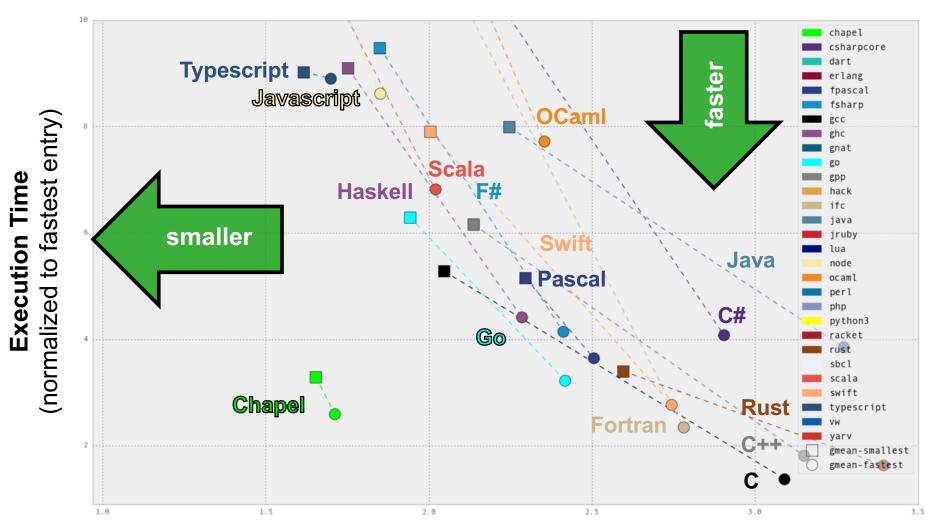




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(Oct 2017 standings, zoomed in)

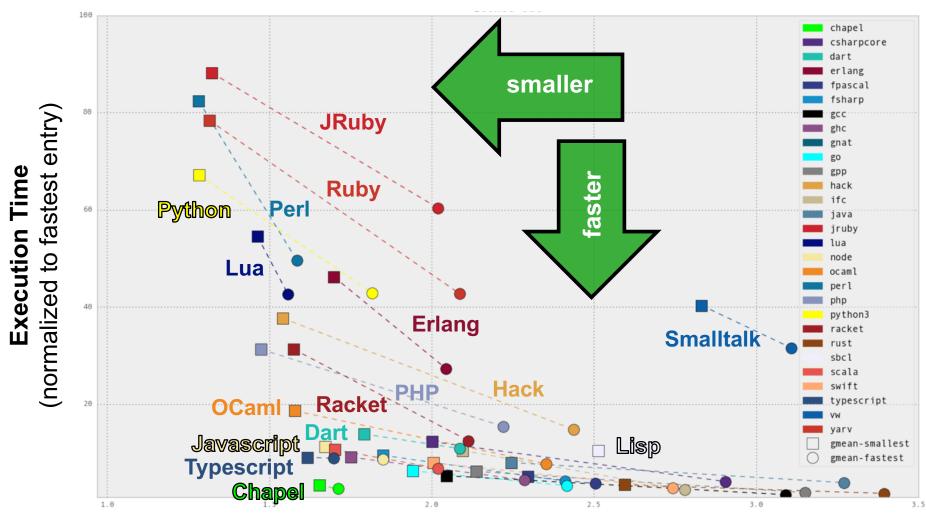






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(Oct 2017 standings)







CLBG: Qualitative Code Comparisons



Can also browse program source code (but this requires actual thought!):

```
proc main() {
 printColorEquations();
  const group1 = [i in 1..popSize1] new Chameneos(i, ((i-1)%3):Color);
 const group2 = [i in 1..popSize2] new Chameneos(i, colors10[i]);
  cobegin {
    holdMeetings(group1, n);
    holdMeetings(group2, n);
  print(group1);
 print(group2);
 for c in group1 do delete c;
 for c in group2 do delete c;
// Print the results of getNewColor() for all color pairs.
proc printColorEquations() {
 for c1 in Color do
    for c2 in Color do
      writeln(c1, " + ", c2, " -> ", getNewColor(c1, c2));
 writeln();
// Hold meetings among the population by creating a shared meeting
// place, and then creating per-chameneos tasks to have meetings.
proc holdMeetings(population, numMeetings) {
 const place = new MeetingPlace(numMeetings);
 coforall c in population do
                                        // create a task per chameneos
   c.haveMeetings(place, population);
  delete place;
```

```
void get affinity(int* is smp, cpu set t* affinity1, cpu set t* affinity2)
   cpu set t
                                active cpus;
   FILE*
   char
                                buf [2048];
   char const*
   int
                                cpu idx;
   int
                                physical id;
   int
                                core id;
   int
                                cpu cores;
   int
                                apic_id;
   size t
                                cpu_count;
   size t
   char const*
                                processor_str
                                                     = "processor";
   size t
                                processor_str_len
                                                     = strlen(processor_str);
   char const*
                                physical id str
                                                     = "physical id";
                                physical id str len = strlen(physical id str);
   size t
   char const*
                                core id str
                                                     = "core id";
   size t
                                core id str len
                                                     = strlen(core id str);
   char const*
                                cpu cores str
                                                     = "cpu cores";
   size t
                                cpu_cores_str_len
                                                    = strlen(cpu_cores_str);
   CPU ZERO(&active cpus);
   sched getaffinity(0, sizeof(active cpus), &active cpus);
   cpu count = 0;
   for (i = 0; i != CPU SETSIZE; i += 1)
       if (CPU ISSET(i, &active cpus))
           cpu_count += 1;
   if (cpu count == 1)
       is smp[0] = 0;
       return;
   is\_smp[0] = 1;
   CPU ZERO(affinity1);
```

excerpt from 1210 gz Chapel entry

excerpt from 2863 gz C gcc entry



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CLBG: Qualitative Code Comparisons



Can also browse program source code (but this requires actual thought!):

```
proc main() {
                                                                                      void get affinity(int* is smp, cpu set t* affinity1, cpu set t* affinity2)
 printColorEquations();
const group1 = [i in 1 pepSize1] new Chameneos(i,
const group2 = [i in 1 popSize2] new Chameneos(i,
                                                                                                                active_cpus;
                                                   cobegin {
                                                                                                                buf [2048];
                                                       holdMeetings(group1, n);
 cobegin {
                                                                                                                cpu idx;
   holdMeetings(group1, n);
                                                                                                                physical id;
   holdMeetings(group2, n);
                                                       holdMeetings(group2, n);
                                                                                                                core id;
                                                                                                                cpu cores;
print(group1);
                                                                                                                apic_id;
                                                                                                                cpu_count;
                                                                                                                processor str
                                                                                                                                = "processor";
 for c in group2 do delete c;
                                                                                                                processor str len
                                                                                                                                = strlen(processor_str);
                                                                                         char const*
                                                                                                                physical id str
                                                                                                                                = "physical id";
                                                                                                                physical id str len = strlen(physical id str);
                                                                                         size t
                                                                                         char const*
                                                                                                                core id str
                                                                                                                                = "core id";
                                                                                                                                      n(core id str);
// Print the results of getNewColor() for all color p
                                              proc holdMeetings(population, numMeetings) {
                                                                                                                                     cores";
                                                                                                                                     n(cpu_cores_str);
proc printColorEquations() {
                                                  const place = new MeetingPlace(numMeetings);
 for cl in Color do
   for c2 in Color do
    writeln(c1, " + ", c2,
                             getNewColor(c1,
 writeln();
                                                  coforall c in population do
                                                                                                                       // creat
                                                     c.haveMeetings(place, population);
// Hold meetings among the population by creating a s
// place, and then creating per-chameneos tasks to ha
                                                  delete place;
proc holdMeetings(population, numMeetings)
 const place = new MeetingPlace(numMeetings);
 coforall c in population do
                                // create a t
   c.haveMeetings(place, population);
                                                                                         is_smp[0] = 1;
 delete place;
                                                                                         CPU ZERO(affinity1);
    excerpt from 1210 gz Chapel entry
                                                                                              excerpt from 2863 gz C gcc entry
```



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CLBG: Qualitative Code Comparisons



Can also browse program source code (but this requires actual thought!):

```
proc main() {
char const*
                             core id str
                                                   = "core id
size t
                             core id str len
                                                  = strlen(co:
                                                  = "cpu core
char const*
                             cpu cores str
                                                  = strlen(cp
size t
                             cpu cores str len
CPU ZERO(&active cpus);
sched getaffinity(0, sizeof(active cpus), &active cpus);
cpu count = 0;
for (i = 0; i != CPU SETSIZE; i += 1)
    if (CPU ISSET(i, &active cpus))
        cpu count += 1;
   (cpu count == 1)
    is smp[0] = 0;
    return;
```

```
void get affinity(int* is smp, cpu set t* affinity1, cpu set t* affinity2)
   cpu_set_t
                                active cpus;
   FILE*
   char
                                buf [2048];
   char const*
   int
                                cpu idx;
   int
                                physical id;
   int
                                core id;
   int
                                cpu cores;
   int
                                apic id;
   size t
                                cpu_count;
   size t
   char const*
                                processor_str
                                                     = "processor";
   size t
                                processor str len
                                                     = strlen(processor_str);
   char const*
                                physical id str
                                                     = "physical id";
                                physical id str len = strlen(physical id str);
   size t
                                                     = strlen(core id str);
   size t
                                core id str len
   char const*
                                cpu cores str
                                                     = "cpu cores";
   size t
                                cpu_cores_str_len
                                                     = strlen(cpu_cores_str);
   CPU_ZERO(&active_cpus);
   sched getaffinity(0, sizeof(active cpus), &active cpus);
   cpu count = 0;
   for (i = 0; i != CPU SETSIZE; i += 1)
       if (CPU ISSET(i, &active cpus))
            cpu_count += 1;
   if (cpu count == 1)
       is smp[0] = 0;
       return;
   is\_smp[0] = 1;
   CPU ZERO(affinity1);
```

excerpt from 1210 gz Chapel entry

excerpt from 2863 gz C gcc entry



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The Chapel Team at Cray (May 2018)







Chapel Community Partners

























Sandia National Laboratories



Yale

(and several others...)

https://chapel-lang.org/collaborations.html



Outline

- √ What is Chapel?
- > Chapel Overview
- Chapel: Then vs. Now

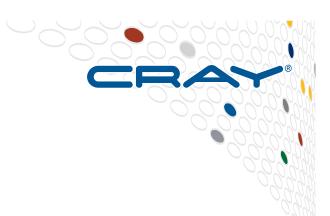
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- Chapel User Profiles
- What's Next?

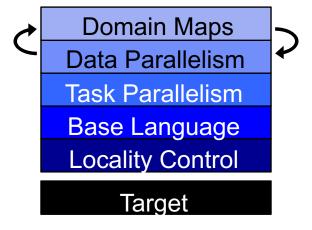


Chapel language feature areas

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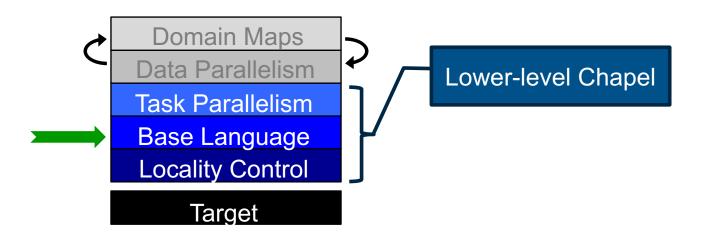
Chapel language concepts





Base Language







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```
iter fib(n) {
  var current = 0,
      next = 1;

  for i in 1..n {
      yield current;
      current += next;
      current <=> next;
    }
}
```

```
config const n = 10;
for f in fib(n) do
  writeln(f);
```

```
0
1
1
2
3
5
8
```





```
Iterators
                           config const n = 10;
iter fib(n)
  var current = 0,
                           for f in fib(n) do
      next = 1;
                             writeln(f);
  for i in 1..n {
    yield current;
                                 1
2
3
5
    current += next;
    current <=> next;
                                 8
```



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ANALYZE

Copyright 2018 Cray Inc.

```
Static type inference for:
```

- arguments
- return types
- variables

```
iter fib(n) {
  var current = 0,
    next = 1;

for i in 1..n {
  yield current;
  current += next;
  current <=> next;
}
}
```

```
config const n = 10;
for f in fib(n) do
  writeln(f);
```

```
0
1
1
2
3
5
8
```



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```
Explicit types also
                     permitted
                            config const n: int = 10;
iter fib(n : \int): int
  var current: int = 0,
                             r f: int in fib(n) do
      next: int = 1;
                             writeln(f);
  for i in 1..n {
    yield current;
                                  1
1
2
3
5
    current += next;
    current <=> next;
                                  8
```



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```
iter fib(n) {
  var current = 0,
     next = 1;

  for i in 1..n {
     yield current;
     current += next;
     current <=> next;
  }
}
```

```
config const n = 10;
for f in fib(n) do
  writeln(f);
```

```
0
1
1
2
3
5
8
```





```
iter fib(n) {
  var current = 0,
    next = 1;

for i in 1..n {
    yield current;
    current += next;
    current <=> next;
}
```

```
config const n = 10;
for (i,f) in zip(0..#n, fib(n)) do
  writeln("fib #", i, " is ", f);
```

Zippered iteration

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...
```





```
operators
                         config const n = 10;
iter fib(n) {
  var current =
                         for (i,f) in zip(0..#n, fib(n)) do
      next = 1
                           writeln("fib #", i, " is ", f);
  for i in 1...n {
                                fib #0 is 0
    yield current;
                                fib #1 is 1
    current += next;
                                fib #2 is 1
    current <=> next;
                                fib #3 is 2
                                fib #4 is 3
                                fib #5 is 5
                                fib #6 is 8
```

Range types and





```
iter fib(n) {
  var current = 0,
    next = 1;

  for i in 1..n {
    yield current;
    current += next;
    current <=> next;
  }
}
```

```
config const n = 10;
for (i,f) in zip(0..#n, fib(n)) do
  writeln("fib #", i, " is ", f);
```

Tuples

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...
```





```
iter fib(n) {
  var current = 0,
    next = 1;

  for i in 1..n {
    yield current;
    current += next;
    current <=> next;
  }
}
```

```
config const n = 10;

for (i,f) in zip(0..#n, fib(n)) do
    writeln("fib #", i, " is ", f);
```

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...
```



Other Base Language Features



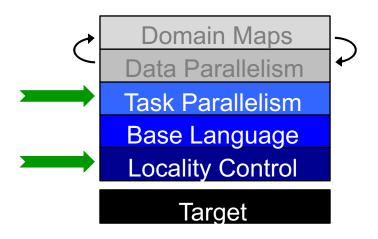
- **Object-oriented features**
- Generic programming / polymorphism
- Procedure overloading / filtering
- Default args, arg intents, keyword-based arg passing
- Argument type queries / pattern-matching
- Compile-time meta-programming
- Modules (namespaces)
- Error-handling
- and more...



Task Parallelism and Locality Control

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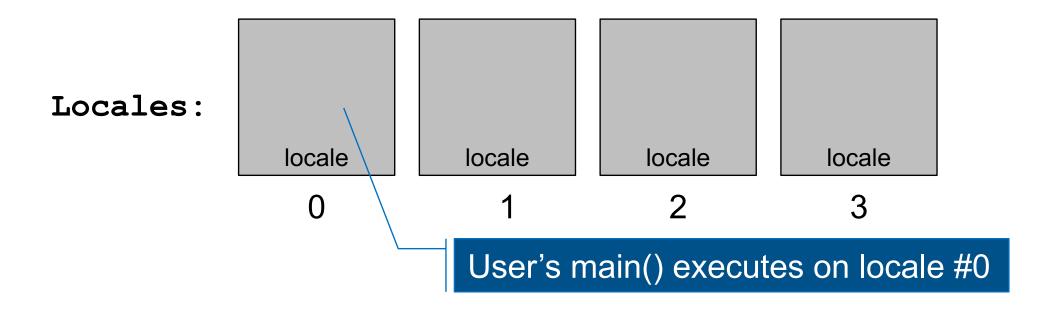






Locales, briefly

- Locales can run tasks and store variables
 - Think "compute node"





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```
taskParallel.chpl
```

```
prompt> chpl taskParallel.chpl
prompt> ./taskParallel
Hello from task 2 of 2 running on n1032
Hello from task 1 of 2 running on n1032
```



Task Parallelism and Locality, by example



```
taskParallel.chpl
                          const numTasks = here.numPUs();
 Abstraction of
                          coforall tid in 1..numTasks do
System Resources
                            writef("Hello from task %n of %n "+
                                    "running on %s\n",
                                    tid, numTasks, here.name);
                 prompt> chpl taskParallel.chpl
                 prompt> ./taskParallel
                 Hello from task 2 of 2 running on n1032
```

Hello from task 1 of 2 running on n1032



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```
prompt> chpl taskParallel.chpl
prompt> ./taskParallel
Hello from task 2 of 2 running on n1032
Hello from task 1 of 2 running on n1032
```







This is a shared memory program

Nothing has referred to remote locales, explicitly or implicitly

taskParallel.chpl

```
const numTasks = here.numPUs();
coforall tid in 1..numTasks do
   writef("Hello from task %n of %n "+
        "running on %s\n",
        tid, numTasks, here.name);
```

```
prompt> chpl taskParallel.chpl
prompt> ./taskParallel
Hello from task 2 of 2 running on n1032
Hello from task 1 of 2 running on n1032
```







```
prompt> chpl taskParallel.chpl
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1033
```



Task Parallelism and Locality, by example



Abstraction of System Resources

```
prompt> chpl taskParallel.chpl
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1033
```



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Task Parallelism and Locality, by example



Control of Locality/Affinity

```
prompt> chpl taskParallel.chpl
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1033
```







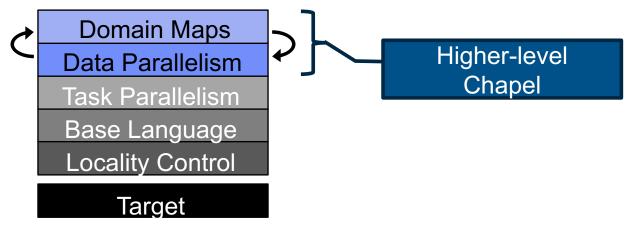


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Data Parallelism in Chapel



Chapel language concepts









```
dataParallel.chpl
```

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl
prompt> ./dataParallel --n=5

1.1 1.3 1.5 1.7 1.9

2.1 2.3 2.5 2.7 2.9

3.1 3.3 3.5 3.7 3.9

4.1 4.3 4.5 4.7 4.9

5.1 5.3 5.5 5.7 5.9
```



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Domains (Index Sets)

dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl
prompt> ./dataParallel --n=5

1.1 1.3 1.5 1.7 1.9

2.1 2.3 2.5 2.7 2.9

3.1 3.3 3.5 3.7 3.9

4.1 4.3 4.5 4.7 4.9

5.1 5.3 5.5 5.7 5.9
```



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Arrays

```
dataParallel.chpl
```

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl
prompt> ./dataParallel --n=5

1.1 1.3 1.5 1.7 1.9

2.1 2.3 2.5 2.7 2.9

3.1 3.3 3.5 3.7 3.9

4.1 4.3 4.5 4.7 4.9

5.1 5.3 5.5 5.7 5.9
```



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Data-Parallel Forall Loops

```
dataParallel.chpl
```

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl
prompt> ./dataParallel --n=5

1.1 1.3 1.5 1.7 1.9

2.1 2.3 2.5 2.7 2.9

3.1 3.3 3.5 3.7 3.9

4.1 4.3 4.5 4.7 4.9

5.1 5.3 5.5 5.7 5.9
```





This is a shared memory program

Nothing has referred to remote locales, explicitly or implicitly

dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl
prompt> ./dataParallel --n=5

1.1 1.3 1.5 1.7 1.9

2.1 2.3 2.5 2.7 2.9

3.1 3.3 3.5 3.7 3.9

4.1 4.3 4.5 4.7 4.9

5.1 5.3 5.5 5.7 5.9
```



Distributed Data Parallelism, by example



Domain Maps (Map Data Parallelism to the System)

```
dataParallel.chpl

use CyclicDist;
config const n = 1000;
var D = {1..n, 1..n}
         dmapped Cyclic(startIdx = (1,1));

var A: [D] real;

forall (i,j) in D do
        A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl
prompt> ./dataParallel --n=5 --numLocales=4
1.1 1.3 1.5 1.7 1.9
2.1 2.3 2.5 2.7 2.9
3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```







dataParallel.chpl

```
use CyclicDist;
config const n = 1000;
var D = {1..n, 1..n}
          dmapped Cyclic(startIdx = (1,1));
var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl
prompt> ./dataParallel --n=5 --numLocales=4

1.1 1.3 1.5 1.7 1.9

2.1 2.3 2.5 2.7 2.9

3.1 3.3 3.5 3.7 3.9

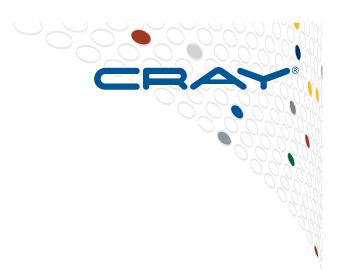
4.1 4.3 4.5 4.7 4.9

5.1 5.3 5.5 5.7 5.9
```



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A Brief History of Chapel



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A Brief History of Chapel: Infancy



Chapel's Infancy: DARPA HPCS (2003–2012)

- ~6–7 FTEs
- Research focus:
 - distinguish locality from parallelism
 - seamlessly mix data- and task-parallelism
 - support user-defined distributed arrays, parallel iterators
- CUG 2013 paper captured post-HPCS project status:

The State of the Chapel Union

Chamberlain, Choi, Dumler, Hildebrandt, Iten, Litvinov, Titus



Crossing the Stream of Adoption







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Crossing the Stream of Adoption: Post-HPCS Barriers







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A Brief History of Chapel: Adolescence



Chapel's Adolescence: "the five-year push" (2013-

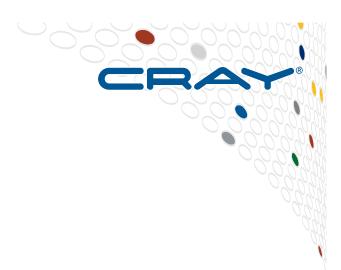
(2013 - 2018)

Motivated by user enthusiasm for Chapel

Then Now

- Development focus:
 - address weak points in HPCS prototype
 - support and grow the Chapel community
- ~13–14 FTEs
- This CUG 2018 talk & paper reports on progress during this time





Chapel Performance: Then vs. Now vs. Reference



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Performance Focus Areas (during 5-year push)



Array Optimizations:

- shifted data optimization (eliminates arbitrary indexing overhead)
- loop-invariant code motion (eliminates meta-data overhead)
- eliminated multiply in indexing for 1D (and innermost dim of 2D+) arrays

Runtime Library Improvements:

- scalable parallel memory allocator
- tasks mapped to affinity aware user-level threads
- native/optimized comm with RDMA and limited software overhead

Optimized Communication:

- compiler locality analysis improvements
- bulk array assignments
- remote-value-forwarding, new distributions, fast-ons, ...



Experimental Methodology



Methodology for the next several slides:

- Resurrected a copy of Chapel 1.7
 - updated it to build with current versions of gcc/g++
- Compared it to Chapel 1.17, released April 2018
- Used today's Cray systems
- Used today's benchmark codes
 - with modest edits for 1.7 in response to language changes



LCALS Serial Kernel



Chapel source:

```
for i in 0..#len do

bvc[i] = cls * (compression[i] + 1.0);
```

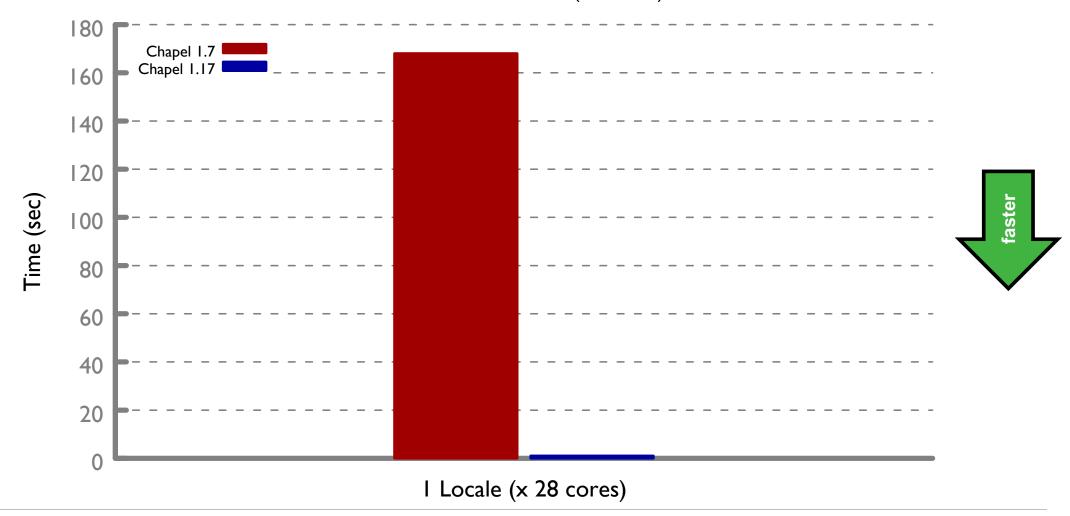
COMPUTE



LCALS Serial Kernel: Chapel Then vs. Now









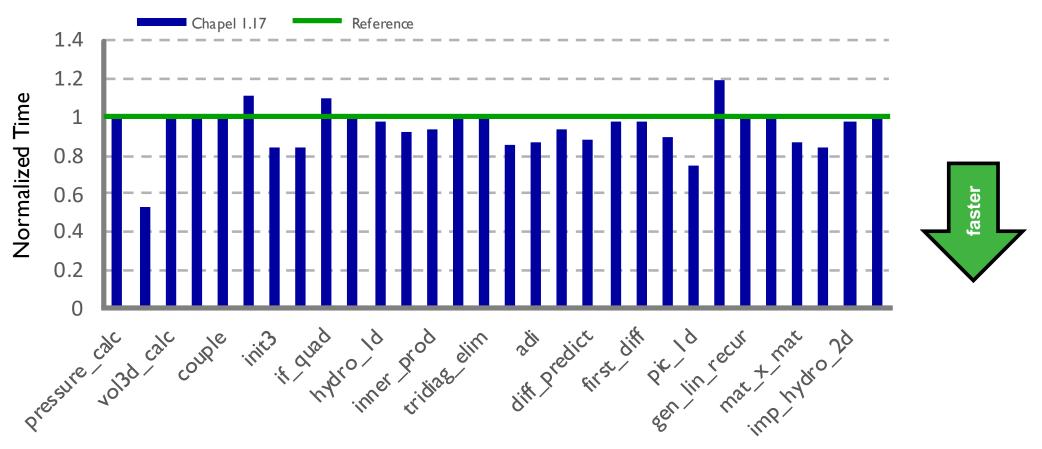
COMPUTE

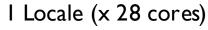
STORE

LCALS Serial Kernels: Chapel Now vs. Ref



LCALS Serial Kernels (Normalized to Ref)



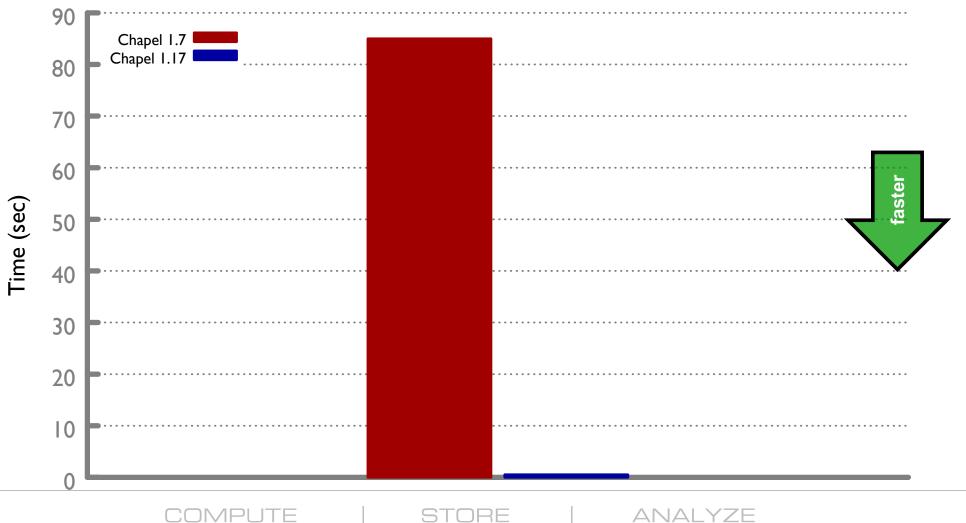




LCALS Parallel Kernel: Chapel Then vs. Now





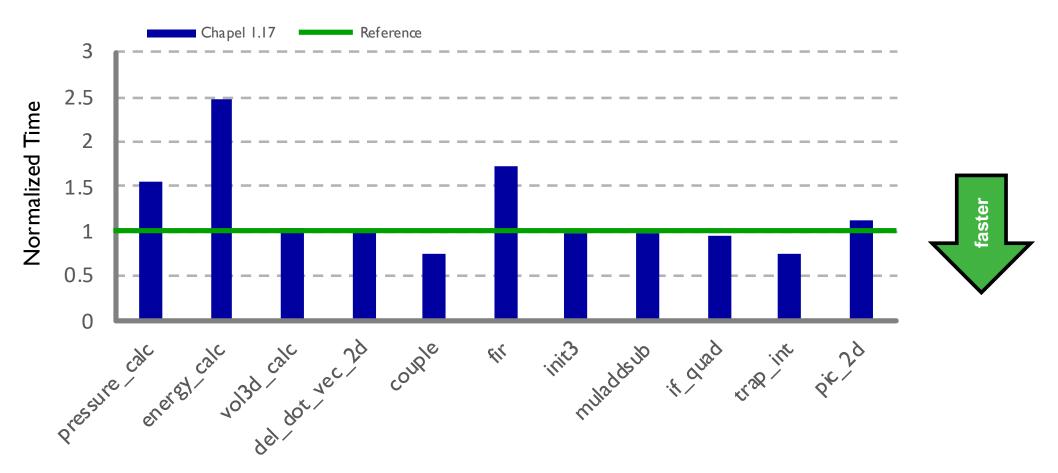




LCALS Parallel Kernels: Chapel Now vs. Ref



LCALS Parallel Kernels (Normalized to Ref)



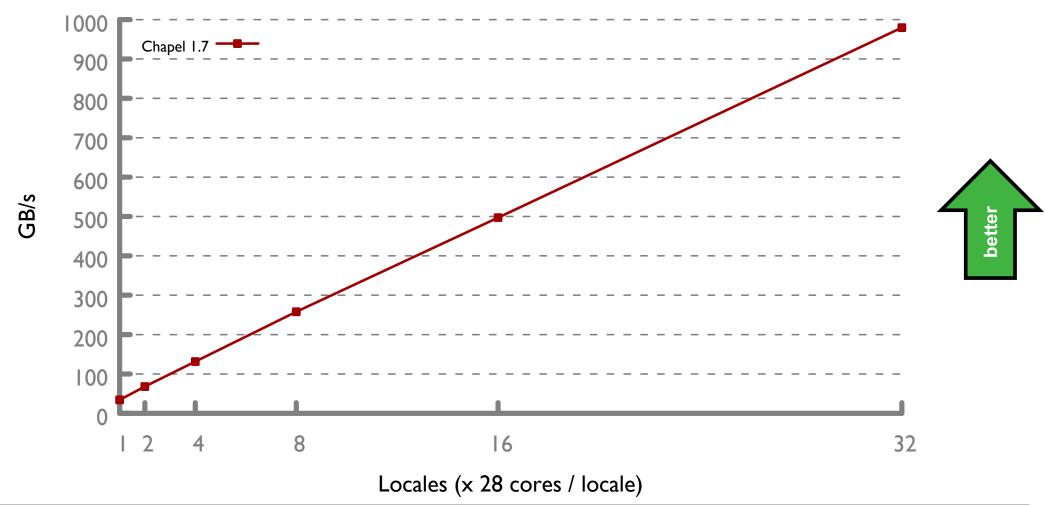
I Locale (x 28 cores)



HPCC STREAM Triad: Chapel Then









COMPUTE |

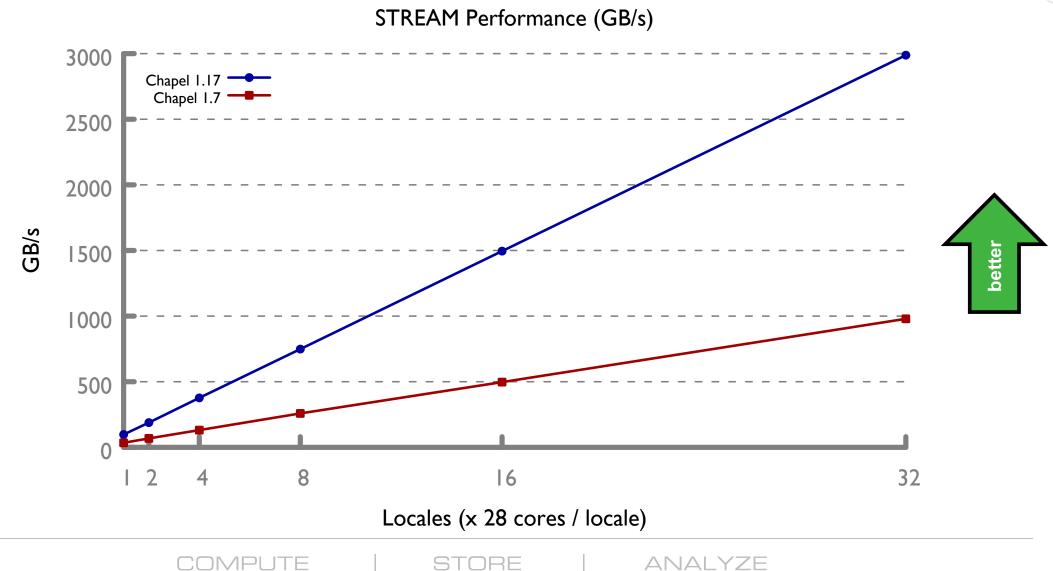
ANALYZE

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STORE

HPCC STREAM Triad: Chapel Then vs. Now



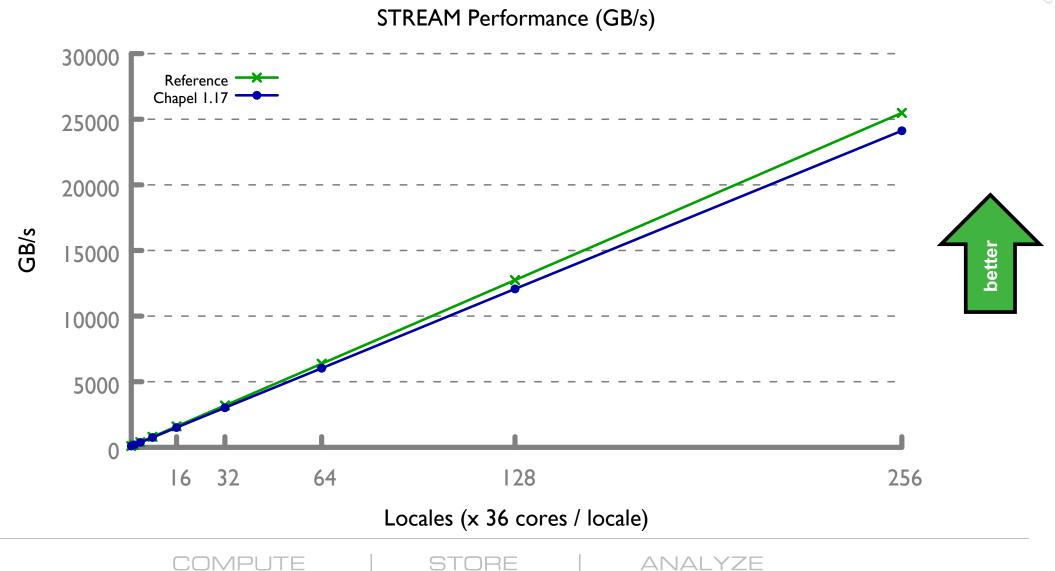




STORE

HPCC STREAM Triad: Chapel Now vs. Ref



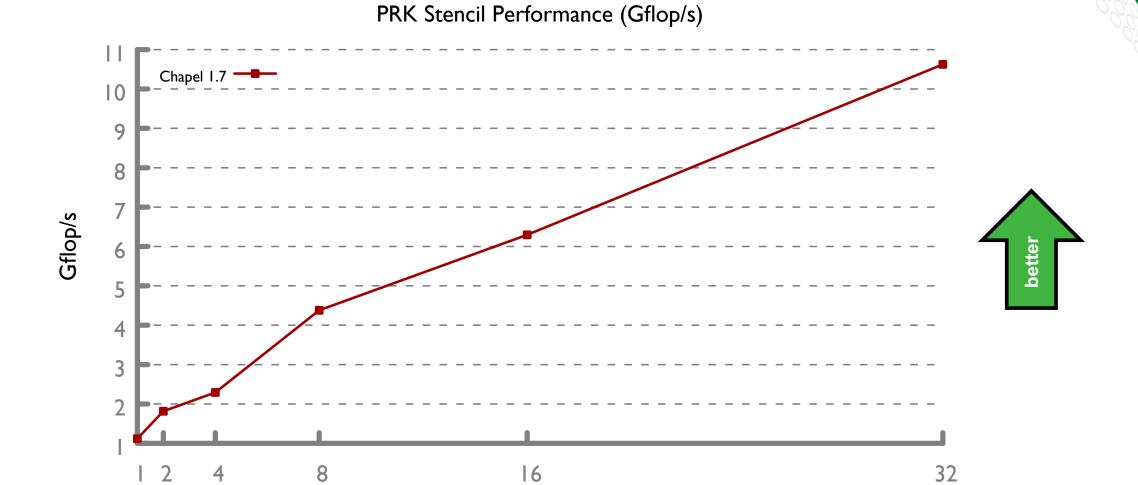




STORE

PRK Stencil: Chapel Then







COMPUTE

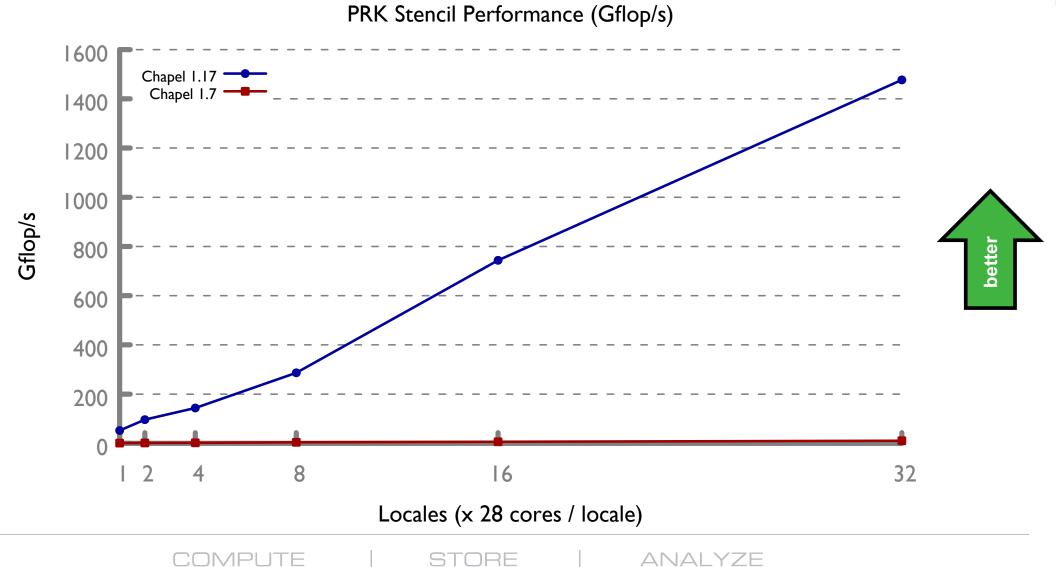
STORE

Locales (x 28 cores / locale)



PRK Stencil: Chapel Then vs. Now



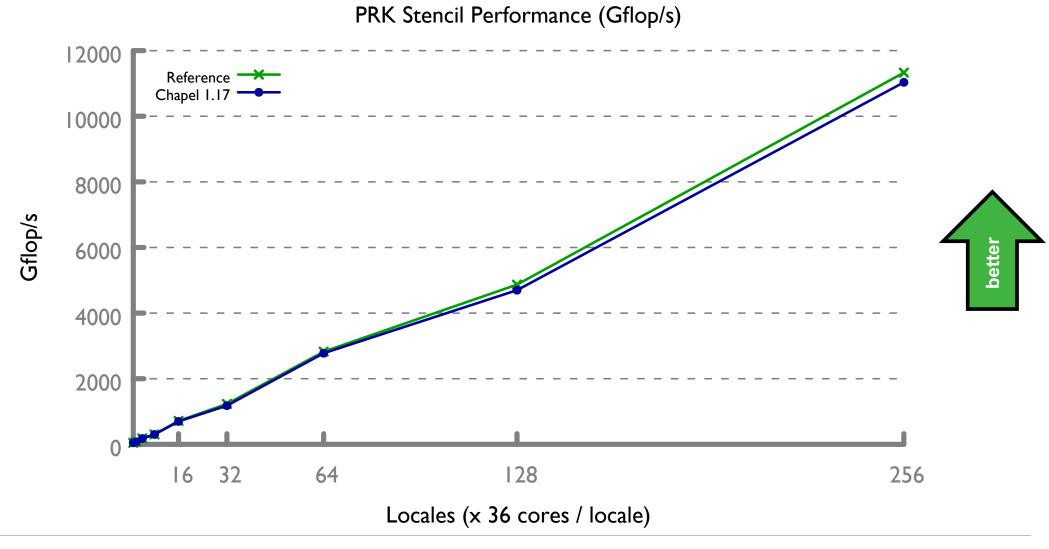




STORE

PRK Stencil: Chapel Now vs. Ref







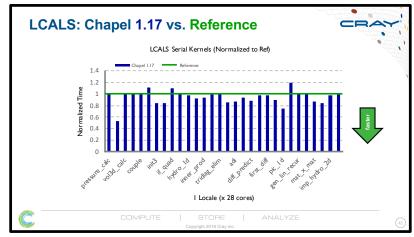
COMPUTE | STORE

ANALYZE

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HPC Patterns: Chapel Now vs. reference





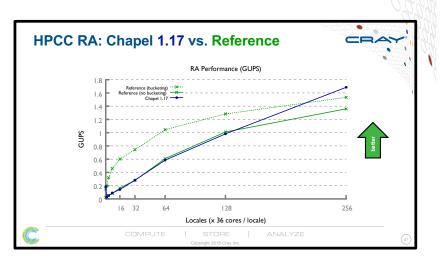
_CALS

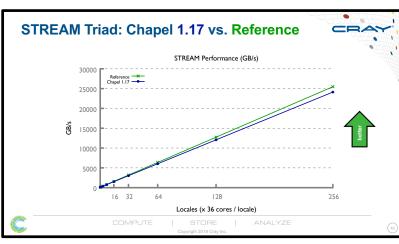
HPCC RA

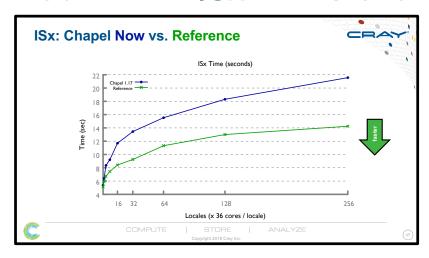
STREAM Triad

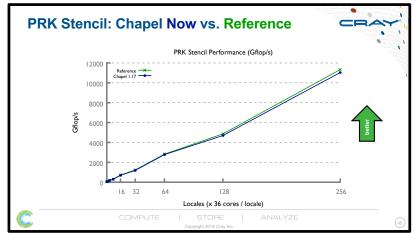
ISx

PRK Stencil











COMPUTE

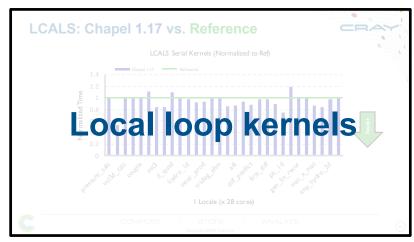
STORE

ANALY

Nightly performance tickers online at: https://chapel-lang.org/perf-nightly.html

HPC Patterns: Chapel Now vs. reference





LCALS

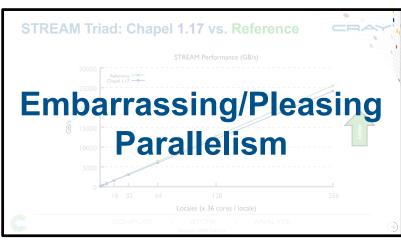
HPCC RA

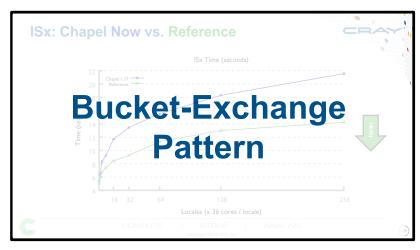
STREAM Triad

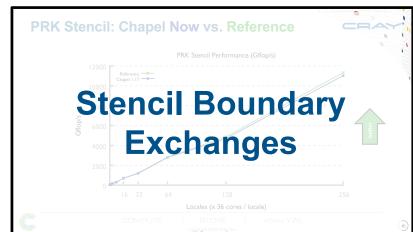
ISx

PRK Stencil





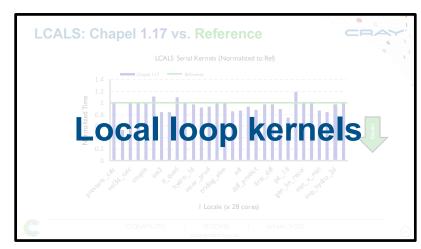






HPC Patterns: Chapel Now vs. reference





LCALS

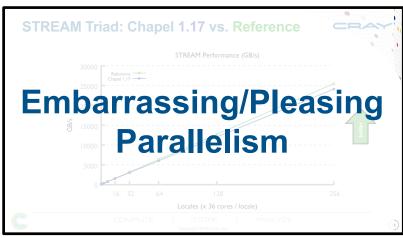
HPCC RA

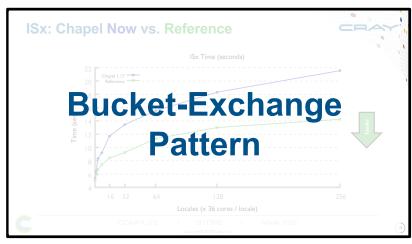
STREAM Triad

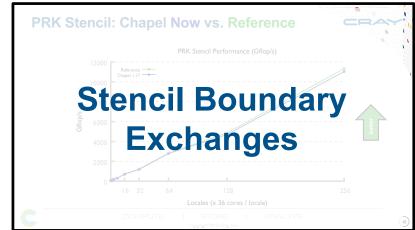
ISx













COMPUTE

ANAL

Nightly performance tickers online at: https://chapel-lang.org/perf-nightly.html

HPCC Random Access Kernel: MPI



```
/* Perform updates to main table. The scalar equivalent is:
                                                                           } else {
                                                                                HPCC InsertUpdate (Ran, WhichPe, Buckets);
    for (i=0; i<NUPDATE; i++) {
                                                                                 pendingUpdates++;
    Ran = (Ran << 1) ^ (((s64Int) Ran < 0) ? POLY : 0);
    Table[Ran & (TABSIZE-1)] ^= Ran;
                                                                               i++;
                                                                             else {
                                                                               MPI Test (&outreq, &have done, MPI STATUS IGNORE);
MPI Irecv(&LocalRecvBuffer, localBufferSize, tparams.dtype64,
                                                                               if (have done)
            MPI ANY SOURCE, MPI ANY TAG, MPI COMM WORLD, &inreg);
                                                                                 outreq = MPI REQUEST NULL;
 while (i < SendCnt) {
                                                                                 pe = HPCC GetUpdates (Buckets, LocalSendBuffer, localBufferSize,
   /* receive messages */
                                                                                                       &peUpdates);
                                                                                 MPI Isend(&LocalSendBuffer, peUpdates, tparams.dtype64, (int)pe,
     MPI Test(&inreq, &have done, &status);
                                                                                           UPDATE TAG, MPI COMM WORLD, &outreq);
     if (have done)
                                                                                 pendingUpdates -= peUpdates;
       if (status.MPI TAG == UPDATE TAG) {
          MPI Get count (&status, tparams.dtype64, &recvUpdates);
          bufferBase = 0;
                                                                           /* send remaining updates in buckets */
          for (j=0; j < recvUpdates; j ++) {
            inmsg = LocalRecvBuffer[bufferBase+j];
                                                                           while (pendingUpdates > 0) {
            LocalOffset = (inmsg & (tparams.TableSize - 1)) -
                                                                             /* receive messages */
                          tparams.GlobalStartMyProc;
            HPCC Table[LocalOffset] ^= inmsg;
                                                                               MPI Test (&inreg, &have done, &status);
                                                                               if (have done) {
        } else if (status.MPI TAG == FINISHED TAG) {
                                                                                 if (status.MPI TAG == UPDATE TAG) {
                                                                                   MPI Get count(&status, tparams.dtype64, &recvUpdates);
         NumberReceiving--;
                                                                                   bufferBase = 0;
         MPI Abort ( MPI COMM WORLD, -1 );
                                                                                   for (j=0; j < recvUpdates; j ++) {</pre>
       MPI Irecv(&LocalRecvBuffer, localBufferSize, tparams.dtype64,
                                                                                     inmsg = LocalRecvBuffer[bufferBase+j];
                  MPI ANY SOURCE, MPI ANY TAG, MPI COMM WORLD, &inreq);
                                                                                     LocalOffset = (inmsg & (tparams.TableSize - 1)) -
                                                                                                    tparams.GlobalStartMyProc;
    } while (have_done && NumberReceiving > 0);
                                                                                     HPCC Table[LocalOffset] ^= inmsg;
    if (pendingUpdates < maxPendingUpdates) {</pre>
     Ran = (Ran << 1) ^ ((s64Int) Ran < ZERO64B ? POLY : ZERO64B);
                                                                                 } else if (status.MPI TAG == FINISHED TAG) {
                                                                                   /* we got a done message. Thanks for playing... */
     GlobalOffset = Ran & (tparams.TableSize-1);
     if ( GlobalOffset < tparams.Top)</pre>
                                                                                   NumberReceiving--;
       WhichPe = (GlobalOffset / (tparams.MinLocalTableSize + 1));
                                                                                   MPI Abort ( MPI COMM WORLD, -1 );
       WhichPe = ( (GlobalOffset - tparams.Remainder) /
                                                                                 MPI Irecv(&LocalRecvBuffer, localBufferSize, tparams.dtype64,
                  tparams.MinLocalTableSize );
      if (WhichPe == tparams.MyProc) {
                                                                                           MPI ANY SOURCE, MPI ANY TAG, MPI COMM WORLD, &inreq);
       LocalOffset = (Ran & (tparams.TableSize - 1)) -
                      tparams.GlobalStartMyProc;
                                                                             } while (have done && NumberReceiving > 0);
        HPCC Table[LocalOffset] ^= Ran;
```

```
MPI Test(&outreq, &have done, MPI STATUS IGNORE);
  if (have done) {
   outreq = MPI REQUEST NULL;
    pe = HPCC GetUpdates (Buckets, LocalSendBuffer, localBufferSize,
                         &peUpdates);
    MPI Isend(&LocalSendBuffer, peUpdates, tparams.dtype64, (int)pe,
              UPDATE TAG, MPI COMM WORLD, &outreg);
    pendingUpdates -= peUpdates;
/* send our done messages */
for (proc count = 0; proc count < tparams.NumProcs; ++proc count) {
  if (proc count == tparams.MyProc) { tparams.finish req[tparams.MyProc] =
                                      MPI REQUEST NULL; continue; }
  /* send garbage - who cares, no one will look at it */
  MPI Isend(@Ran, 0, tparams.dtype64, proc count, FINISHED TAG,
           MPI COMM WORLD, tparams.finish req + proc count);
/* Finish everyone else up... */
while (NumberReceiving > 0) {
  MPI Wait (&inreg, &status);
  if (status.MPI TAG == UPDATE TAG) {
   MPI Get count (&status, tparams.dtype64, &recvUpdates);
    bufferBase = 0;
    for (j=0; j < recvUpdates; j ++) {
      inmsg = LocalRecvBuffer[bufferBase+j];
      LocalOffset = (inmsg & (tparams.TableSize - 1)) -
                     tparams.GlobalStartMyProc;
      HPCC Table[LocalOffset] ^= inmsg;
  } else if (status.MPI TAG == FINISHED TAG) {
   /* we got a done message. Thanks for playing... */
   NumberReceiving--;
  } else {
   MPI Abort ( MPI COMM WORLD, -1 );
  MPI Irecv(&LocalRecvBuffer, localBufferSize, tparams.dtype64,
            MPI ANY SOURCE, MPI ANY TAG, MPI COMM WORLD, &inreq);
MPI Waitall( tparams.NumProcs, tparams.finish req, tparams.finish statuses);
```



COMPLITE

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HPCC Random Access Kernel: MPI



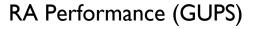
```
Chapel Kernel
/* Perform updates to main table. The scalar equivalent is:
 for (i=0; i<NUPDATE; i++) {
                               forall ( , r) in zip(Updates, RAStream()) do
  Ran = (Ran << 1) ^ (((s64Int) Ran < 0) ? POLY : 0);
  Table[Ran & (TABSIZE-1)] \(^=\) Ran;
                                  T[r & indexMask] ^= r;
                                             MPI Comment
          Perform updates to main table. The scalar equivalent is:
        *
                for (i=0; i<NUPDATE; i++) {
                   Ran = (Ran << 1) ^ (((s64Int) Ran < 0) ? POLY : 0);
                   Table[Ran & (TABSIZE-1)] ^= Ran;
```

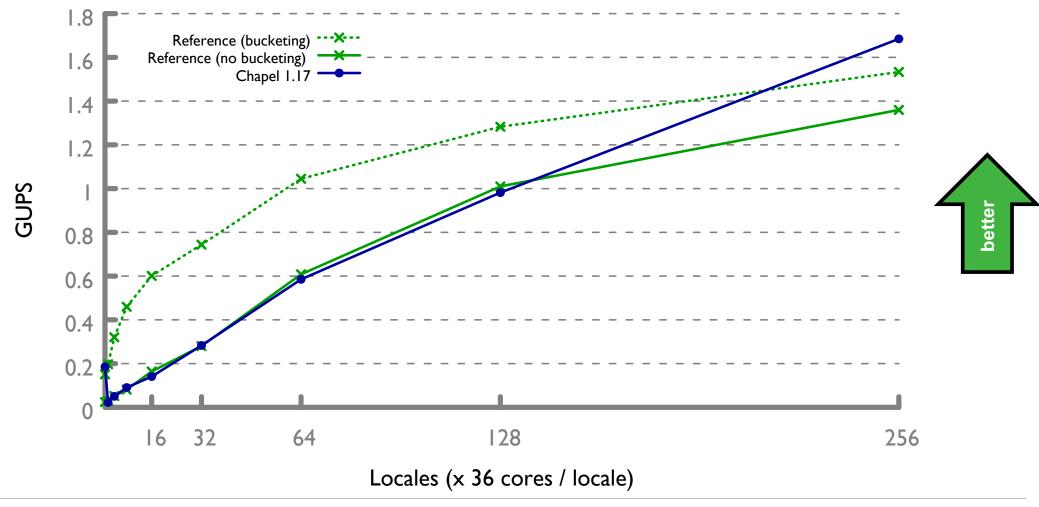


 \triangle

HPCC RA: Chapel Now vs. Ref









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Memory Leaks: Then vs. Now (skipped at CUG due to time constraints)

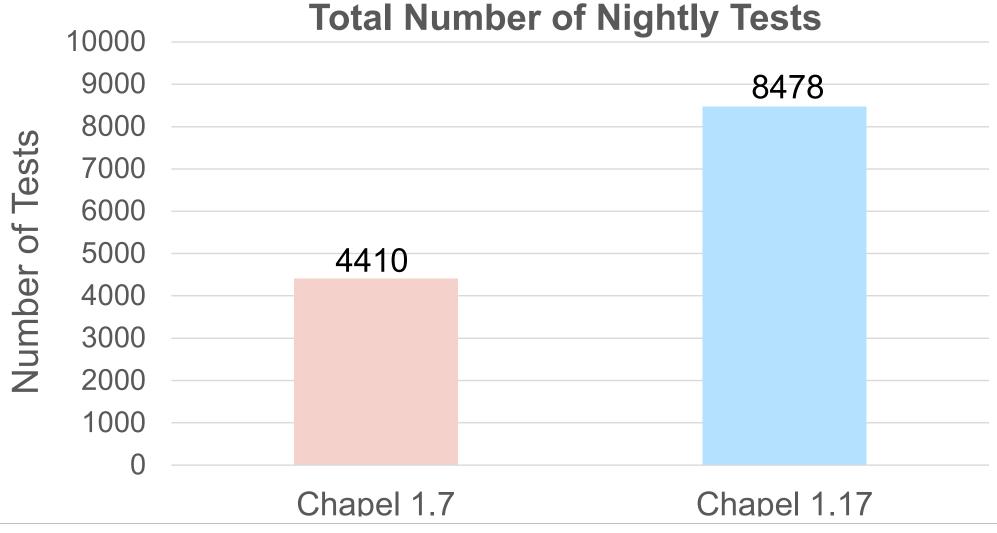




Memory Leaks: Chapel Then vs. Now

COMPUTE

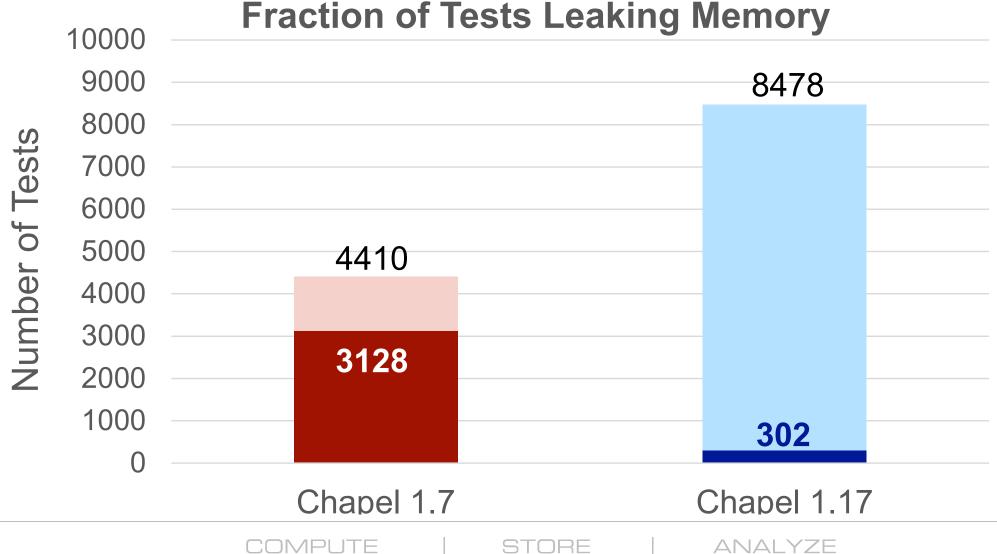






Memory Leaks: Chapel Then vs. Now







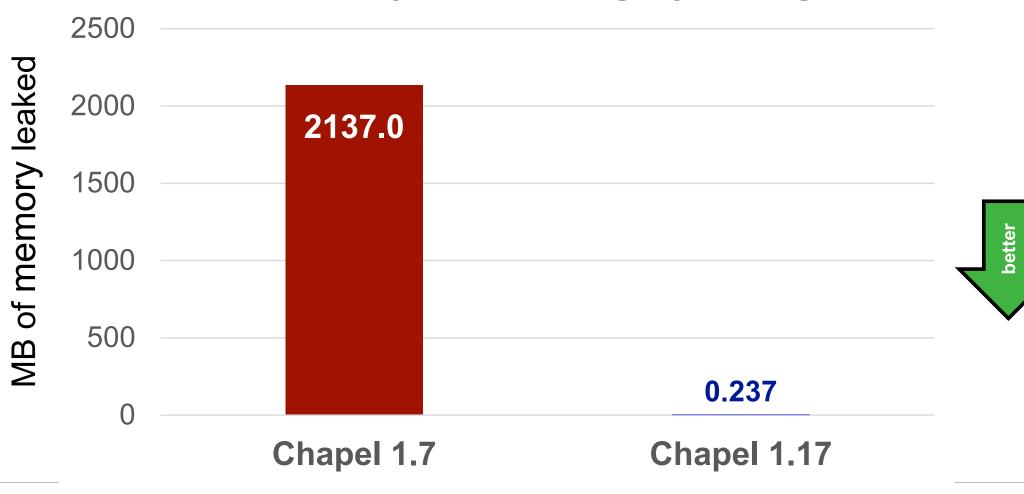
ANALYZE

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Memory Leaks: Chapel Then vs. Now









COMPUTE

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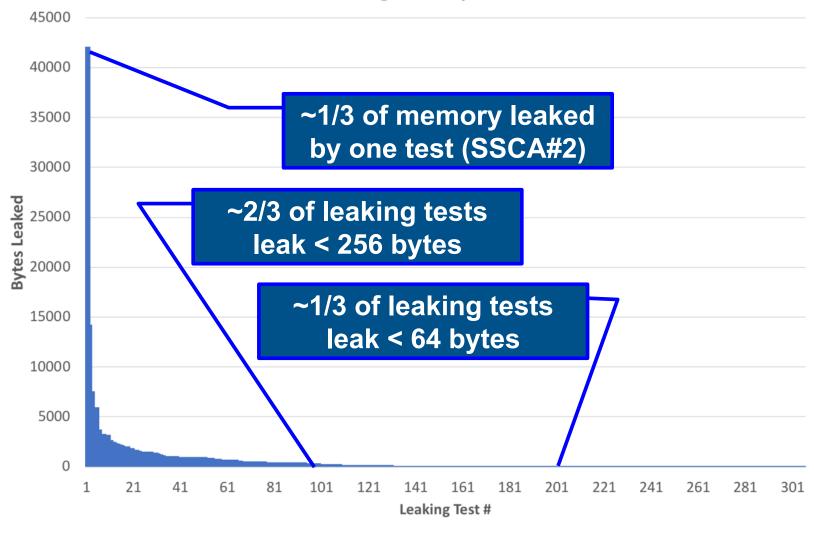


Memory Leaks: Remaining Leaks

COMPUTE





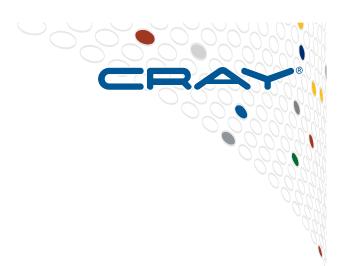




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Chapel Language: Then vs. Now



Language: Then



Parallelism and Locality: Generally in good shape

not many changes here since HPCS

Base Language: Left much to be desired

lots of focus here since HPCS



Language: Now



Parallelism and Locality

- introduced task intents to reduce chances of race conditions
- and user-defined locale models to support new node architectures

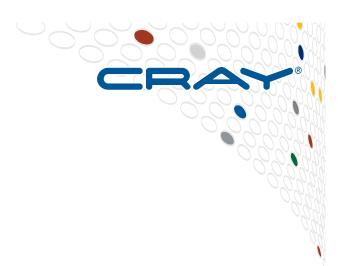
Base Language

- fixed a number of problems with object-oriented programming
 - records: poor memory management discipline
 - classes: problems with generic classes, class hierarchies
- made strings usable
- added error-handling features
- made namespace improvements

COMPUTE

(and much more...)





Chapel Ecosystem: Then vs. Now



Documentation: Then

After HPCS:

a PDF language specification

a Quick Reference sheet

- a number of READMEs
- ~22 primer examples

Chapel Language Specification Version 0.93

Cray Inc 901 Fifth Avenue, Suite 1000

April 18, 2013



```
Quick Start
                                                          Expression Precedence and Associativi
  Task Parallel Primer
 /
/ This primer illustrates Chapel's parallel tasking features,
/ namely the begin, cobegin, and coforall statements.
writeln("1: ### The begin statement ###");
  / The begin statement spawns a thread of execution that is independent
 / of the current (main) thread of execution.
egin writeln("1: output from spawned task");
  The main thread of execution continues on to the next statement.
  There is no guarantee as to which statment will execute first.
 riteln("1: output from main task");
writeln("2: ### The cobegin statement ###"):
 / For more structured behavior, the cobegin statement can be used to
  spawm a block of tasks, one for each statement. Control continues after the cobegin block, but only after all the tasks within the
  cobegin block have completed.
  writeln("2: output from spawned task 1");
 writeln("2: output from spawned task 2");
  / The output from within the cobegin statement will always precede the
// following output from the main thread of execution.
writeln("2: output from main task");
writeln("3: ### The cobegin statement with nested begin statements ###");
 / If any begin statements are used within a cobegin statement,
  the thread of execution does not a wait for those begin statements
```

Chapel Quick Reference

label outer for ... break; ar break oute

Default Values for Formal Arguments proc foo(argl: int, argl: real = 3.14);



COMPUTE

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ANALYZE

/ to complete.

writeln("3: output from spawned task 1");

writeln("3: output from spawned task 2");

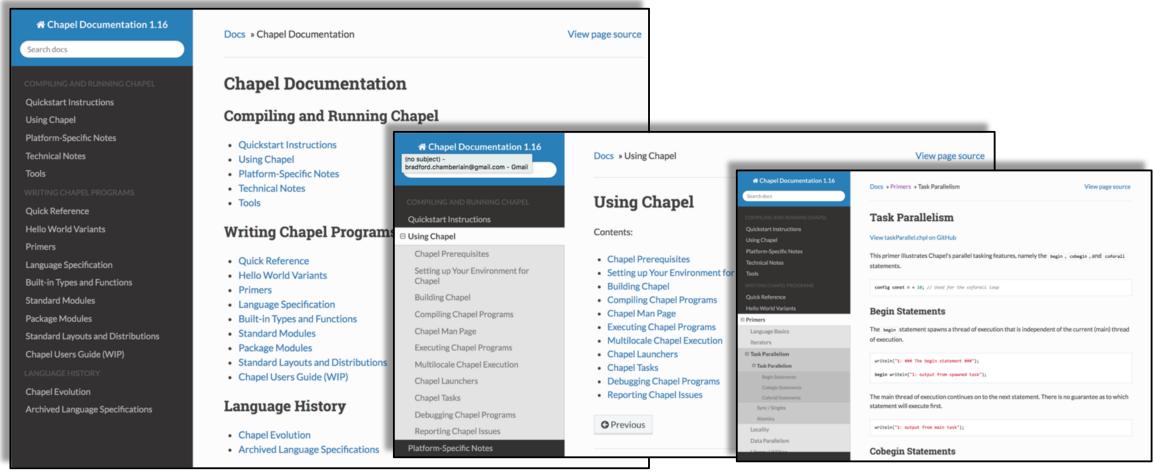
uu-:---F1 taskParallel.chpl Top L1 (Chapel/l Abbrev)-----



Documentation: Now



Now: 200+ modern, hyperlinked, web-based documentation pages





COMPUTE

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Libraries: Then

After HPCS: ~25 library modules

documented via source comments, if at all:

```
File Edit Options Buffers Tools chpl Help
// Random Module
// This standard module contains a random number generator based on
// the one used in the NPB benchmarks. Tailoring the NPB comments to
// this code, we can say the following:
    This generator returns uniform pseudorandom real values in the
   range (0, 1) by using the linear congruential generator
      x_{k+1} = a x_k \pmod{2**46}
    where 0 < x_k < 2**46 and 0 < a < 2**46. This scheme generates
    2**44 numbers before repeating. The seed value must be an odd 64-bit integer in the range (1, 2^46). The generated values are
    normalized to be between 0 and 1, i.e., 2**(-46) * x_k.
    This generator should produce the same results on any computer
   with at least 48 mantissa bits for real(64) data.
  1. We would like to support general serial and parallel iterators
// on the RandomStream class, but this is not possible with our
// current parallel iterator framework.
// 2. The random number generation functionality in this module is
// currently restricted to 64-bit real, 64-bit imag, and 128-bit
// complex values. This should be extended to other primitive types
 / for which this would make sense. Coercions are insufficient.
// 3. Can the multiplier 'arand' be moved into the RandomStream class
// so that it can be changed by a user of this class.
// 4. By default, the random stream seed is initialized based on the
// current time in microseconds, allowing for some degree of
// randomness. The intent of the SeedGenerator enumerated type is to
// provide a menu of options for initializing the random stream seed,
// but only one option is implemented to date.
// It is the intent that once Chapel supports the notion of 'private',
-uu-:---F1 Random.chpl Top L1 (Chapel/l Abbrev)-----
```

COMPLITE

```
File Edit Options Buffers Tools chpl Help
  Copyright (c) 2004-2013, Cray Inc. (See LICENSE file for more details)
 xtern type qio_regexp_t;
 xtern record qio_regexp_options_t {
 var utf8:bool;
 var posix:bool;
 var literal:bool;
 var nocapture:bool;
 // These ones can be set inside the regexp
 var ignorecase:bool; // (?i)
 var multiline:bool; // (?m)
 var dotnl:bool; // (?s)
 var nongreedy:bool; // (?U)
 ctern proc qio_regexp_null():qio_regexp_t;
 ctern proc qio_regexp_init_default_options(ref options:qio_regexp_options_t);
 ctern proc qio_regexp_create_compile(str:string, strlen:int(64), ref options:q\)
io_regexp_options_t, ref compiled:qio_regexp_t);
       proc qio_regexp_create_compile_flags(str:string, strlen:int(64), flags:s
 ring, flagslen:int(64), isUtf8:bool, ref compiled:qio_regexp_t);
      proc qio_regexp_create_compile_flags_2(str:c_ptr, strlen:int(64), flags:\)
c_ptr, flagslen:int(64), isUtf8:bool, ref compiled:qio_regexp_t);
extern proc qio_regexp_retain(ref compiled:qio_regexp_t);
 xtern proc qio_regexp_release(ref compiled:qio_regexp_t);
 xtern proc qio_regexp_get_options(ref regexp:qio_regexp_t, ref options: qio_re\
gexp_options_t);
 ctern proc qio_regexp_get_pattern(ref regexp:qio_regexp_t, ref pattern: string)
 xtern proc qio_regexp_get_ncaptures(ref regexp:qio_regexp_t):int(64);
 xtern proc qio_regexp_ok(ref regexp:qio_regexp_t):bool;
 xtern proc qio_regexp_error(ref regexp:qio_regexp_t):string;
 xtern const QIO_REGEXP_ANCHOR_UNANCHORED:c_int;
 xtern const QIO_REGEXP_ANCHOR_START:c_int;
 xtern const QIO_REGEXP_ANCHOR_BOTH:c_int;
 xtern record qio_regexp_string_piece_t {
 var offset:int(64); // counting from 0, -1 means "NULL"
 var len:int(64);
 ktern proc qio_regexp_string_piece_isnull(ref sp:qio_regexp_string_piece_t):bo'
-uu-:---F1 Regexp.chpl Top L1 (Chapel/l Abbrev)--
```

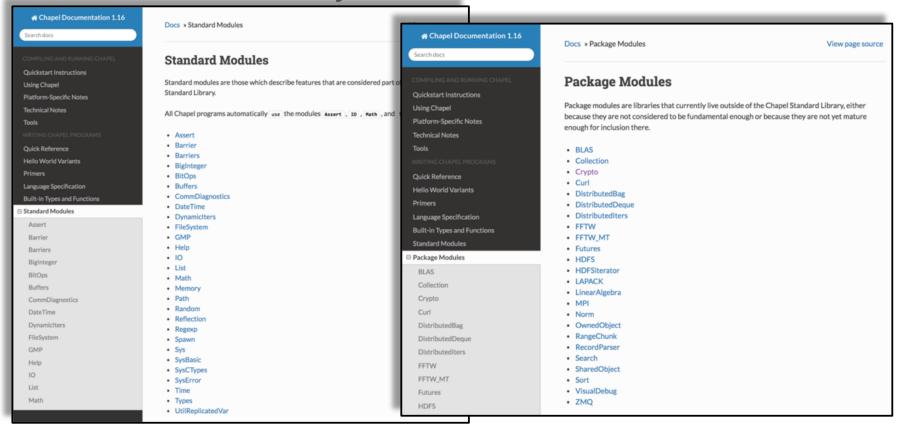


Libraries: Now

Now: ~60 library modules

web-documented, many user-contributed

COMPUTE





Libraries: Now



Math: FFTW, BLAS, LAPACK, LinearAlgebra, Math

Inter-Process Communication: MPI, ZMQ (ZeroMQ)

Parallelism: Futures, Barrier, DynamicIters

Distributed Computing: DistributedIters, DistributedBag, DistributedDeque, Block, Cyclic, Block-Cyclic, ...

File Systems: FileSystem, Path, HDFS

Others: BigInteger, BitOps, Crypto, Curl, DateTime, Random, Reflection, Regexp, Search, Sort, Spawn, ...



Tools: Then

After HPCS:

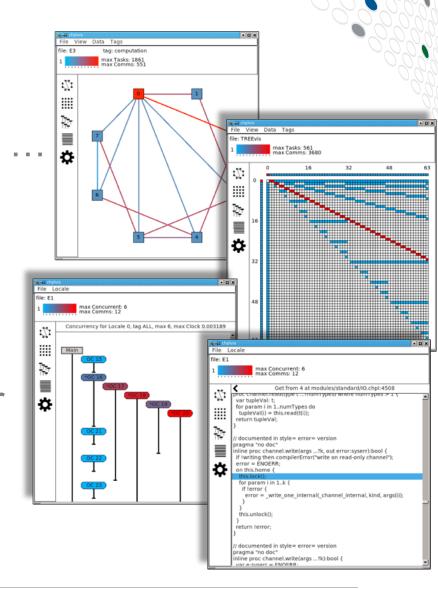
- highlighting modes for emacs and vim
- chpldoc: documentation tool (early draft)



Tools: Now

Now:

- highlighting modes for emacs, vim, atom, ...
- chpldoc: documentation tool
- mason: package manager
- c2chapel: interoperability aid
- bash tab completion: command-line help
- chplvis: performance visualizer / debugger





Then vs. Now: And so much more...



Interoperability:

passing arrays & functions to C, working with C pointers, ...

Development process:

GitHub, Jenkins, Travis, interactive nightly performance graphs...

Social media: Twitter, Facebook, YouTube

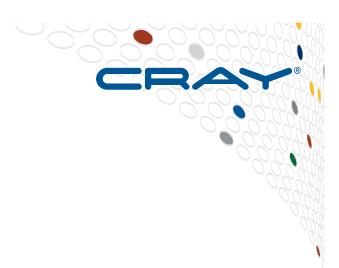
User support: GitHub issues, StackOverflow, Gitter, email

Web presence: CLBG, Try It Online, CyberDojo, ...

Memory Leaks: significantly reduced

CHIUW: annual community workshop





Chapel User Profiles

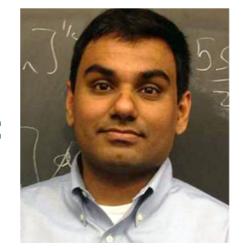




Chapel User Profiles



Current Users:



Time-to-science Cosmologist



Commercial Al Scientist

Potential Users:



Genomic Researcher



DOE Scientist



Chapel User Profiles



Current Users:

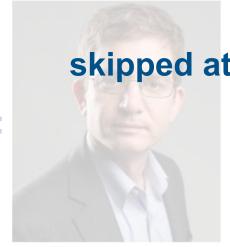


Time-to-science Cosmologist



Commercial Al Scientist

Potential Users:



Genomic Researcher



DOE Scientist



User Profile: Time-to-Science Cosmologist





Name: Nikhil Padmanabhan

Title: Associate Professor of Physics and Astronomy, Yale University

Computations: Surveys of galaxies to constrain cosmological models, n-body simulations of gravity

Why Chapel? "My interests in Chapel developed from a desire to have a lower barrier to writing parallel codes. In particular, I often find myself writing prototype codes (often serial), but then need to scale these codes to run on large numbers of simulations/datasets. Chapel allows me to smoothly transition from serial to parallel codes with a minimal number of changes.

"Another important issue for me is "my time to solution" (some measure of productivity vs performance). Raw performance is rarely the only consideration."



User Profile: Commercial AI Scientist





Name: Brian Dolan DEEP 6 A

Title: Co-Founder and Chief Scientist of Deep 6 Al

Computations: Natural language processing, Al and ML applications, network analysis, community detection, reinforcement learning in the form of Deep Q-Networks

Why Chapel? "I have used Fortran, R, Java and Python extensively. If I had to give up Chapel, I would probably move to C++. I prefer Chapel due to the extreme legibility and performance. We have abandoned Python on large problems for performance reasons.

"We've now developed thousands of lines of Chapel code and a half dozen open source libraries for things like database connectivity, numerical libraries, graph processing, and even a REST framework. We've done this because AI is about to face an HPC crisis, and the folks at Chapel understand the intersection of usability and scalability."



Potential User Profile: Genomic Researcher





Name: Jonathan Dursi

Title: Senior Research Associate, The Hospital for Sick Children, Toronto

Computations: Human genomics, bioinformatics, and medical informatics

Why Chapel? "My interest in Chapel lies in its potential for bioinformatics tools that are currently either written in elaborately crafted, threaded but single node, C++ code, or in Python. Either has advantages and disadvantages (performance vs rapid development cycles), but neither has a clear path to cross-node computation, for performance as well as larger memory and memory bandwidth. Chapel has the potential to have some of the best of both worlds in terms of C++ and Python, as well as having a path to distributed memory."



Potential User Profile: DOE Scientist

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Name: Anshu Dubey

Title: Computer Scientist, Argonne National Laboratory

Computations: Design and development of Multiphysics software that can serve multiple science domains; solvers for PDEs and ODEs

Why Chapel? "In Multiphysics applications separation of concerns and use of high level abstractions is critical for sustainable software. Chapel combines language features that would enable this for clean implementation.

"HPC Scientific software is made more complex than it needs to be because the only language designed for scientific work, Fortran, is losing ground for various reasons. Its object oriented features are clunky and make it nearly as unsuitable as other languages for scientific work. Chapel appears to be parallel and modern Fortran done better, therefore has the potential to become a more suitable language."



Chapel and Productivity

Chapel aims to be as...

- ...programmable as Python
- ...fast as Fortran
- ...scalable as MPI, SHMEM, or UPC
- ...portable as C
- ...flexible as C++
- ...fun as [your favorite language]



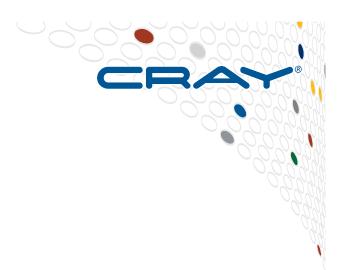








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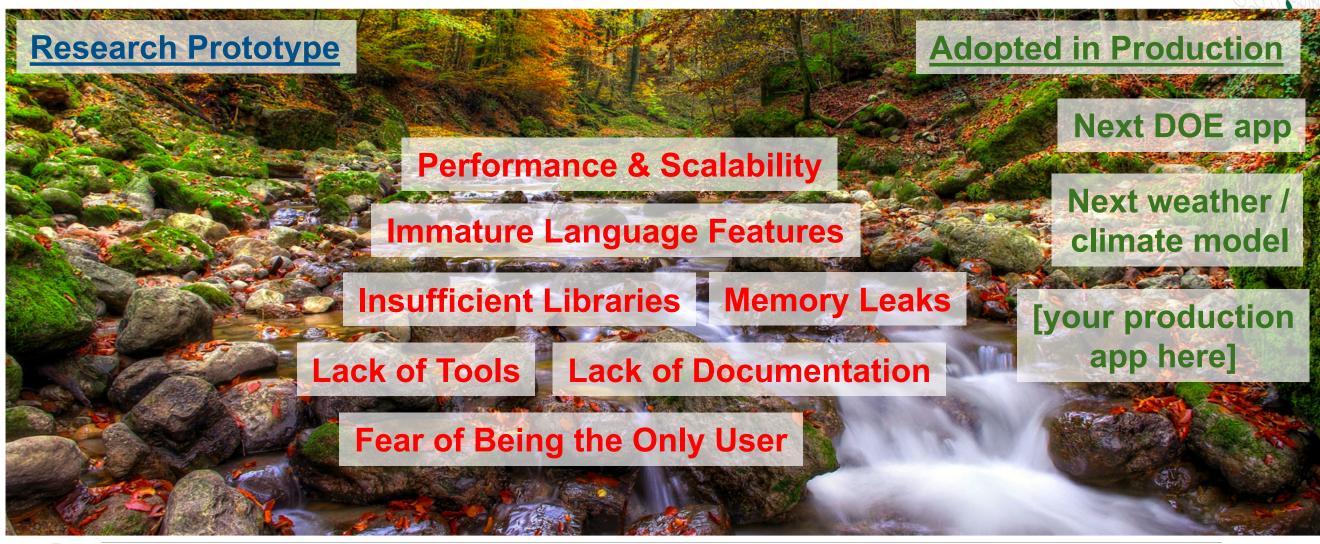






Crossing the Stream of Adoption







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Crossing the Stream of Adoption



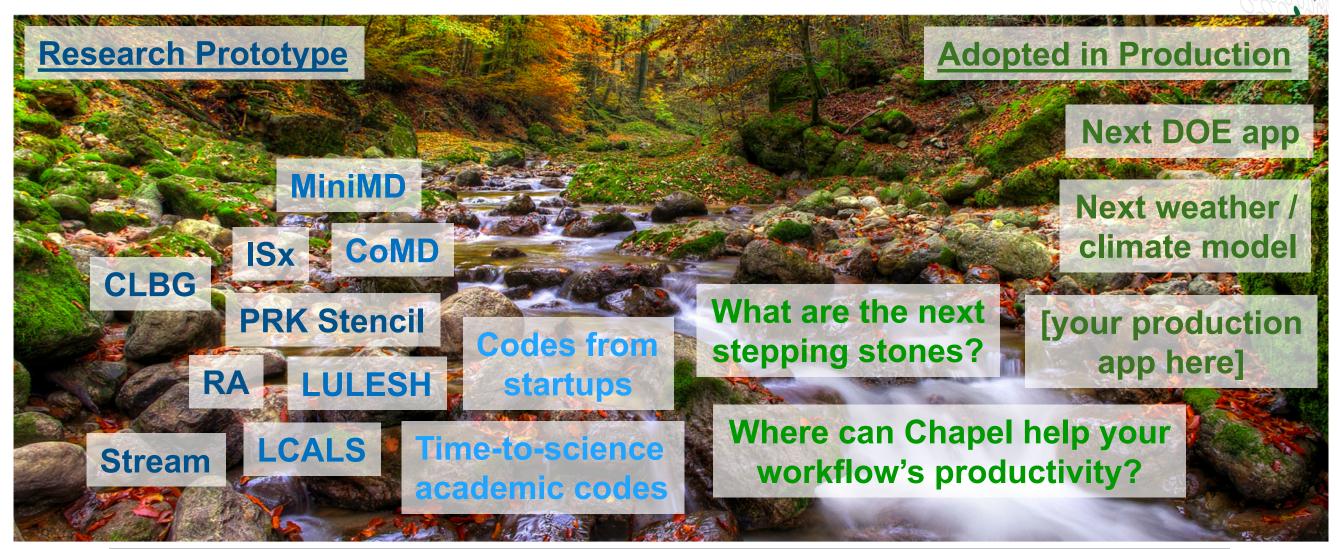




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Crossing the Stream of Adoption







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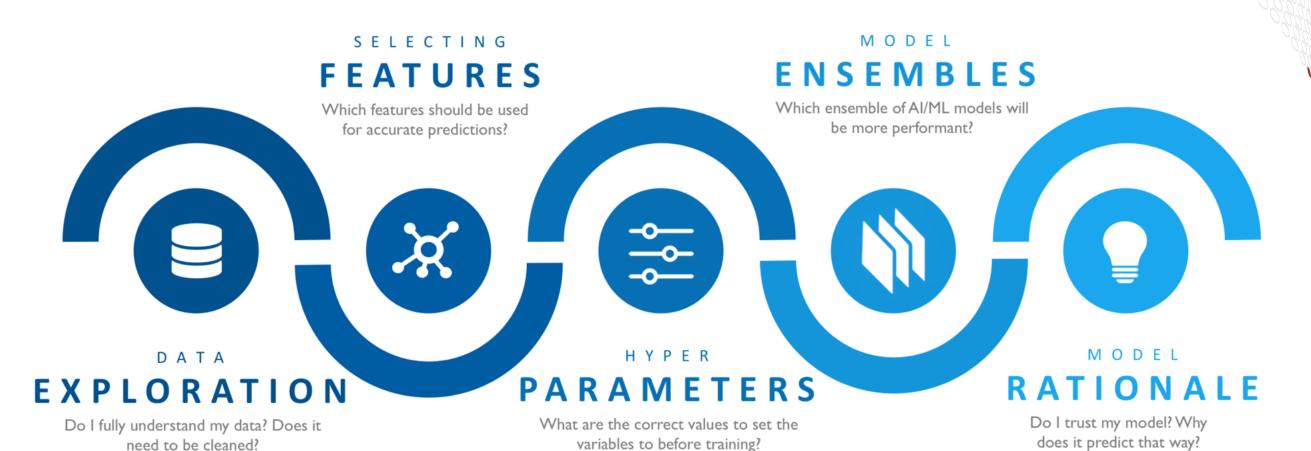
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Discovery Roadblocks

Data Science Pain Points





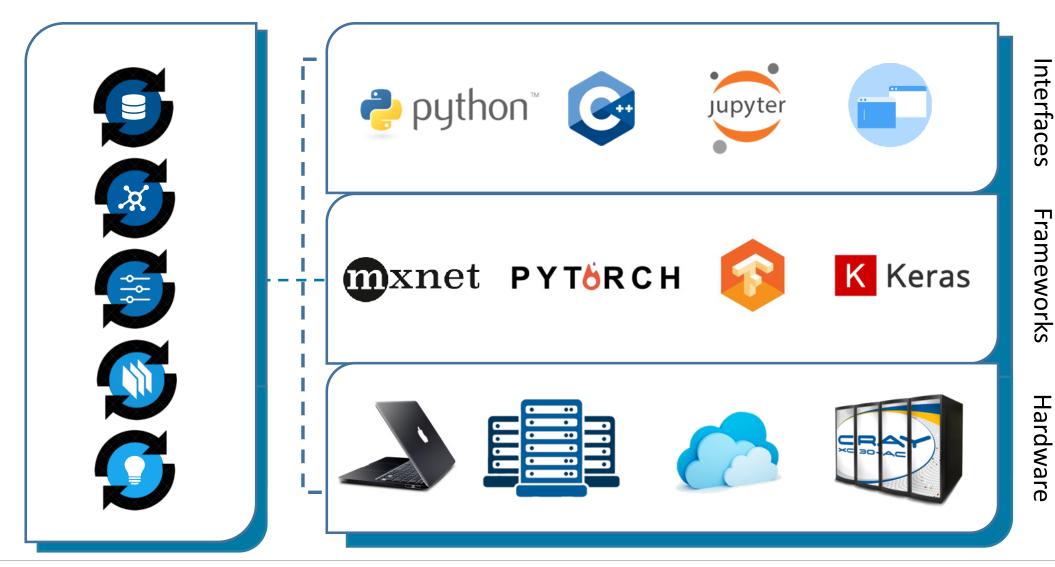
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Chapel Al Ecosystem







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Sample Chapel Al Workflow



User works from within a Jupyter notebook



- Uses Chapel to ingest large HDF5 data files
 - read in parallel
 - transformed / analyzed during ingestion
 - stored in a distributed Dataframe



- Starts working on model locally on laptop
- As confidence in model grows, tunes it at scale
 - feature selection
 - hyperparameter optimization







Chapel's college years: plans for 2018-2021

- Language Core
- Interoperability / Usability

- **Portability**
- **Data Ingestion**
- **Chapel Al**





Chapel's college years: plans for 2018-2021

- **Language Core**
 - Language stabilization: avoid backward-breaking changes
 - Sparse array improvements, partial reductions, delete-free features, ...
 - Additional performance and scalability improvements
- Interoperability / Usability

- **Portability**
- **Data Ingestion**
- Chapel Al





Chapel's college years: plans for 2018-2021

- Language Core
- Interoperability / Usability
 - Python / C++ interoperability
 - Support for Jupyter notebooks / REPL
- **Portability**
- **Data Ingestion**
- Chapel Al





Chapel's college years: plans for 2018-2021

- Language Core
- Interoperability / Usability
- Portability
 - LLVM back-end
 - Target Libfabric/OFI
 - Target GPUs
 - Cloud computing support
- Data Ingestion
- Chapel Al





Chapel's college years: plans for 2018-2021

- Language Core
- Interoperability / Usability
- Portability
- Data Ingestion
 - Support HDF5, NetCDF, CSV, ...
 - Transform-on-ingest
 - Distributed DataFrames support

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Chapel Al





Chapel's college years: plans for 2018-2021

- Language Core
- Interoperability / Usability
- **Portability**
- **Data Ingestion**
- **Chapel Al**
 - Hyperparameter optimization
 - Deep Learning



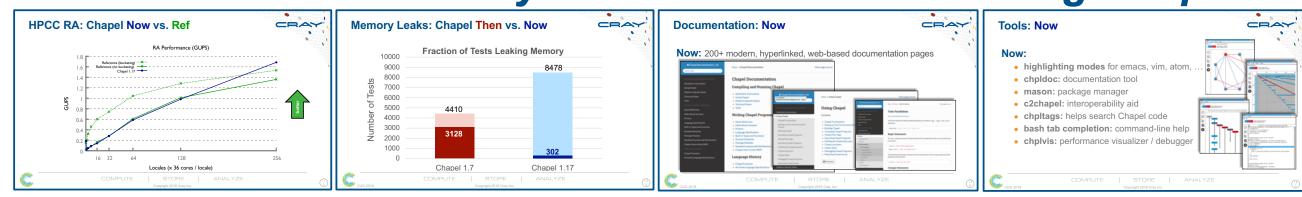


Summary



Chapel has made huge strides over the past five years

We've addressed many historical barriers to using Chapel



We're continuing our work to support and improve Chapel

We're looking for the next generation of Chapel users, as well as concrete use cases for AI / ML

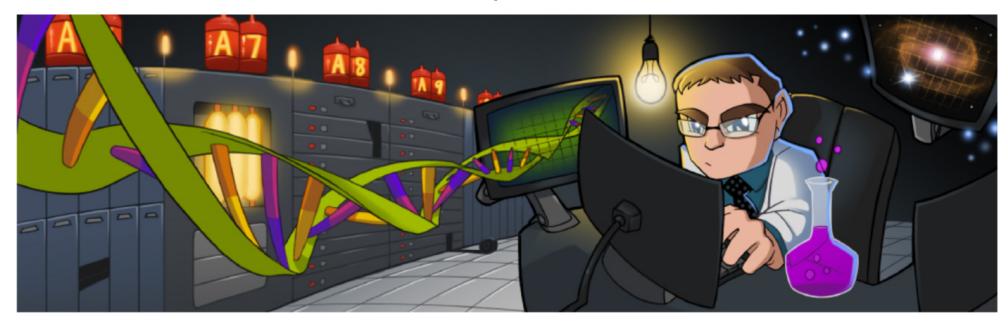


CHIUW 2017 Keynote



Chapel's Home in the Landscape of New Scientific Computing Languages (and what it can learn from the neighbours)

Jonathan Dursi, The Hospital for Sick Children, Toronto





Quote from CHIUW 2017 keynote



"My opinion as an outsider...is that Chapel is important, Chapel is mature, and Chapel is just getting started.

"If the scientific community is going to have frameworks...that are actually designed for our problems, they're going to come from a project like Chapel.

"And the thing about Chapel is that the set of all things that are 'projects like Chapel' is 'Chapel."

-Jonathan Dursi

Chapel's Home in the New Landscape of Scientific Frameworks
(and what it can learn from the neighbours)

CHIUW 2017 keynote

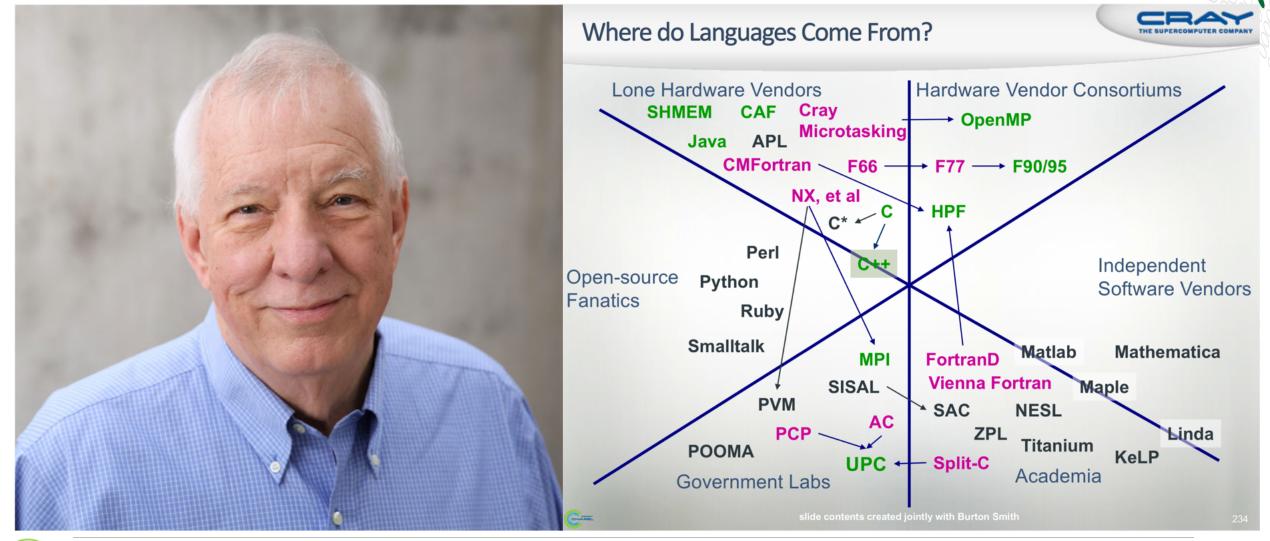
https://ljdursi.github.io/CHIUW2017 / https://www.youtube.com/watch?v=xj0rwdLOR4U



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Dedicated to the Memory of Burton Smith

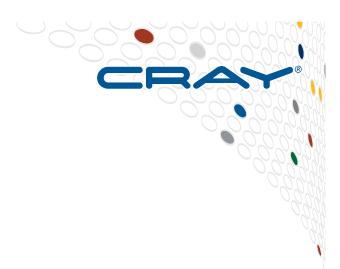






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Chapel Resources





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Chapel Central

https://chapel-lang.org

- downloads
- documentation
- resources
- presentations
- papers



Home What is Chapel?

What's New?

How Can I Learn Chapel? **Contributing to Chapel**

Documentation

Download Chapel Try It Now Release Notes

User Resources Educator Resources
Developer Resources

Social Media / Blog Posts

Presentations Tutorials
Publications and Papers

Contributors / Credits
Research / Collaborations

chapel-lang.org chapel_info@cray.com



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The Chapel Parallel Programming Language

What is Chapel?

Chapel is a modern programming language that is...

- parallel: contains first-class concepts for concurrent and parallel computation
- productive: designed with programmability and performance in mind
- · portable: runs on laptops, clusters, the cloud, and HPC systems
- · scalable: supports locality-oriented features for distributed memory systems
- · open-source: hosted on GitHub, permissively licensed

New to Chapel?

As an introduction to Chapel, you may want to...

- read a blog article or book chapter
- · watch an overview talk or browse its slides
- · download the release
- browse sample programs
- · view other resources to learn how to trivially write distributed programs like this:

```
// use the Cyclic distribution library
use CyclicDist;
config const n = 100;
                         // use --n=<val> when executing to override this default
forall i in {1..n} dmapped Cyclic(startIdx=1) do
 writeln("Hello from iteration ", i, " of ", n, " running on node ", here.id);
```

What's Hot?

- Chapel 1.17 is now available—download a copy or browse its release notes
- The advance program for CHIUW 2018 is now available—hope to see you there!
- Chapel is proud to be a Rails Girls Summer of Code 2018 organization
- Watch talks from ACCU 2017, CHIUW 2017, and ATPESC 2016 on YouTube
- . Browse slides from SIAM PP18, NWCPP, SeaLang, SC17, and other recent talks
- Also see: What's New?



Chapel Social Media (no account required)



http://twitter.com/ChapelLanguage

http://facebook.com/ChapelLanguage

https://www.youtube.com/channel/UCHmm27bYjhknK5mU7ZzPGsQ/





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Chapel Community

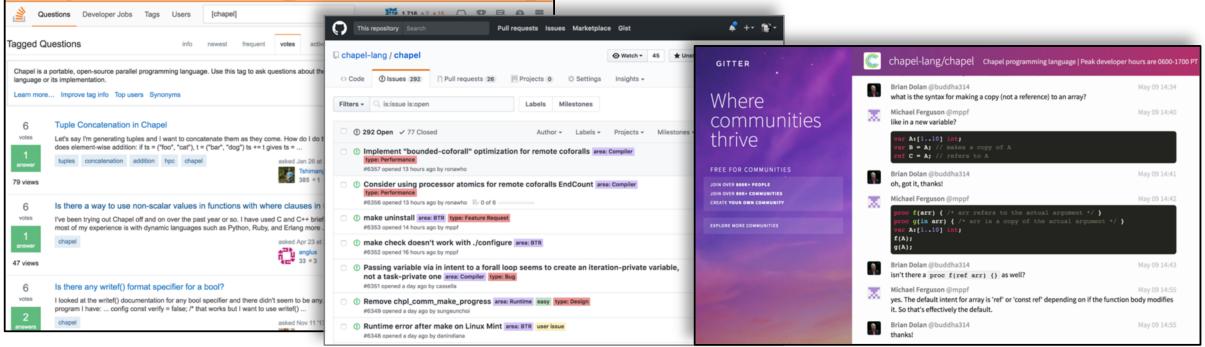
https://stackoverflow.com/questions/tagged/chapel

https://github.com/chapel-lang/chapel/issues

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https://gitter.im/chapel-lang/chapel

chapel-announce@lists.sourceforge.net



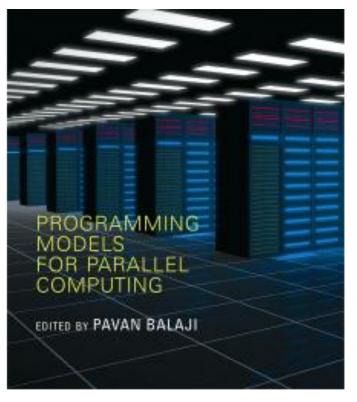


Suggested Reading (healthy attention spans)



Chapel chapter from **Programming Models for Parallel Computing**

- a detailed overview of Chapel's history, motivating themes, features
- published by MIT Press, November 2015
- edited by Pavan Balaji (Argonne)
- chapter is also available <u>online</u>



Other Chapel papers/publications available at https://chapel-lang.org/papers.html



Suggested Reading (short attention spans)



CHIUW 2017: Surveying the Chapel Landscape, Cray Blog, July 2017.

a run-down of recent events (as of 2017)

Chapel: Productive Parallel Programming, Cray Blog, May 2013.

a short-and-sweet introduction to Chapel

Six Ways to Say "Hello" in Chapel (parts 1, 2, 3), Cray Blog, Sep-Oct 2015.

a series of articles illustrating the basics of parallelism and locality in Chapel

Why Chapel? (parts 1, 2, 3), Cray Blog, Jun-Oct 2014.

 a series of articles answering common questions about why we are pursuing Chapel in spite of the inherent challenges

[Ten] Myths About Scalable Programming Languages, IEEE TCSC Blog

(index available on chapel-lang.org "blog posts" page), Apr-Nov 2012.

 a series of technical opinion pieces designed to argue against standard reasons given for not developing high-level parallel languages



Where to...



Submit bug reports:

GitHub issues for chapel-lang/chapel: public bug forum chapel_bugs@cray.com: for reporting non-public bugs

Ask User-Oriented Questions:

StackOverflow: when appropriate / other users might care Gitter (chapel-lang/chapel): community chat with archives chapel-users@lists.sourceforge.net: user discussions

Discuss Chapel development

chapel-developers@lists.sourceforge.net: developer discussions GitHub issues for chapel-lang/chapel: for feature requests, design discussions

Discuss Chapel's use in education

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chapel-education@lists.sourceforge.net: educator discussions

Directly contact Chapel team at Cray: chapel_info@cray.com



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