

The Exascale Programming Challenge and Chapel's Response

Brad Chamberlain, Chapel Team, Cray Inc. SICM² Parallel Computing Workshop

March 29th, 2014



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* time permitting

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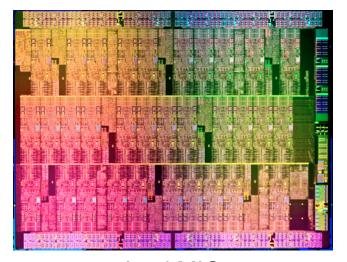
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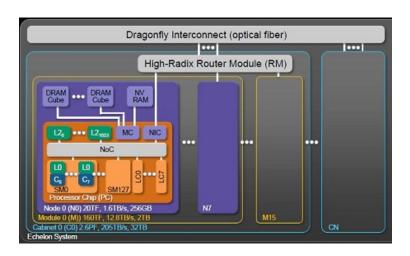
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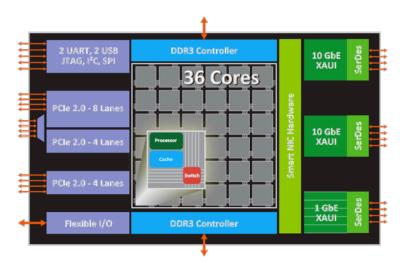
Prototypical Next-Gen Processor Technologies



Intel MIC



AMD APU



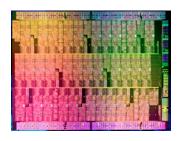
Nvidia Echelon

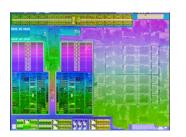
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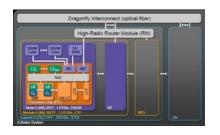


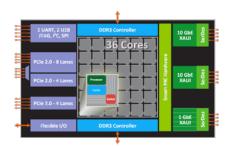
General Trends in These Architectures











- Increased hierarchy and/or sensitivity to locality
- Potentially heterogeneous processor/memory types
 - ⇒ Next-gen programmers will have a lot more to think about at the node level than in the past







Because HPC has adopted programming models that...

...have poor support for parallel work decomposition and scheduling

...have poor support for array layouts and distributed data structures

...tend to be closely tied to the architectural capabilities they target

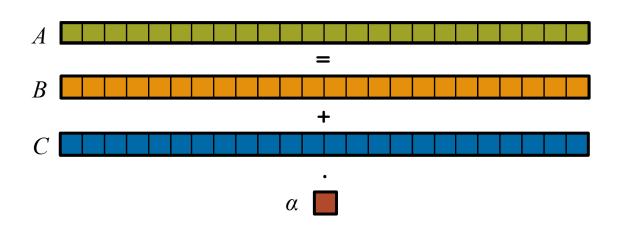


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Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m$, $A_i = B_i + \alpha \cdot C_i$

In pictures:



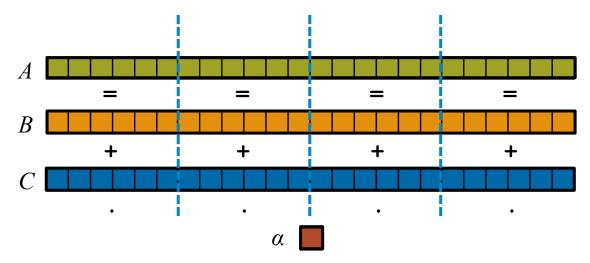


n

Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m$, $A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel:





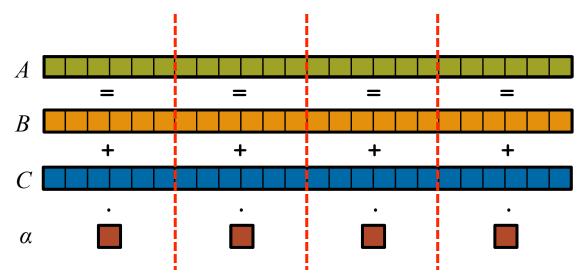
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Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m$, $A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel (distributed memory):

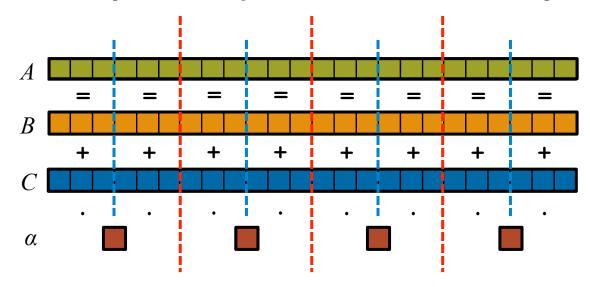




Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel (distributed memory multicore):





STREAM Triad: MPI



```
#include <hpcc.h>
static int VectorSize;
static double *a, *b, *c;
int HPCC StarStream(HPCC Params *params) {
  int myRank, commSize;
  int rv, errCount;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size ( comm, &commSize );
 MPI Comm rank ( comm, &myRank );
  rv = HPCC Stream( params, 0 == myRank);
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM,
   0, comm );
  return errCount;
int HPCC Stream(HPCC Params *params, int doIO) {
  register int j;
  double scalar;
 VectorSize = HPCC LocalVectorSize( params, 3,
   sizeof(double), 0);
  a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC( double, VectorSize );
  c = HPCC XMALLOC( double, VectorSize );
```

```
if (!a || !b || !c) {
  if (c) HPCC free(c);
  if (b) HPCC free(b);
  if (a) HPCC free(a);
  if (doIO) {
    fprintf( outFile, "Failed to allocate memory (%d).
 \n", VectorSize );
    fclose( outFile );
  return 1;
for (j=0; j<VectorSize; j++) {</pre>
 b[j] = 2.0;
  c[j] = 0.0;
scalar = 3.0;
for (j=0; j<VectorSize; j++)</pre>
```



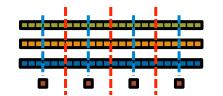
a[i] = b[j] + scalar*c[j];

HPCC free(c);

HPCC free (b);

HPCC free(a);

STREAM Triad: MPI+OpenMP





MPI + OpenMP

```
#include <hpcc.h>
                                                       if (!a || !b || !c) {
#ifdef OPENMP
                                                         if (c) HPCC free(c);
#include <omp.h>
                                                         if (b) HPCC free(b);
#endif
                                                         if (a) HPCC free(a);
static int VectorSize;
                                                         if (doI0) {
static double *a, *b, *c;
                                                            fprintf( outFile, "Failed to allocate memory (%d).
                                                         \n", VectorSize );
int HPCC StarStream(HPCC Params *params) {
                                                            fclose( outFile );
  int myRank, commSize;
  int rv, errCount;
                                                         return 1;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size ( comm, &commSize );
                                                     #ifdef OPENMP
 MPI Comm rank ( comm, &myRank );
                                                     #pragma omp parallel for
                                                      #endif
  rv = HPCC Stream( params, 0 == myRank);
                                                       for (j=0; j<VectorSize; j++) {</pre>
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM,
                                                         b[j] = 2.0;
   0, comm );
                                                         c[j] = 0.0;
  return errCount;
                                                       scalar = 3.0;
int HPCC Stream(HPCC Params *params, int doIO) {
                                                     #ifdef OPENMP
  register int j;
                                                     #pragma omp parallel for
  double scalar:
                                                      #endif
                                                       for (j=0; j<VectorSize; j++)</pre>
 VectorSize = HPCC LocalVectorSize( params, 3,
                                                         a[j] = b[j] + scalar*c[j];
   sizeof(double), 0 );
                                                       HPCC free(c);
  a = HPCC XMALLOC( double, VectorSize );
                                                       HPCC free (b);
 b = HPCC XMALLOC( double, VectorSize );
                                                       HPCC free(a);
  c = HPCC XMALLOC( double, VectorSize );
```



STREAM Triad: MPI+OpenMP vs. CUDA

MPI + OpenMP

```
#ifdef _OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Params *params) {
   int myRank, commSize;
   int rv, errCount;
   MPI_Comm comm = MPI_COMM_WORLD;

MPI_Comm_size( comm, &commSize );
   MPI_Comm_rank( comm, &myRank );

rv = HPCC_Stream( params, 0 == myRank);
   MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM, 0, comm );
```

CUDA

```
#define N 2000000
int main() {
  float *d_a, *d_b, *d_c;
  float scalar;

  cudaMalloc((void**)&d_a, sizeof(float)*N);
  cudaMalloc((void**)&d_b, sizeof(float)*N);
  cudaMalloc((void**)&d_c, sizeof(float)*N);
```

HPC suffers from too many distinct notations for expressing parallelism and locality

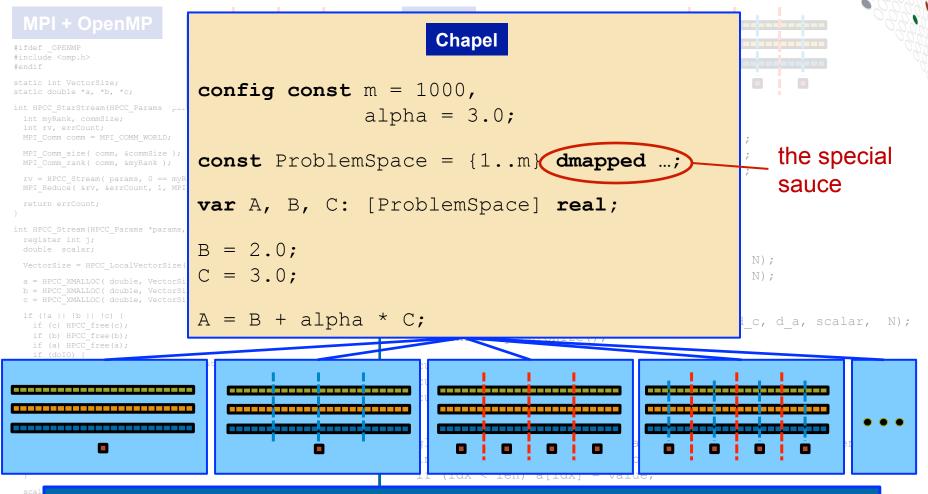
```
register int j;
 double scalar;
 VectorSize = HPCC LocalVectorSize( params, 3, sizeof(double), 0 );
 a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC ( double, VectorSize );
 c = HPCC XMALLOC( double, VectorSize );
 if (!a || !b || !c) {
   if (c) HPCC free(c);
   if (b) HPCC free(b);
   if (a) HPCC free(a);
   if (doIO) {
     fprintf( outFile, "Failed to allocate memory (%d).\n", VectorSize );
     fclose(outFile):
   return 1;
#ifdef OPENMP
#pragma omp parallel for
#endif
 for (j=0; j<VectorSize; j++) {
  b[i] = 2.0;
   c[j] = 0.0;
 scalar = 3.0;
#ifdef OPENMP
#pragma omp parallel for
#endif
 for (j=0; j<VectorSize; j++)
   a[j] = b[j] + scalar*c[j];
 HPCC free(c);
 HPCC free (b);
```



HPCC free(a);

STREAM Triad: Chapel

#pragr



<u>Philosophy:</u> Good language design can tease details of locality and parallelism away from an algorithm, permitting the compiler, runtime, applied scientist, and HPC expert to each focus on their strengths.

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1288 lines of source code

plus 266 lines of comments

487 blank lines

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(the corresponding C+MPI+OpenMP version is nearly 4x bigger)



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This is trunk/test/release/examples/benchmarks/lulesh/*.chpl in the

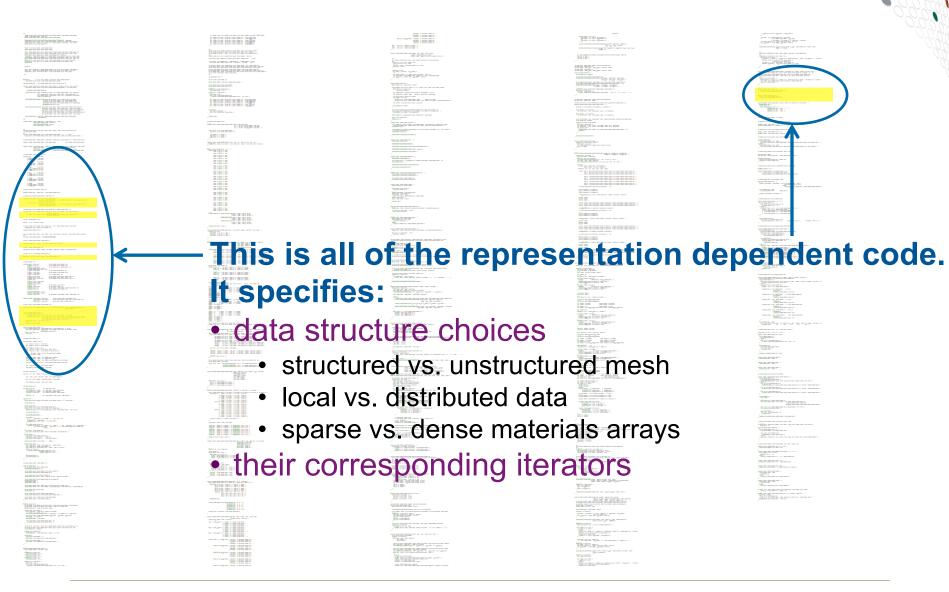
SourceForge repository, as of r22745 (2/16/14).



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LULESH in Chapel





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Why so many programming models?



HPC has traditionally given users...

- ...low-level, control-centric programming models
- ...ones that are closely tied to the underlying hardware
- ...ones that support only a single type of parallelism

Type of HW Parallelism	Programming Model	Unit of Parallelism
Inter-node	MPI	executable
Intra-node/multicore	OpenMP/pthreads	iteration/task
Instruction-level vectors/threads	pragmas	iteration
GPU/accelerator	CUDA/OpenCL/OpenACC	SIMD function/task

benefits: lots of control; decent generality; easy to implement downsides: lots of user-managed detail; brittle to changes



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What is Chapel?



- An emerging parallel programming language
 - Design and development led by Cray Inc.
 - in collaboration with academia, labs, industry
 - version 1.8 had19 contributors from 8 organizations and 5 countries
 - Initiated under the DARPA HPCS program
- Being developed as open (BSD) software at SourceForge

A work-in-progress



Chapel's Targets



- Target Architectures:
 - multicore desktops and laptops
 - commodity clusters and the cloud
 - HPC systems from Cray and other vendors
 - *in-progress:* exascale-era architectures
- Chapel's overall goal: Improve programmer productivity



What does "Productivity" mean to you?



Recent Graduate:

"something similar to what I used in school: Python, Matlab, Java, ..."

Seasoned HPC Programmer:

want full control/performance

"that sugary stuff which I don't need because I was born to suffer"

Computational Scientist:

"something that lets me focus on my parallel computational algorithms without having to wrestle with architecture-specific details"

Chapel Team:

"something that lets the computational scientist express what they want, without taking away the control an HPC programmer would want, implemented in a language as attractive as recent graduates want."



Three Chapel Successes



Effectively separating algorithms from system mappings

- user-defined array layouts and distributions
 - alg: "I'd like an array of this type over this index set"
 - map: "how should this array be distributed? stored locally?"
- user-defined parallel iterators
 - alg: "forall ...", whole-array operations, reductions, ...
 - map: "how many tasks? how to divide the iterations?"
- seamless integration of data and task parallelism

Distinct concepts for parallelism and locality*

- "SPMD-only" and "shared memory-only" are restrictive to begin with
- I believe they're non-starters in an exascale world

Withstanding the Naysayers

 we've generated cautious optimism in a community that's never had a productive language; and that has seen many, many failed attempts

(* I don't mean to suggest that Chapel was the first to do this—we weren't—simply that I believe it to be so crucial as to deserve silver)



Three Chapel Challenges



Performance

 the downside of permitting so much to be user-defined is that there's a bigger gap to close compared to the status quo

Reaching a Tipping Point in Acceptance/Utilization

- Chapel has lots of wallflower fans—how to get them invested?
- and when?

Rapidly Responding to Emerging Architectures

- Chapel is designed to be forward portable, but effort is still required
- ability to respond quickly would increase attractiveness



How Can Scientists Help?



- Secure Time/Resources for Studying and Evaluating Promising Emerging Technologies
 - no need to be more comprehensive than you desire
 - but if you don't like the status quo, invest some time in an alternative
- Communicate your wishlists to new languages like Chapel
 - "protect me from architectural changes" is a reasonable one
 - but surely you've got others?
- Be patient
 - no truly productive HPC-ready language is going to show up overnight without warning
- Do something more constructive than stating the obvious
 - yes, adoption of new languages is a difficult challenge
 - do you want to be part of the grumbling crowd, or part of the solution?

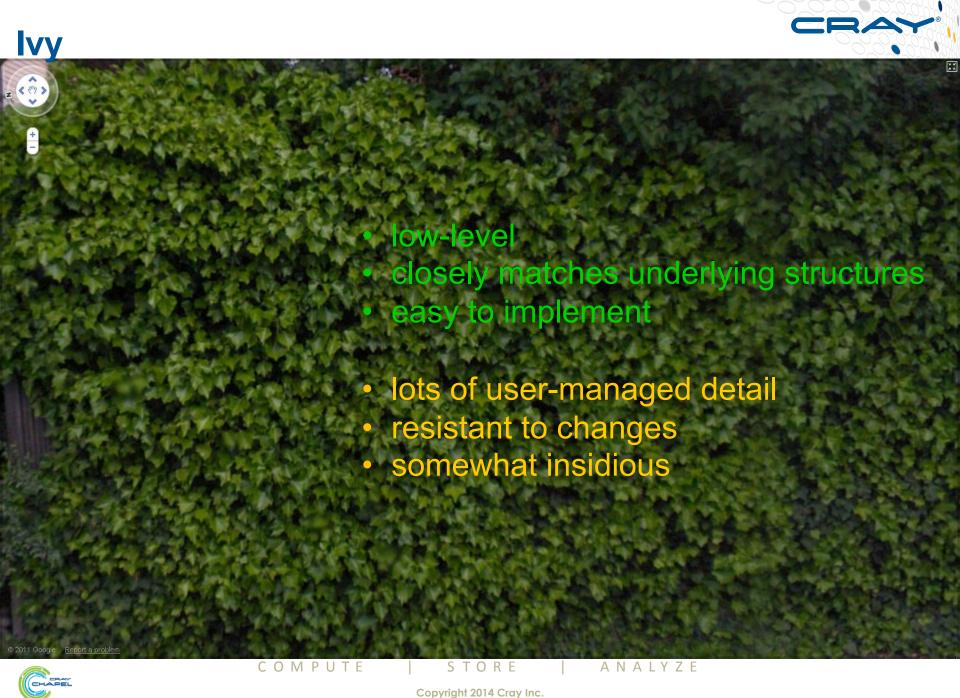


A Seattle Corner









Trees







Landscaping Quotes from the HPC community



Early HPCS years:

"The HPC community tried to plant a tree once. It didn't survive. Nobody should ever bother planting one again."

"Why plant a tree if you can't be assured it will thrive?"

"Why would anyone ever want anything other than ivy?"

"We're in the business of building treehouses that last 40 years; we can't afford to build one in the branches of your sapling."



Landscaping Quotes from the HPC community



Early HPCS years (for the analogy-challenged):

"The HPC community tried to develop a HLL once. It didn't survive. Nobody should ever bother developing one again."

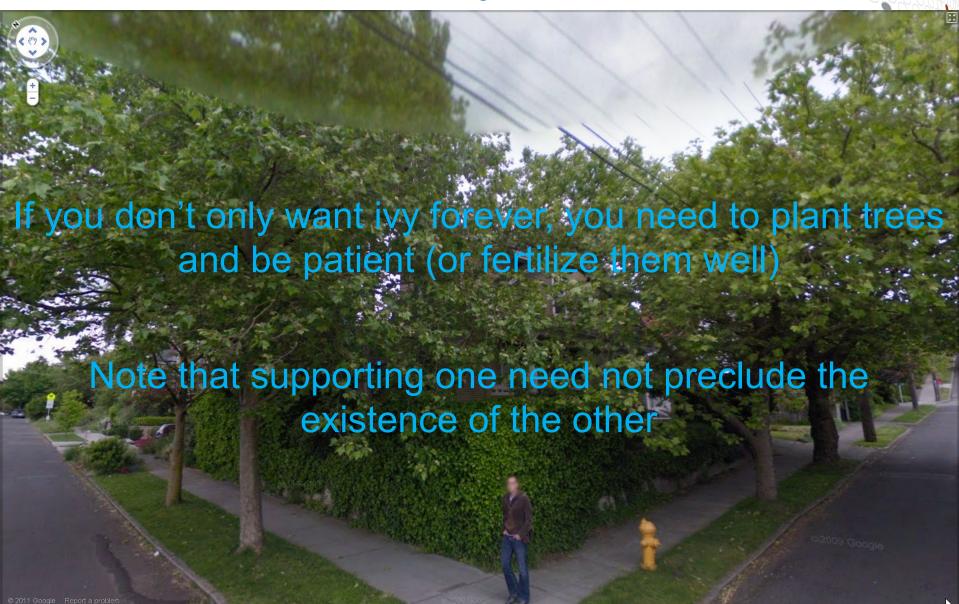
"Why develop a language you can't be assured it will thrive?"

"Why would anyone ever want anything other than MPI+X?"

"We're in the business of writing applications that last 40 years; we can't afford to risk writing one in an emerging language."



A Corner in Seattle: Takeaways





Challenges for Computer Scientists



What are the abstractions that...

- give the application scientists the abstractions they want?
 - could be realized as DSLs, APIs, ADTs, ...
- support mapping down to multiple implementation choices
 - e.g., MPI+X as a safety net; Chapel as an investment in a better future

• How do we collaborate effectively?

- there aren't many parallel language folks, and we each have our own
- lone wolf researcher is seductive: independent, full control, full credit
 - but, we didn't reach the moon via dozens of partially-built rockets



A Note on Interoperability



 If your language only supports one array format, it's only going to be efficiently interoperable with a small set of languages

• Via its user-defined array distributions and layouts, Chapel enables universal *in situ* interoperation



A Note on Parallel Education



When teaching parallel programming, I like to cover:

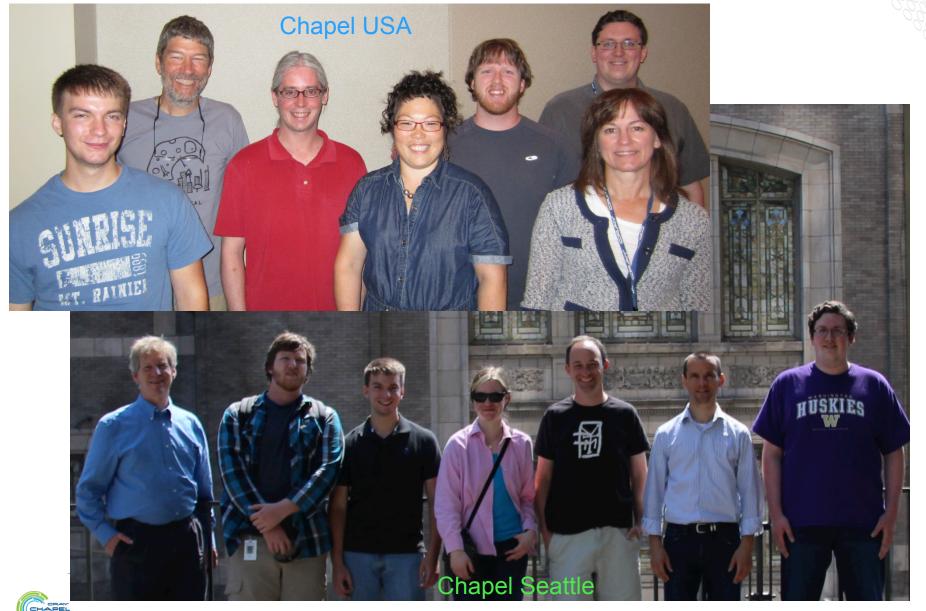
- data parallelism
- task parallelism
- concurrency
- synchronization
- locality/affinity
- deadlock, livelock, and other pitfalls
- performance tuning
- ...

I don't think there's been a good language out there...

- for teaching all of these things
- for teaching some of these things well at all
- until now: We believe Chapel can fill a crucial gap here
 (see http://chapel.cray.com/education.html for more information and http://cs.washington.edu/education/courses/csep524/13wi/ for my use of Chapel in class)



The Cray Chapel Team (Summer 2013)



Chapel is a collaborative effort... join us!





















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For More Information: Online Resources



Chapel project page: http://chapel.cray.com

overview, papers, presentations, language spec, ...

Chapel SourceForge page: https://sourceforge.net/projects/chapel/

• release downloads, public mailing lists, code repository, ...

Mailing Aliases:

- chapel_info@cray.com: contact the team at Cray
- chapel-announce@lists.sourceforge.net: announcement list
- chapel-users@lists.sourceforge.net: user-oriented discussion list
- chapel-developers@lists.sourceforge.net: developer discussion
- chapel-education@lists.sourceforge.net: educator discussion
- chapel-bugs@lists.sourceforge.net: public bug forum

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For More Information: Suggested Reading



Overview Papers:

- <u>The State of the Chapel Union</u> [slides], Chamberlain, Choi, Dumler, Hildebrandt, Iten, Litvinov, Titus. CUG 2013, May 2013.
 - a high-level overview of the project summarizing the HPCS period
- <u>A Brief Overview of Chapel</u>, Chamberlain (pre-print of a chapter for A Brief Overview of Parallel Programming Models, edited by Pavan Balaji, to be published by MIT Press in 2014).
 - a more detailed overview of Chapel's history, motivating themes, features

Blog Articles:

- [Ten] Myths About Scalable Programming Languages, Chamberlain. IEEE Technical Committee on Scalable Computing (TCSC) Blog, (https://www.ieeetcsc.org/activities/blog/), April-November 2012.
 - a series of technical opinion pieces designed to rebut standard arguments against the development of high-level parallel languages



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