

Chapel: Productive Parallel Programming at Scale

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Northwest C++ Users Group
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Cray: The Supercomputing Company

1972: Seymour Cray founded Cray Research

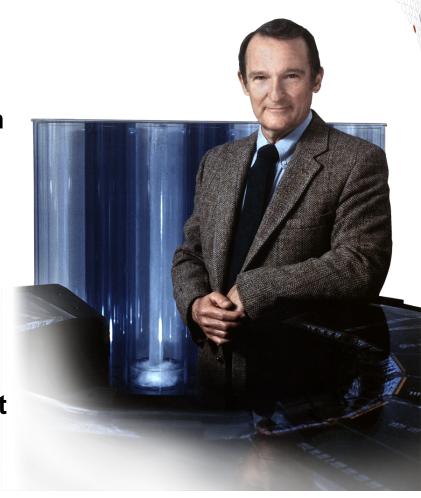
2000: Tera purchased Cray Research from SGI and formed Cray Inc.

corporate headquarters based in Seattle, WA

Technology Focus Areas:

- Computation
- Storage
- Analytics

Vision: Provide the systems and tools that our customers need to solve the world's hardest problems.







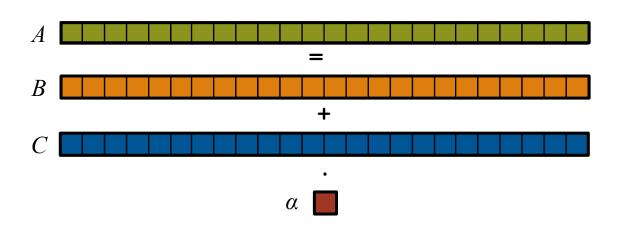
High Performance Computing (HPC) Programming Models by Example



Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

In pictures:



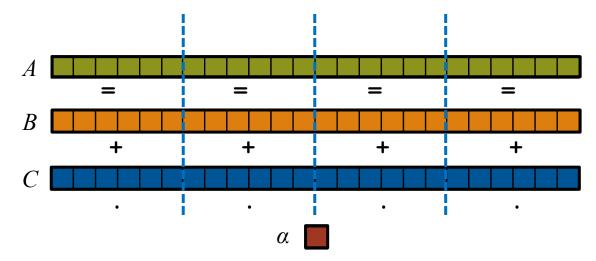


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Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel:

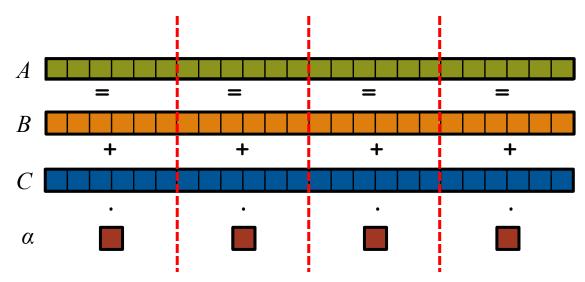




Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel (distributed memory):

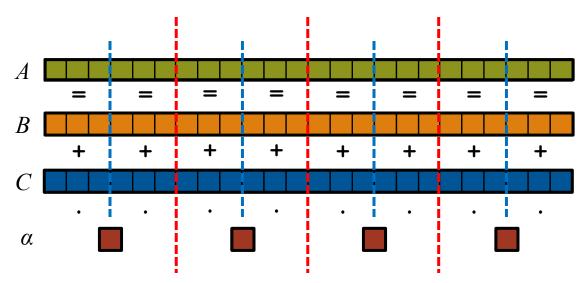




Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel (distributed memory multicore):





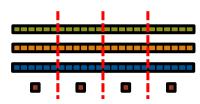
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STREAM Triad: MPI



```
#include <hpcc.h>
```

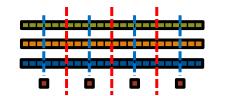
```
static int VectorSize;
static double *a, *b, *c;
int HPCC StarStream(HPCC Params *params) {
 int myRank, commSize;
 int rv, errCount;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size( comm, &commSize );
 MPI Comm rank( comm, &myRank );
 rv = HPCC Stream( params, 0 == myRank);
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM,
   0, comm);
 return errCount;
int HPCC Stream(HPCC Params *params, int doIO) {
 register int j;
 double scalar;
 VectorSize = HPCC LocalVectorSize( params, 3,
   sizeof(double), 0);
  a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC( double, VectorSize );
                                                      HPCC free(a);
 c = HPCC XMALLOC( double, VectorSize );
```



```
if (!a || !b || !c) {
  if (c) HPCC free(c);
  if (b) HPCC free(b);
  if (a) HPCC free(a);
  if (doIO) {
    fprintf( outFile, "Failed to allocate memory
  (%d).\n", VectorSize);
    fclose( outFile );
  return 1;
for (j=0; j<VectorSize; j++) {</pre>
 b[\dot{j}] = 2.0;
  c[j] = 0.0;
scalar = 3.0;
for (j=0; j<VectorSize; j++)</pre>
  a[j] = b[j] + scalar*c[j];
HPCC free(c);
HPCC free (b);
```



STREAM Triad: MPI+OpenMP





MPI + OpenMP

```
#include <hpcc.h>
                                                        if (!a || !b || !c) {
#ifdef OPENMP
                                                          if (c) HPCC free(c);
#include <omp.h>
                                                          if (b) HPCC free(b);
#endif
                                                          if (a) HPCC free(a);
                                                          if (doIO) {
static int VectorSize;
static double *a, *b, *c;
                                                            fprintf( outFile, "Failed to allocate memory
                                                          (%d).\n", VectorSize);
int HPCC StarStream(HPCC Params *params) {
                                                            fclose( outFile );
  int myRank, commSize;
  int rv, errCount;
                                                          return 1;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size( comm, &commSize );
                                                      #ifdef OPENMP
 MPI Comm rank( comm, &myRank );
                                                      #pragma omp parallel for
                                                      #endif
  rv = HPCC Stream( params, 0 == myRank);
                                                        for (j=0; j<VectorSize; j++) {</pre>
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM,
                                                         b[\dot{j}] = 2.0;
   0, comm);
                                                          c[j] = 0.0;
  return errCount;
                                                        scalar = 3.0;
int HPCC Stream(HPCC Params *params, int doIO) {
                                                      #ifdef OPENMP
  register int j;
                                                      #pragma omp parallel for
  double scalar;
                                                      #endif
                                                        for (j=0; j<VectorSize; j++)</pre>
 VectorSize = HPCC LocalVectorSize( params, 3,
                                                          a[j] = b[j] + scalar*c[j];
   sizeof(double), 0);
                                                        HPCC free(c);
  a = HPCC XMALLOC( double, VectorSize );
                                                        HPCC free(b);
 b = HPCC XMALLOC( double, VectorSize );
                                                        HPCC free(a);
  c = HPCC XMALLOC( double, VectorSize );
```



STREAM Triad: MPI+OpenMP vs. CUDA

MPI + OpenMP #ifdef OPENMP #include <omp.h> #endif static int VectorSize; static double *a, *b, *c;

```
int HPCC StarStream(HPCC Params *params) {
 int myRank, commSize;
 int rv, errCount;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size ( comm, &commSize );
 MPI Comm rank ( comm, &myRank );
 rv = HPCC Stream( params, 0 == myRank);
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM, 0, comm );
```

return errCount:

```
VectorSize = HPCC LocalVectorSize( params, 3, sizeof(double), 0 );
 a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC( double, VectorSize );
 c = HPCC XMALLOC( double, VectorSize );
 if (!a || !b || !c) {
   if (c) HPCC free(c);
   if (b) HPCC free(b);
   if (a) HPCC_free(a);
   if (doIO) {
     fprintf( outFile, "Failed to allocate memory (%d).\n", VectorSize );
     fclose( outFile );
   return 1;
#ifdef OPENMP
#pragma omp parallel for
#endif
 for (j=0; j<VectorSize; j++) {
   b[j] = 2.0;
   c[j] = 0.0;
 scalar = 3.0;
#ifdef OPENMP
#pragma omp parallel for
#endif
 for (j=0; j<VectorSize; j++)</pre>
   a[j] = b[j] + scalar*c[j];
 HPCC free(c);
 HPCC free (b);
 HPCC free(a):
 return 0:
```

CUDA

dim3 dimBlock (128)

```
#define N
                2000000
int main() {
  float *d a, *d b, *d c;
  float scalar;
  cudaMalloc((void**)&d a, sizeof(float)*N);
  cudaMalloc((void**)&d b, sizeof(float)*N);
 cudaMalloc((void**)&d c, sizeof(float)*N);
```

HPC suffers from too many distinct notations for expressing parallelism and locality

```
set array<<<dimGrid,dimBlock>>>(d b, .5f, N);
 set array<<<dimGrid,dimBlock>>>(d c, .5f, N);
 scalar=3.0f;
 STREAM Triad<<<dimGrid,dimBlock>>>(d b, d_c, d_a, scalar, N);
 cudaThreadSvnchronize();
 cudaFree(d a);
 cudaFree(d b);
 cudaFree(d c);
 global void set array(float *a, float value, int len) {
 int idx = threadIdx.x + blockIdx.x * blockDim.x;
 if (idx < len) a[idx] = value;
global void STREAM Triad( float *a, float *b, float *c,
                             float scalar, int len) {
 int idx = threadIdx.x + blockIdx.x * blockDim.x;
 if (idx < len) c[idx] = a[idx]+scalar*b[idx];</pre>
```



Why so many programming models?



HPC tends to approach programming models bottom-up:

Given a system and its core capabilities...

- ...provide features that can access the available performance.
 - portability? generality? programmability? ...not strictly required.

Type of HW Parallelism	Programming Model	Unit of Parallelism
Inter-node	MPI	executable
Intra-node/multicore	OpenMP / pthreads	iteration/task
Instruction-level vectors/threads	pragmas	iteration
GPU/accelerator	CUDA / Open[CL MP ACC]	SIMD function/task

benefits: lots of control; decent generality; easy to implement downsides: lots of user-managed detail; brittle to changes



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Motivation for Chapel



Q: Can a single language be...

- ...as productive as Python?
- ...as fast as Fortran?
- ...as portable as C?
- ...as scalable as MPI?
- ...as fun as <your favorite language here>?

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A: We believe so.



The Challenge



Q: So why don't we have such languages already?

A: Technical challenges?

while they exist, we don't think this is the main issue...

A: Due to a lack of...

- ...long-term efforts
- ...resources
- ...community will
- ...co-design between developers and users
- ...patience

Chapel is our attempt to reverse this trend



What is Chapel?



Chapel: A productive parallel programming language

- extensible
- portable
- open-source
- a collaborative effort
- a work-in-progress

Goals:

- Support general parallel programming
 - "any parallel algorithm on any parallel hardware"
- Make parallel programming far more productive



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What does "Productivity" mean to you?



Recent Graduates:

"something similar to what I used in school: Python, Matlab, Java, ..."

Seasoned HPC Programmers:

"that sugary stuff that I don't need because I was born to suffer"
want full control
to ensure performance"

Computational Scientists:

"something that lets me express my parallel computations without having to wrestle with architecture-specific details"

Chapel Team:

"something that lets computational scientists express what they want, without taking away the control that HPC programmers need, implemented in a language as attractive as recent graduates want."



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Rewinding a few slides...

MPI + OpenMP

```
#ifdef _OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;
int HPCC_StarStream(HPCC_Params *params) {
   int myRank, commSize;
   int rv, errCount;

MPI_Comm comm = MPI_COMM_WORLD;

MPI_Comm size( comm, $commSize );
MPI_Comm_rank( comm, $myRank );

rv = HPCC_Stream( params, 0 == myRank);
MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM, 0, comm );
return errCount;
```

```
#define N 2000000
int main() {
  float *d a, *d b, *d c;
```

```
cudaMalloc((void**)&d_a, sizeof(float)*N);
cudaMalloc((void**)&d_b, sizeof(float)*N);
cudaMalloc((void**)&d_c, sizeof(float)*N);
```

HPC suffers from too many distinct notations for expressing parallelism and locality

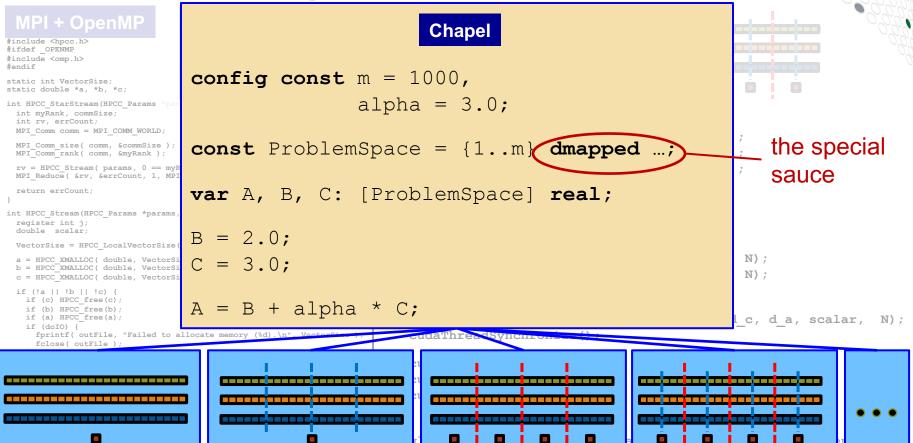
dim3 dimBlock(128)

float scalar;

```
VectorSize = HPCC LocalVectorSize( params, 3, sizeof(double), 0 );
 a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC( double, VectorSize );
 c = HPCC XMALLOC( double, VectorSize );
 if (!a || !b || !c) {
   if (c) HPCC free(c);
   if (b) HPCC free(b);
   if (a) HPCC_free(a);
   if (doIO) {
     fprintf( outFile, "Failed to allocate memory (%d).\n", VectorSize );
     fclose( outFile );
   return 1;
#ifdef OPENMP
#pragma omp parallel for
#endif
 for (j=0; j<VectorSize; j++) {
   b[j] = 2.0;
   c[j] = 0.0;
 scalar = 3.0;
#ifdef OPENMP
#pragma omp parallel for
#endif
 for (j=0; j<VectorSize; j++)</pre>
   a[j] = b[j] + scalar*c[j];
 HPCC free(c);
 HPCC free (b);
 HPCC free(a):
 return 0:
```



STREAM Triad: Chapel



<u>Philosophy:</u> Good, *top-down* language design can tease system-specific implementation details away from an algorithm, permitting the compiler, runtime, applied scientist, and HPC expert to each focus on their strengths.



Outline

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- ✓ Motivation for Chapel
- Survey of Chapel Concepts
- Chapel Project and Characterizations
- Chapel Resources



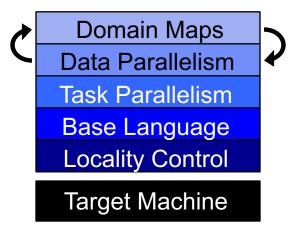
Chapel's Multiresolution Philosophy



Multiresolution Design: Support multiple tiers of features

- higher levels for programmability, productivity
- lower levels for greater degrees of control

Chapel language concepts



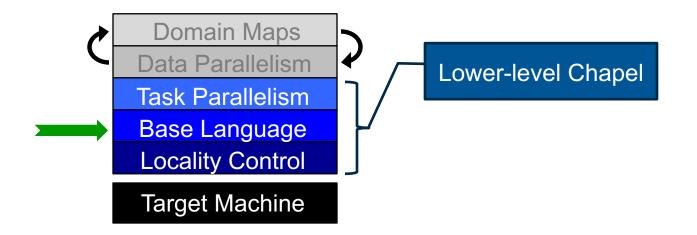
- build the higher-level concepts in terms of the lower
- permit the user to intermix layers arbitrarily



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Base Language







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```
iter fib(n) {
 var current = 0,
      next = 1;
  for i in 1...n {
    yield current;
    current += next;
    current <=> next;
```

```
config const n = 10;
for f in fib(n) do
 writeln(f);
```

```
3
5
8
```





CLU-style iterators

```
var current = 0,
    next = 1;

for i in 1..n {
    yield current;
    current += next;
    current <=> next;
}
```

```
config const n = 10;
for f in fib(n) do
  writeln(f);
```

```
0
1
1
2
3
5
8
```



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Configuration declarations
(to avoid command-line argument parsing)
./a.out -n=1000000

```
iter fib(n) {
  var current = 0,
    next = 1;

for i in 1..n {
    yield current;
    current += next;
    current <=> next;
}
```

```
config const n = 10;

for f in fib(n) do
  writeln(f);
```

```
0
1
1
2
3
5
8
```



```
Static type inference for:
                 arguments
                 return types
                variables
                            config const n = 10;
iter fib(n')
  var current' = 0,
                            for f'in fib(n) do
      next = 1;
                              writeln(f);
  for i in 1..n {
    yield current;
    current += next;
    current <=> next;
                                   3
                                   5
                                   8
```



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Zippered iteration

```
iter fib(n) {
 var current = 0,
      next = 1;
  for i in 1..n {
    yield current;
    current += next;
    current <=> next;
```

```
config const n = 10;
for (i,f) in zip(0..#n, fib(n)) do
 writeln("fib #", i, " is ", f);
```

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
```



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```
iter fib(n) {
  var current = 0
    next = 1;

for i in 1..n {
    yield current;
    current += next;
    current <=> next;
}
```

```
Range types and operators
```

```
config const n = 10;
for (i,f) in zip(0..#n, fib(n)) do
  writeln("fib #", i, " is ", f);
```

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...
```



```
iter fib(n) {
  var current = 0,
    next = 1;

for i in 1..n {
    yield current;
    current += next;
    current <=> next;
}
```

```
config const n = 10;
for (i,f) in zip(0..#n, fib(n)) do
  writeln("fib #", i, " is ", f);
```

tuples

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...
```



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```
iter fib(n) {
  var current = 0,
    next = 1;

for i in 1..n {
    yield current;
    current += next;
    current <=> next;
}
```

```
config const n = 10;

for (i,f) in zip(0..#n, fib(n)) do
  writeln("fib #", i, " is ", f);
```

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...
```



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Other Base Language Features

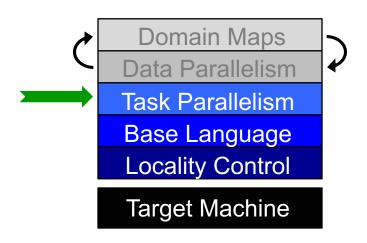


- interoperability features
- OOP (value- and reference-based)
- overloading, where clauses
- argument intents, default values, match-by-name
- compile-time features for meta-programming
 - e.g., compile-time functions to compute types, values; reflection
- modules (for namespace management)
- rank-independent programming features



Task Parallelism







Task Parallelism: Begin Statements



```
// create a fire-and-forget task for a statement
begin writeln("hello world");
writeln("goodbye");
```

Possible outputs:

hello world goodbye

goodbye hello world



Task Parallelism: Coforall Loops



```
// create a task per iteration
coforall t in 0..#numTasks {
  writeln("Hello from task ", t, " of ", numTasks);
} // implicit join of the numTasks tasks here
writeln("All tasks done");
```

Sample output:

```
Hello from task 2 of 4
Hello from task 0 of 4
Hello from task 3 of 4
Hello from task 1 of 4
All tasks done
```



Task Parallelism: Data-Driven Synchronization

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- atomic variables: support atomic operations
 - e.g., compare-and-swap; atomic sum, multiply, etc.
 - similar to C/C++
- sync variables: store full-empty state along with value
 - by default, reads/writes block until full/empty, leave in opposite state



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Bounded Buffer Producer/Consumer Example

```
begin producer();
consumer();
// 'sync' types store full/empty state along with value
var buff$: [0..#buffersize] sync real;
proc producer() {
  var i = 0;
  for ... {
    i = (i+1) % buffersize;
   buff$[i] = ...; // writes block until empty, leave full
proc consumer() {
  var i = 0;
  while ... {
    i= (i+1) % buffersize;
    ...buff$[i]...; // reads block until full, leave empty
```



Other Task Parallel Concepts

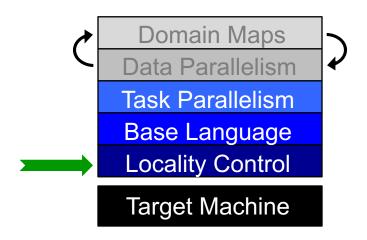


- cobegins: create tasks using compound statements
- single variables: like sync variables, but write-once
- sync statements: join unstructured tasks
- serial statements: conditionally squash parallelism



Locality Control







The Locale Type



Definition:

- Abstract unit of target architecture
- Supports reasoning about locality
 - defines "here vs. there" / "local vs. remote"
- Capable of running tasks and storing variables
 - i.e., has processors and memory

Typically: A compute node (multicore processor or SMP)



Getting started with locales



Specify # of locales when running Chapel programs

```
% a.out --numLocales=8
```

Chapel provides built-in locale variables

main() starts execution as a task on locale #0



Locale Operations



Locale methods support queries about the target system:

```
proc locale.physicalMemory(...) { ... }
proc locale.numCores { ... }
proc locale.id { ... }
proc locale.name { ... }
```

On-clauses support placement of computations:

```
writeln("on locale 0");

on Locales[1] do
  writeln("now on locale 1");

writeln("on locale 0 again");
```

```
on A[i,j] do
  bigComputation(A);
on node.left do
  search(node.left);
```



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Parallelism and Locality: Orthogonal in Chapel



```
coforall i in 1..msgs do
  writeln("Hello from task ", i);
```

This is a distributed, but serial program:

```
writeln ("Hello from locale 0!");
on Locales[1] do writeln("Hello from locale 1!");
on Locales[2] do writeln ("Hello from locale 2!");
```

This is a distributed parallel program:

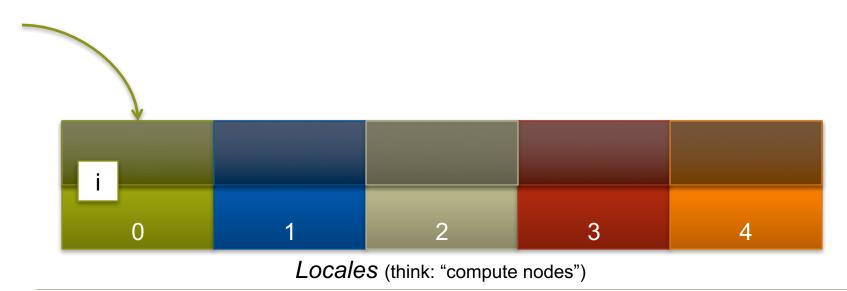
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```
coforall i in 1..msgs do
  on Locales[i%numLocales] do
   writeln ("Hello from task ", i,
            " running on locale ", here.id);
```



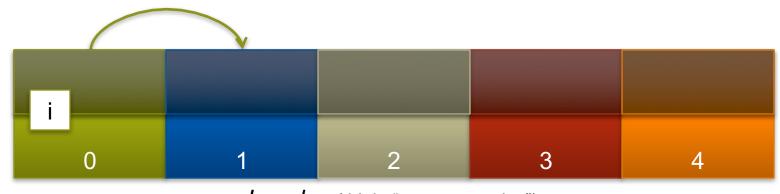


var i: int;





```
var i: int;
on Locales[1] {
```



Locales (think: "compute nodes")



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```
var i: int;
on Locales[1] {
  var j: int;
```

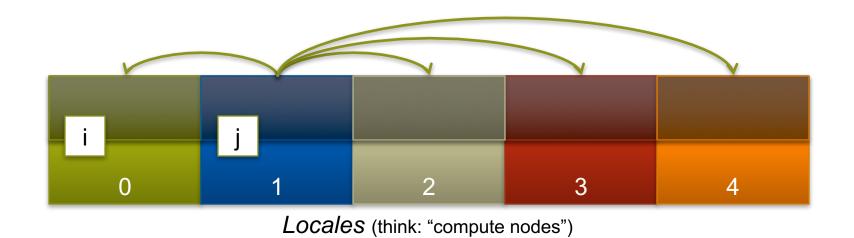


Locales (think: "compute nodes")



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```
var i: int;
on Locales[1] {
  var j: int;
  coforall loc in Locales {
    on loc {
```





COMPUTE



```
var i: int;
on Locales[1] {
  var j: int;
  coforall loc in Locales {
    on loc {
       var k: int;
                                             k
                                                         k
                                          3
                    Locales (think: "compute nodes")
```



```
var i: int;
on Locales[1] {
  var j: int;
  coforall loc in Locales {
     on loc {
       var k: int;
       k = 2*i + j;
          OK to access i, j, and k
                                        = 2*i +
            wherever they live
                                                            k
                     Locales (think: "compute nodes")
```



```
var i: int;
on Locales[1] {
  var j: int;
  coforall loc in Locales {
     on loc {
       var k: int;
       k = 2*i + j;
        here, i and j are remote, so
        the compiler + runtime will
                                            2*i +
           transfer their values
                                                 k
            k
                                                              k
                                           (j)
                      Locales (think: "compute nodes")
```



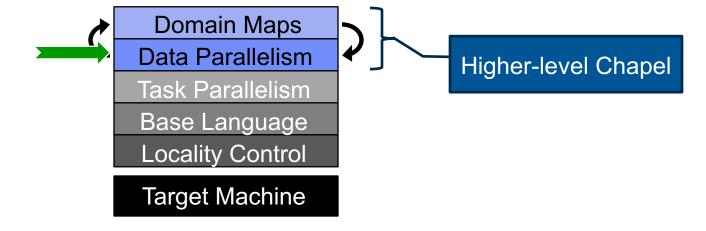
Chapel: Locality queries

```
var i: int;
on Locales[1] {
  var j: int;
  coforall loc in Locales {
     on loc {
       var k: int;
                 // query the locale on which this task is running
       ...here...
       ...j.locale... // query the locale on which j is stored
                                                k
                                                             k
                     Locales (think: "compute nodes")
```



Data Parallelism







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Data Parallelism By Example: STREAM Triad





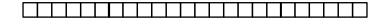
var A, B, C: [ProblemSpace] real;



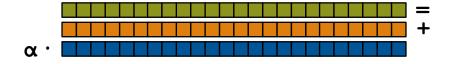


Data Parallelism By Example: STREAM Triad





var A, B, C: [ProblemSpace] real;



A = B + alpha * C; // equivalent to the zippered forall



Other Data Parallel Features



Rich Domain/Array Types:

- multidimensional
- strided
- sparse
- associative
- Slicing: Refer to subarrays using ranges/domains

```
... A[2..n-1, lo..#b] ...
... A[ElementsOfInterest] ...
```

Promotion: Call scalar functions with array arguments

```
... pow(A, B)... // equivalent to: forall (a,b) in zip(A,B) do pow(a,b)
```

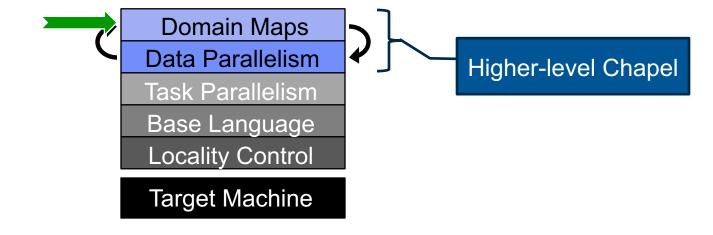
Reductions/Scans: Apply operations across collections

```
... + reduce A ...
... myReduceOp reduce A ...
```



Domain Maps







STREAM Triad: Chapel (multicore)



COMPUTE

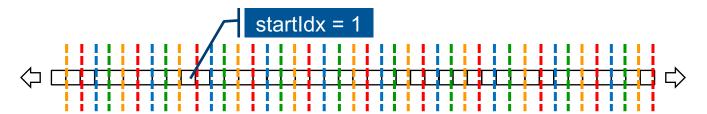
No domain map specified \Rightarrow use default layout

- current locale owns all domain indices and array values
- computation will execute using local processors only



STORE

STREAM Triad: Chapel (multilocale, cyclic)



```
const ProblemSpace = {1..m}
```

dmapped Cyclic(startIdx=1);



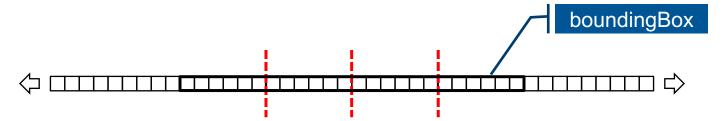
var A, B, C: [ProblemSpace] real;



$$A = B + alpha * C;$$

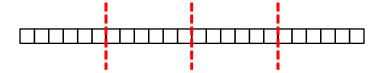


STREAM Triad: Chapel (multilocale, blocked)

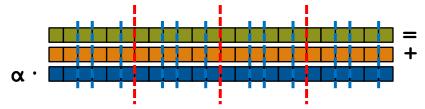


const ProblemSpace = {1..m}

dmapped Block(boundingBox={1..m});



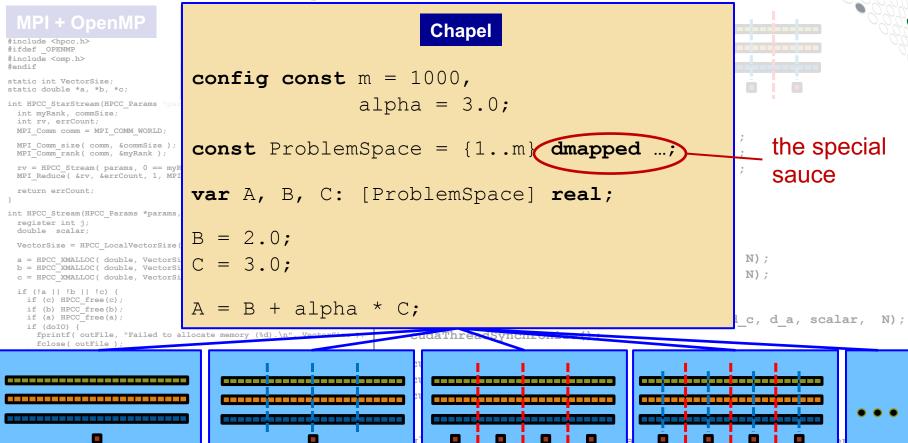
var A, B, C: [ProblemSpace] real;



A = B + alpha * C;



STREAM Triad: Chapel



<u>Philosophy:</u> Good, *top-down* language design can tease system-specific implementation details away from an algorithm, permitting the compiler, runtime, applied scientist, and HPC expert to each focus on their strengths.



Chapel is Extensible



Advanced users can create their own...

- ...domain maps (array layouts and distributions)...
- ...parallel loop schedules...
- ...models of the target architecture...

...as Chapel code, without modifying the compiler.

Why? To create a future-proof language.

This has been our main R&D challenge: How to create a language that does not lock these policies into the implementation without sacrificing performance?



Language Summary



Parallel programmers deserve better programming models

Higher-level programming models can help insulate algorithms from parallel implementation details

- yet, without necessarily abdicating control
- Chapel does this via its multiresolution design

We believe Chapel can greatly improve productivity

- ...for current and emerging parallel architectures
- ...for HPC users as well as mainstream uses of parallelism



Outline

CRAY

- ✓ Motivation for Chapel
- ✓ Survey of Chapel Concepts
- Chapel Project and Characterizations
- Chapel Resources



Chapel is Portable



Chapel is designed to be hardware-independent

The current release requires:

- a C/C++ compiler
- a *NIX environment (Linux, OS X, BSD, Cygwin, ...)
- POSIX threads
- RDMA, MPI, or UDP (for distributed memory execution)

Chapel can run on...

- ...laptops and workstations
- ...commodity clusters
- ...the cloud
- ...HPC systems from Cray and other vendors
- ...modern processors like Intel Xeon Phi, GPUs*, etc.

* = academic work only; not yet supported in the official release



ORE I ANALYZE

Chapel is Open-Source



- Chapel's development is hosted at GitHub
 - https://github.com/chapel-lang
- Chapel is licensed as Apache v2.0 software
- Instructions for download + install are online

COMPUTE

see http://chapel.cray.com/download.html to get started



The Chapel Team at Cray (Summer 2016)







Chapel is a Collaborative Effort





















National Laboratory





Sandia National Laboratories



(and several others...)

http://chapel.cray.com/collaborations.html



I ANALYZE

Chapel is a Work-in-Progress



- Currently being picked up by early adopters
 - over two releases, 3000+ downloads per year
 - Users who try it generally like what they see



A notable early adopter



Chapel in the (Cosmological) Wild

1:00 - 2:00

Nikhil Padmanabhan, Yale University Professor, Physics & Astronomy

Abstract: This talk aims to present my personal experiences using Chapel in my research. My research interests are in observational cosmology; more specifically, I use large surveys of galaxies to constrain the evolution of the



Search





Chapel Parallel Programming Language

Videos

Playlists

Channels



CHIUW 2016 keynote: "Chapel in the (Cosmological) Wild", Nikhil Padmanabhan

Chapel Parallel Programming Language 1 month ago • 86 views

This is Nikhil Padmanabhan's keynote talk from CHIUW 2016: the 3rd Annual Chapel Implementers and Users workshop. The slides are availabl...



ORE I ANALYZE

Chapel is a Work-in-Progress



- Currently being picked up by early adopters
 - Last two releases got ~3500 downloads total in a year
 - Users who try it generally like what they see
- Most current features are functional and working well
 - some areas need improvements, e.g., error-handling, constructors
- Performance varies, but is continually improving
 - shared memory performance is typically competitive with C+OpenMP
 - distributed memory performance tends to be more hit-and-miss
- We are actively working to address these lacks



Outline

CRAY

- ✓ Motivation for Chapel
- ✓ Survey of Chapel Concepts
- ✓ Chapel Project and Characterizations
- Chapel Resources



Chapel Websites



Project page: http://chapel.cray.com

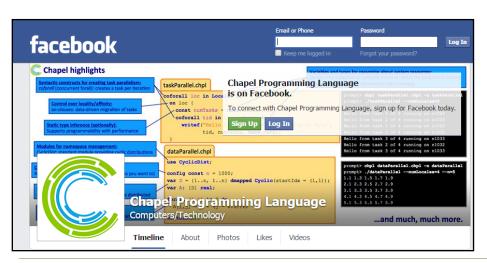
overview, papers, presentations, language spec, ...

GitHub: https://github.com/chapel-lang

download Chapel; browse source repository; contribute code

Facebook: https://www.facebook.com/ChapelLanguage

Twitter: https://twitter.com/ChapelLanguage





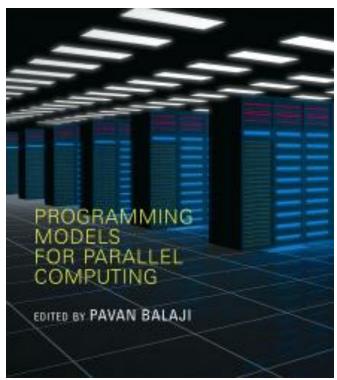


Suggested Reading



Chapel chapter from **Programming Models for Parallel Computing**

- a detailed overview of Chapel's history, motivating themes, features
- published by MIT Press, November 2015
- edited by Pavan Balaji (Argonne)
- chapter is now also available <u>online</u>



Other Chapel papers/publications available at http://chapel.cray.com/papers.html



Chapel Blog Articles



Chapel: Productive Parallel Programming, Cray Blog, May 2013.

a short-and-sweet introduction to Chapel

Chapel Springs into a Summer of Code, Cray Blog, April 2016.

coverage of recent events

Six Ways to Say "Hello" in Chapel (parts 1, 2, 3), Cray Blog, Sep-Oct 2015.

a series of articles illustrating the basics of parallelism and locality in Chapel

Why Chapel? (parts 1, 2, 3), Cray Blog, Jun-Oct 2014.

 a series of articles answering common questions about why we are pursuing Chapel in spite of the inherent challenges

[Ten] Myths About Scalable Programming Languages, IEEE TCSC Blog (index available on chapel.cray.com "blog articles" page), Apr-Nov 2012.

 a series of technical opinion pieces designed to argue against standard reasons given for not developing high-level parallel languages



Chapel Mailing Lists



low-traffic / read-only:

chapel-announce@lists.sourceforge.net: announcements about Chapel

community lists:

chapel-users@lists.sourceforge.net: user-oriented discussion list chapel-developers@lists.sourceforge.net: developer discussions chapel-education@lists.sourceforge.net: educator discussions chapel-bugs@lists.sourceforge.net: public bug forum

(subscribe at SourceForge: http://sourceforge.net/p/chapel/mailman/)

To mail the Cray team:

chapel_info@cray.com: contact the team at Cray chapel_bugs@cray.com: for reporting non-public bugs

or use IRC (#chapel on chat.freenode.net) or StackOverflow





Current Events: Computer Language Benchmark Game



Computer Language Benchmarks Game



Chapel was recently added to the game:

As of Oct 17th:

for performance:

1 top entries: pidigits

2 top-5 entries: meteor, thread-ring

2 top-10 entries: fannkuch-redux, chameneos-redux

3 top-20 entries: n-body, spectral-

norm, binary-trees

for code compactness:

2 top entries: n-body, thread-ring

2 top-5 entries: spectral-norm, pidigits

COMPUTE

4 top-20 entries: mandelbrot, regexdna, chameneos-redux, meteor

The Computer Language Benchmarks Game

64-bit quad core data set

Will your toy benchmark program be faster if you write it in a different programming language? It depends how you write it!

Which programs are fast?

Which are succinct? Which are efficient?

Ada	<u>C</u>	Chape	<u> </u>	Cloj	ure	<u>C#</u>	<u>C++</u>
Dart	Er	lang	F#	For	tran	Go	Hack
<u>Haskell</u>		Java	Jav	JavaScript		Lisp	Lua
<u>OCaml</u>		Pascal	<u>Perl</u>		PHP	Ру	thon
Racket		Ruby	JRuby		Rus	t S	cala
Smalltalk			Swi	ift	TypeScript		



Computer Language Benchmarks Game



rite

Chapel was recently added to the game:

Oct 17th-

The Computer Language Benchmarks Game

64-bit quad core data set

We want easy answers, but easy answers are often incomplete or wrong. You and I know, there's more we should understand:

stories

details

fast?

conclusions

{ for researchers }

4 top-20 entries: mandelbrot, regexdna, chameneos-redux, meteor Racket Ruby JRuby Rust Scala

Smalltalk Swift TypeScript





Chapel:Productive Parallel Programming at Scale

Questions?

