



Chapel: Overview and Features for Heterogeneity

Brad Chamberlain, Chapel Team, Cray Inc.

Heterogeneous C++ Working Group

February 5, 2018



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Safe Harbor Statement



This presentation may contain forward-looking statements that are based on our current expectations. Forward looking statements may include statements about our financial guidance and expected operating results, our opportunities and future potential, our product development and new product introduction plans, our ability to expand and penetrate our addressable markets and other statements that are not historical facts. These statements are only predictions and actual results may materially vary from those projected. Please refer to Cray's documents filed with the SEC from time to time concerning factors that could affect the Company and these forward-looking statements.



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Plan for this morning

- Chapel context
- Overview of example core Chapel features
- Additional detail on features related to heterogeneity
- Feel free to interrupt for Questions / Discussion



Chapel and Heterogeneity in a Nutshell



- **Chapel's design supports hardware heterogeneity**
 - and in a way that is user-extensible without compiler changes
- **That said, in practice...**
 - ...virtually all of our work has been on homogeneous cases
 - ...we haven't spend much time on many of the hardest cases
 - e.g., no FPGA work, less GPU work than we'd like



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What is Chapel?



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What is Chapel?



Chapel: A productive parallel programming language

- portable
- open-source
- a collaborative effort



Goals:

- Support general parallel programming
 - “any parallel algorithm on any parallel hardware”
- Make parallel programming at scale far more productive



Scalable Parallel Programming Concerns



Q: What do HPC programmers need from a language?

A: *Serial Code:* Software engineering and performance

Parallelism: What should execute simultaneously?

Locality: Where should those tasks execute?

Mapping: How to map the program to the system?

Separation of Concerns: Decouple these issues

*Chapel is a language designed to address these needs
from first principles*



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Chapel and Other Languages

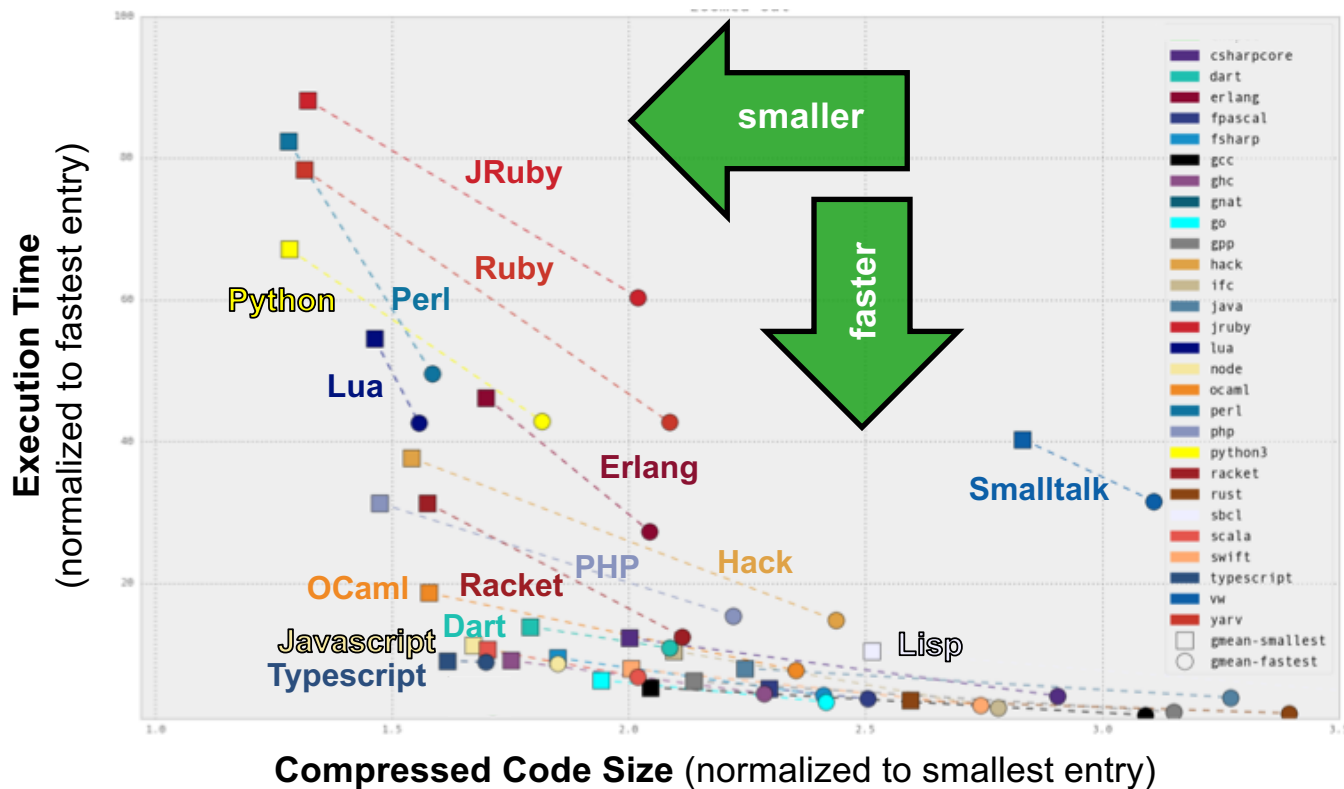
Chapel strives to be as...

- ...**programmable** as Python
- ...**fast** as Fortran
- ...**scalable** as MPI, SHMEM, or UPC
- ...**portable** as C
- ...**flexible** as C++
- ...**fun** as [your favorite programming language]



CLBG Cross-Language Summary

(Oct 2017 standings)



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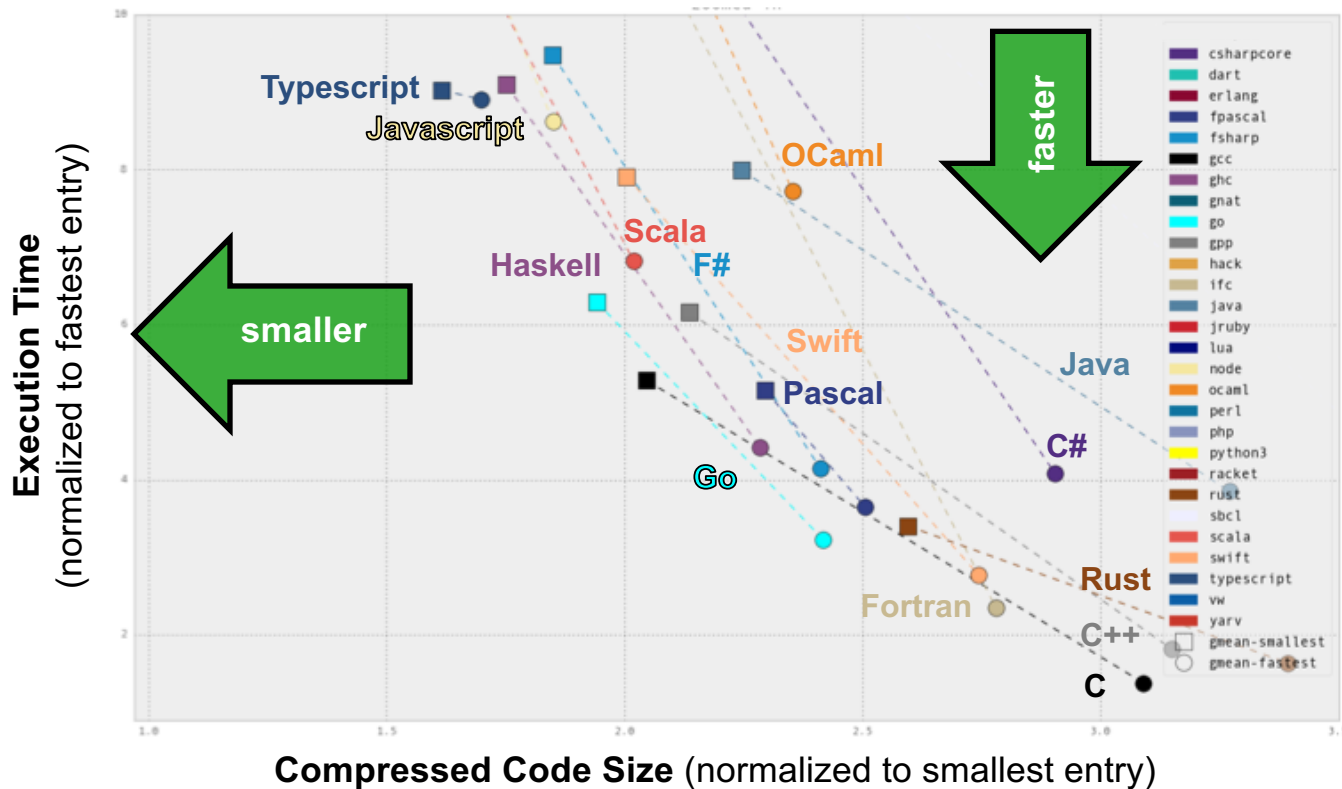
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CLBG Cross-Language Summary

(Oct 2017 standings, zoomed in)



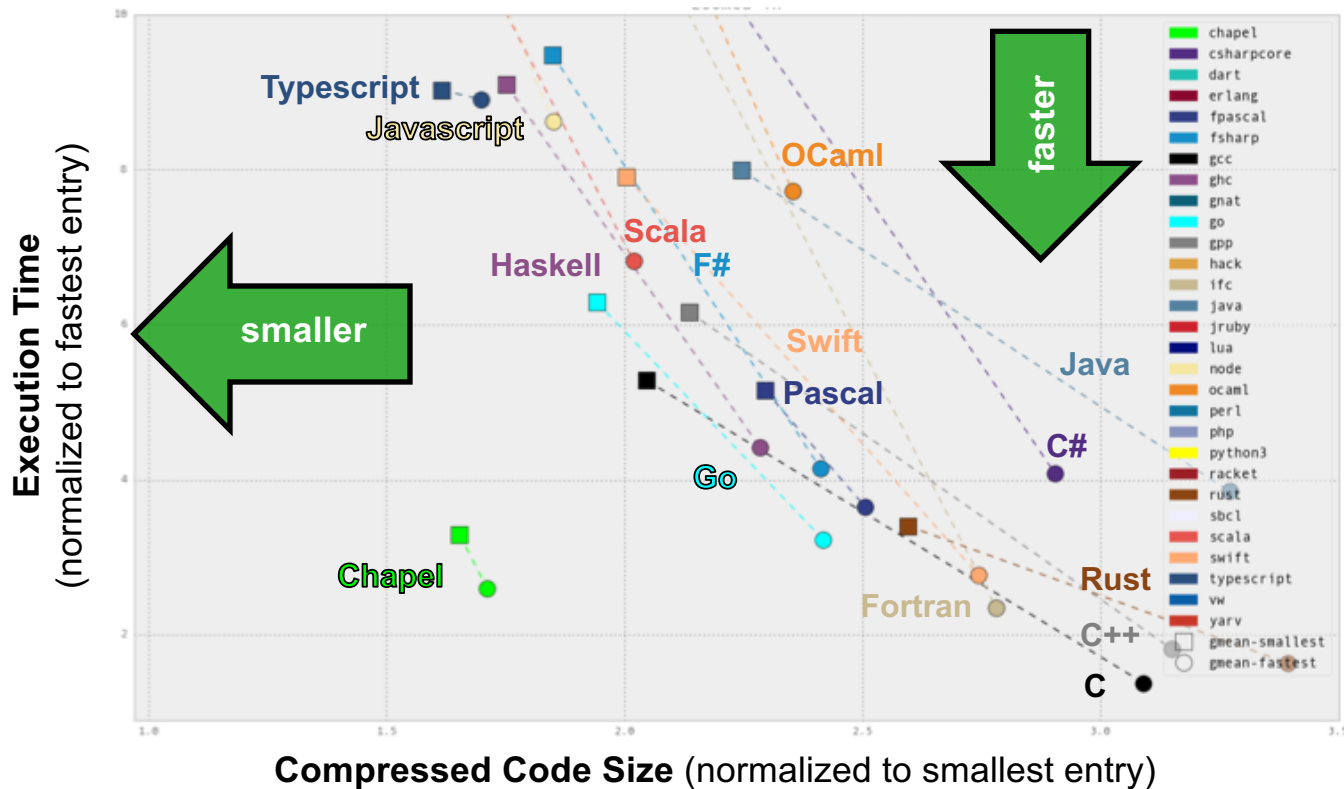
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CLBG Cross-Language Summary

(Oct 2017 standings, zoomed in)



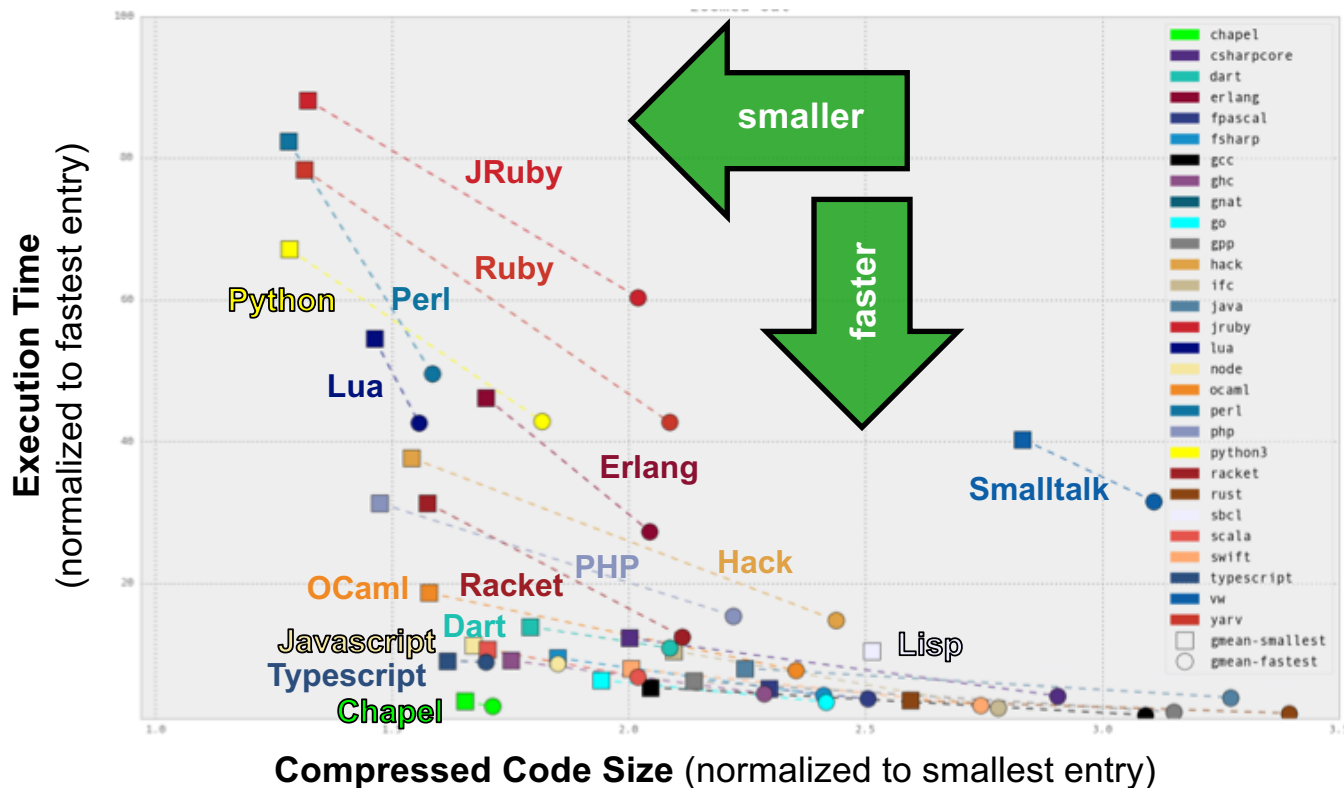
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CLBG Cross-Language Summary

(Oct 2017 standings)



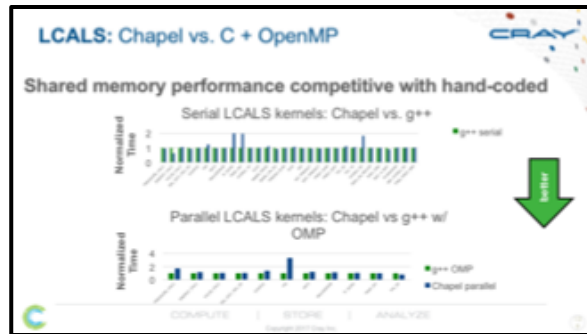
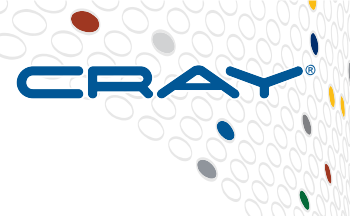
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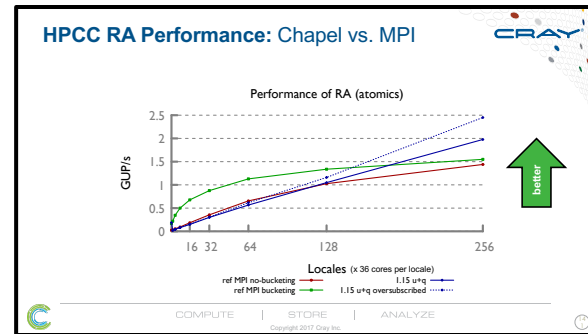


Chapel Performance: HPC Benchmarks



LCALS

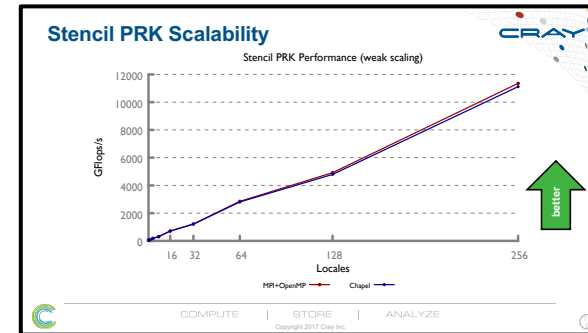
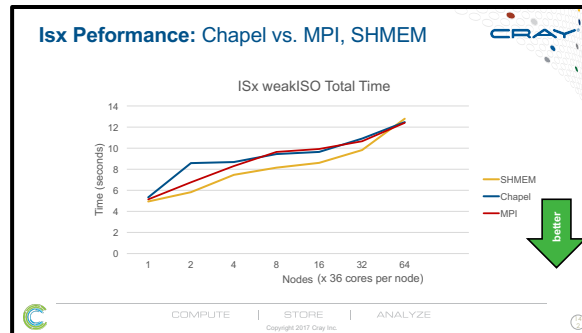
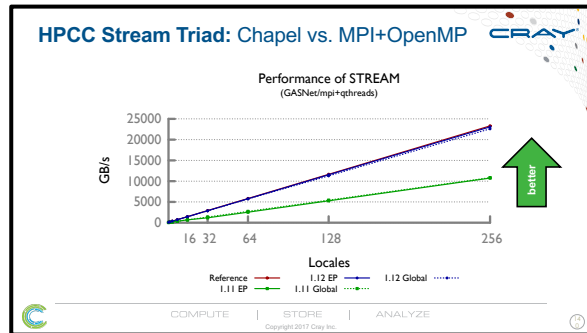
HPCC RA



STREAM
Triad

ISx

PRK
Stencil



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Nightly performance graphs online
at: <https://chapel-lang.org/perf>

The Chapel Team at Cray (May 2017)



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Chapel Community Partners



Lawrence Berkeley
National Laboratory



Sandia National Laboratories



Yale

(and several others...)

<https://chapel-lang.org/collaborations.html>



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Introduction to Chapel, by Example



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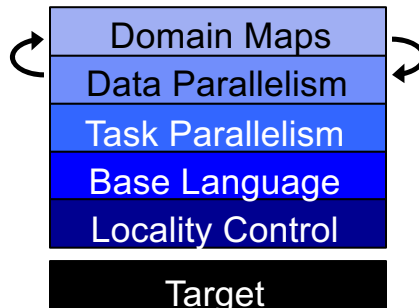
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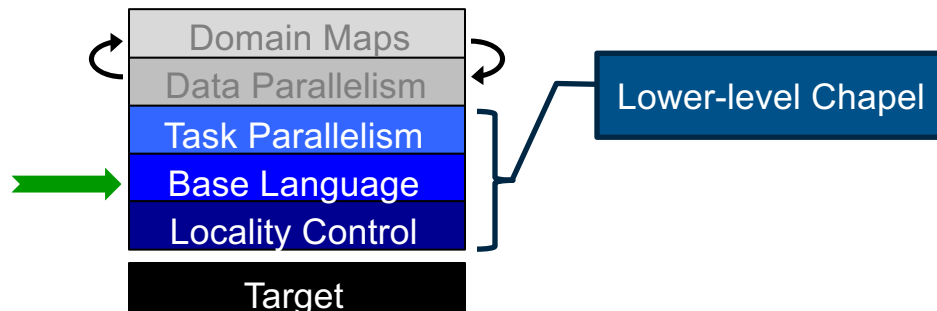
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Chapel language feature areas

Chapel language concepts



Base Language



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Base Language Features, by example

```
iter fib(n) {  
  var current = 0,  
      next = 1;  
  
  for i in 1..n {  
    yield current;  
    current += next;  
    current <=> next;  
  }  
}
```

```
config const n = 10;  
  
for f in fib(n) do  
  writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```



Base Language Features, by example



Configuration declarations
(to avoid command-line argument parsing)
`./a.out --n=1000000`

```
iter fib(n) {  
  var current = 0,  
      next = 1;  
  
  for i in 1..n {  
    yield current;  
    current += next;  
    current <=> next;  
  }  
}
```

```
config const n = 10;  
  
for f in fib(n) do  
  writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```



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Base Language Features, by example

Modern iterators

```
iter fib(n) {  
  var current = 0,  
      next = 1;  
  
  for i in 1..n {  
    yield current;  
    current += next;  
    current <=> next;  
  }  
}
```

```
config const n = 10;  
  
for f in fib(n) do  
  writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```





Base Language Features, by example

- Static type inference for:
- arguments
 - return types
 - variables

```
iter fib(n) {  
  var current = 0,  
      next = 1;  
  
  for i in 1..n {  
    yield current;  
    current += next;  
    current <=> next;  
  }  
}
```

```
config const n = 10;  
  
for f in fib(n) do  
  writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```



Base Language Features, by example

Zippered iteration

```
iter fib(n) {  
  var current = 0,  
      next = 1;  
  
  for i in 1..n {  
    yield current;  
    current += next;  
    current <=> next;  
  }  
}
```

```
config const n = 10;  
  
for (i,f) in zip(0..#n, fib(n)) do  
  writeln("fib #", i, " is ", f);
```

```
fib #0 is 0  
fib #1 is 1  
fib #2 is 1  
fib #3 is 2  
fib #4 is 3  
fib #5 is 5  
fib #6 is 8  
...
```



Base Language Features, by example



Range types and operators

```
iter fib(n) {  
  var current = 0;  
  next = 1;  
  
  for i in 1..n {  
    yield current;  
    current += next;  
    current <=> next;  
  }  
}
```

```
config const n = 10;  
  
for (i,f) in zip(0..#n, fib(n)) do  
  writeln("fib #", i, " is ", f);
```

```
fib #0 is 0  
fib #1 is 1  
fib #2 is 1  
fib #3 is 2  
fib #4 is 3  
fib #5 is 5  
fib #6 is 8  
...
```



Base Language Features, by example



tuples

```
iter fib(n) {  
    var current = 0,  
        next = 1;  
  
    for i in 1..n {  
        yield current;  
        current += next;  
        current <=> next;  
    }  
}
```

```
config const n = 10;  
  
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fib #0 is 0  
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...
```



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Base Language Features, by example

```
iter fib(n) {  
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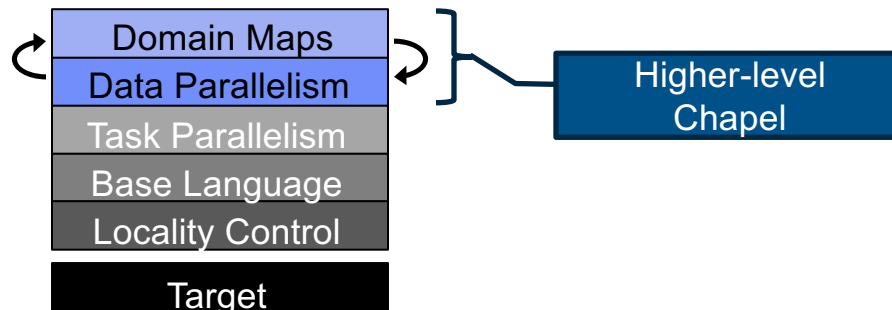
```
fib #0 is 0  
fib #1 is 1  
fib #2 is 1  
fib #3 is 2  
fib #4 is 3  
fib #5 is 5  
fib #6 is 8  
...
```



Data Parallelism in Chapel



Chapel language concepts



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Data Parallelism, by example

dataParallel.chpl

```
config const n = 1000;  
var D = {1..n, 1..n};  
  
var A: [D] real;  
forall (i,j) in D do  
  A[i,j] = i + (j - 0.5)/n;  
writeln(A);
```

```
prompt> chpl dataParallel.chpl -o dataParallel  
prompt> ./dataParallel --n=5  
1.1 1.3 1.5 1.7 1.9  
2.1 2.3 2.5 2.7 2.9  
3.1 3.3 3.5 3.7 3.9  
4.1 4.3 4.5 4.7 4.9  
5.1 5.3 5.5 5.7 5.9
```



Data Parallelism, by example



Domains (Index Sets)

dataParallel.chpl

```
config const n = 1000;  
var D = {1..n, 1..n};  
  
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4.1 4.3 4.5 4.7 4.9  
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```



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Data Parallelism, by example



Arrays

dataParallel.chpl

```
config const n = 1000;  
var D = {1..n, 1..n};  
  
var A: [D] real;  
forall (i,j) in D do  
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5.1 5.3 5.5 5.7 5.9
```



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Data Parallelism, by example

Data-Parallel Forall Loops

dataParallel.chpl

```
config const n = 1000;  
var D = {1..n, 1..n};  
  
var A: [D] real;  
forall (i,j) in D do  
  A[i,j] = i + (j - 0.5)/n;  
writeln(A);
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5.1 5.3 5.5 5.7 5.9
```



Data Parallelism, by example

This is a shared memory program

Nothing has referred to remote locales, explicitly or implicitly

dataParallel.chpl

```
config const n = 1000;  
var D = {1..n, 1..n};  
  
var A: [D] real;  
forall (i,j) in D do  
  A[i,j] = i + (j - 0.5)/n;  
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5.1 5.3 5.5 5.7 5.9
```

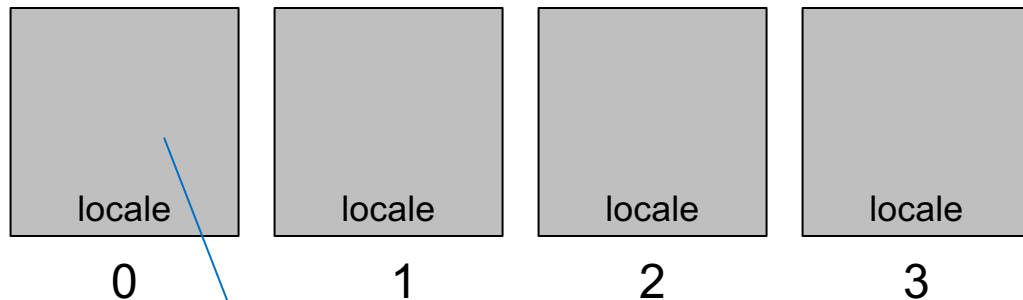


Locales



- Unit of the target system useful for reasoning about locality
 - Each locale can run tasks and store variables
 - Has processors and memory (or can defer to something that does)
 - For most HPC systems, locale == compute node

Locales:



User's main() executes on locale #0



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Data Parallelism, by example

This is a shared memory program

Nothing has referred to remote locales, explicitly or implicitly

dataParallel.chpl

```
config const n = 1000;  
var D = {1..n, 1..n};  
  
var A: [D] real;  
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3.1 3.3 3.5 3.7 3.9  
4.1 4.3 4.5 4.7 4.9  
5.1 5.3 5.5 5.7 5.9
```



Distributed Data Parallelism, by example

Domain Maps
(Map Data Parallelism to the System)

dataParallel.chpl

```
use CyclicDist;  
config const n = 1000;  
var D = {1..n, 1..n}  
      dmapped Cyclic(startIdx = (1,1));  
var A: [D] real;  
forall (i,j) in D do  
  A[i,j] = i + (j - 0.5)/n;  
writeln(A);
```

```
prompt> chpl dataParallel.chpl -o dataParallel  
prompt> ./dataParallel --n=5 --numLocales=4  
1.1 1.3 1.5 1.7 1.9  
2.1 2.3 2.5 2.7 2.9  
3.1 3.3 3.5 3.7 3.9  
4.1 4.3 4.5 4.7 4.9  
5.1 5.3 5.5 5.7 5.9
```



Distributed Data Parallelism, by example

magic? HPF-like?

descriptive?

Not in the slightest...

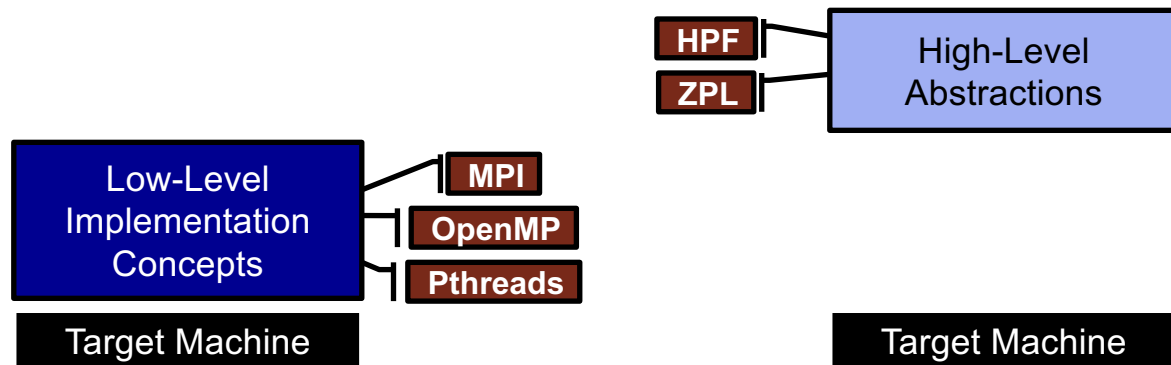
- Lowering of code is well-defined
- User can control details
- Part of Chapel's *multiresolution philosophy*...

dataParallel.chpl

```
use CyclicDist;
config const n = 1000;
var D = {1..n, 1..n}
        dmapped Cyclic(startIdx = (1,1));
var A: [D] real;
forall (i,j) in D do
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3.1 3.3 3.5 3.7 3.9
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5.1 5.3 5.5 5.7 5.9
```


Chapel's Multiresolution Design: Motivation



“Why is everything so tedious/difficult?”

“Why don’t my programs trivially port to new systems?”

“Why don’t I have more control?”

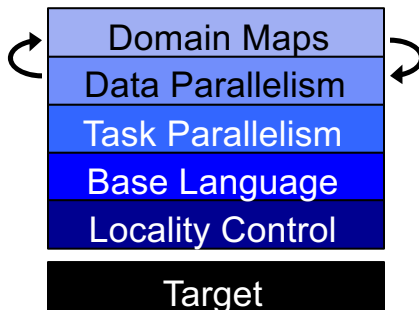


Chapel's Multiresolution Philosophy



Multiresolution Design: Support multiple tiers of features

- higher levels for programmability, productivity
- lower levels for greater degrees of control



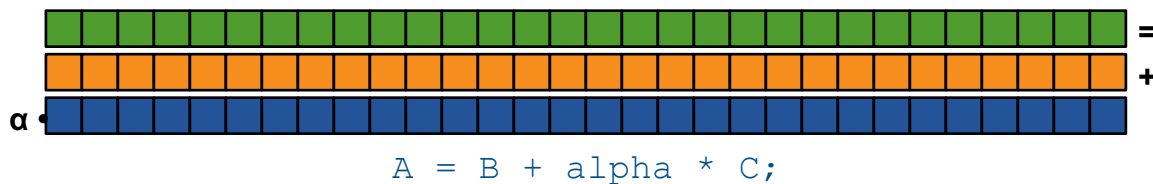
- build the higher-level concepts in terms of the lower
- permit users to intermix layers arbitrarily



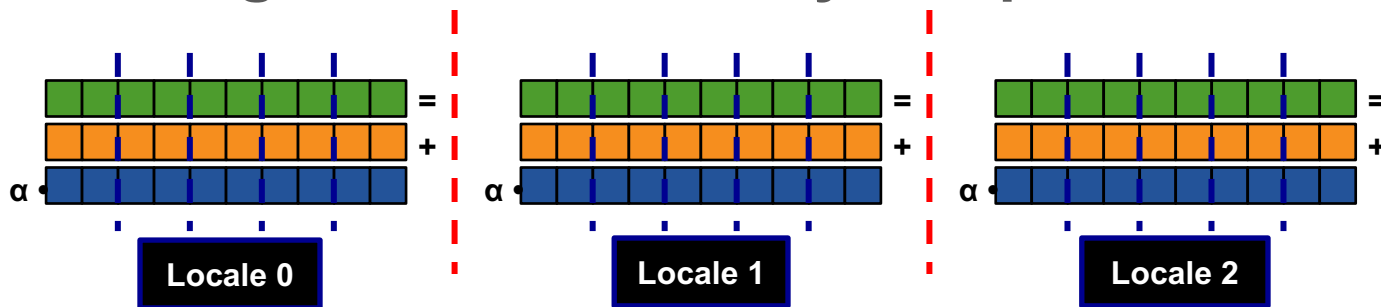
Domain Maps: A Multiresolution Feature



Domain maps are “recipes” that instruct the compiler how to map the global view of a computation...



...to the target locales' memory and processors:



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Authoring Domain Maps



- **Users can write their own domain maps**

- Implemented within Chapel itself
- Create an object type per concept:
 - The domain map itself
 - A domain
 - An array
- Make them satisfy a standard interface
 - e.g., arrays must support iteration, random access, etc.
- Compiler targets this interface in implementing the language
- Goal: make the language flexible, future-proof

Note: *all* Chapel arrays are implemented this way



Distributed Data Parallelism, by example

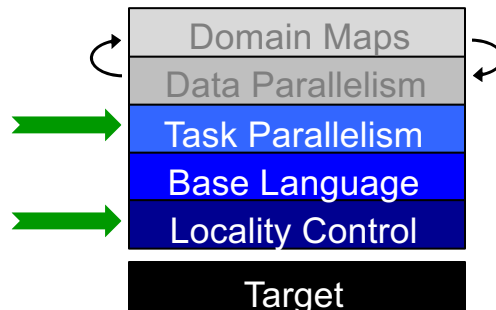
dataParallel.chpl

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var D = {1..n, 1..n}  
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4.1 4.3 4.5 4.7 4.9  
5.1 5.3 5.5 5.7 5.9
```



Task Parallelism and Locality Control



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Task Parallelism and Locality, by example

taskParallel.chpl

```
coforall loc in Locales do
  on loc {
    const numTasks = here.numPUs();
    coforall tid in 1..numTasks do
      writef("Hello from task %n of %n "+
            "running on %s\n",
            tid, numTasks, here.name);
  }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



Task Parallelism and Locality, by example

Abstraction of
System Resources

taskParallel.chpl

```
coforall loc in Locales do
  on loc {
    const numTasks = here.numPUs();
    coforall tid in 1..numTasks do
      writef("Hello from task %n of %n "+
            "running on %s\n",
            tid, numTasks, here.name);
    }
}
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



Task Parallelism and Locality, by example

High-Level
Task Parallelism

taskParallel.chpl

```
coforall loc in Locales do
  on loc {
    const numTasks = here.numPUs();
    coforall tid in 1..numTasks do
      writef("Hello from task %n of %n " +
            "running on %s\n",
            tid, numTasks, here.name);
    }
  }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



Task Parallelism and Locality, by example

taskParallel.chpl

```
coforall loc in Locales do
  on loc {
    const numTasks = here.numPUs();
    coforall tid in 1..numTasks do
      writef("Hello from task %n of %n "+
            "running on %s\n",
            tid, numTasks, here.name);
    }
```

Control of Locality/Affinity

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



Task Parallelism and Locality, by example

Abstraction of
System Resources

taskParallel.chpl

```
coforall loc in Locales do
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            "running on %s\n",
            tid, numTasks, here.name);
  }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



Task Parallelism and Locality, by example

High-Level
Task Parallelism

taskParallel.chpl

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coforall loc in Locales do
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  }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
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Hello from task 1 of 2 running on n1033
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Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



Task Parallelism and Locality, by example

Not seen here:

Data-centric task coordination
via atomic and full/empty vars

taskParallel.chpl

```
coforall loc in Locales do
  on loc {
    const numTasks = here.numPUs();
    coforall tid in 1..numTasks do
      writef("Hello from task %n of %n " +
            "running on %s\n",
            tid, numTasks, here.name);
    }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



Task Parallelism and Locality, by example

taskParallel.chpl

```
coforall loc in Locales do
  on loc {
    const numTasks = here.numPUs();
    coforall tid in 1..numTasks do
      writef("Hello from task %n of %n "+
            "running on %s\n",
            tid, numTasks, here.name);
    }
}
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



Parallelism and Locality: Distinct in Chapel



- This is a **parallel**, but local program:

```
coforall i in 1..msgs do
  writeln("Hello from task ", i);
```

- This is a **distributed**, but serial program:

```
writeln("Hello from locale 0!");
on Locales[1] do writeln("Hello from locale 1!");
on Locales[2] do writeln("Hello from locale 2!");
```

- This is a **distributed parallel** program:

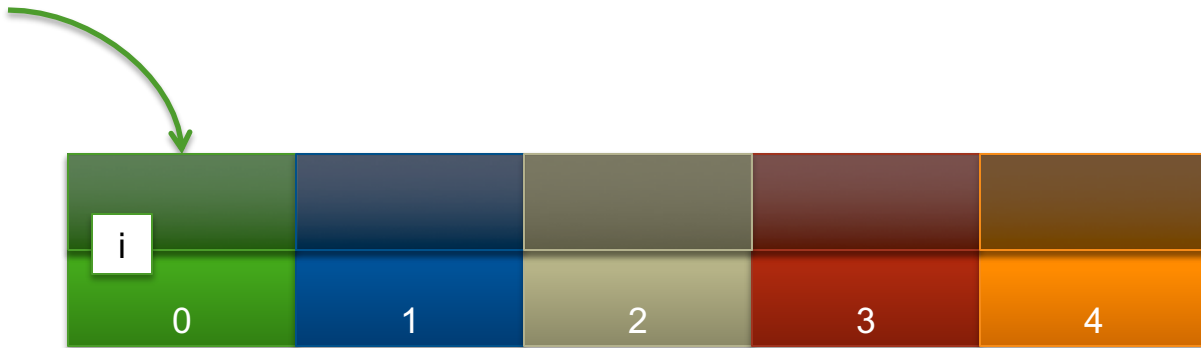
```
coforall i in 1..msgs do
  on Locales[i%numLocales] do
    writeln("Hello from task ", i,
           " running on locale ", here.id);
```



Chapel: Scoping and Locality



```
var i: int;
```



Locales (think: “compute nodes”)

COMPUTE

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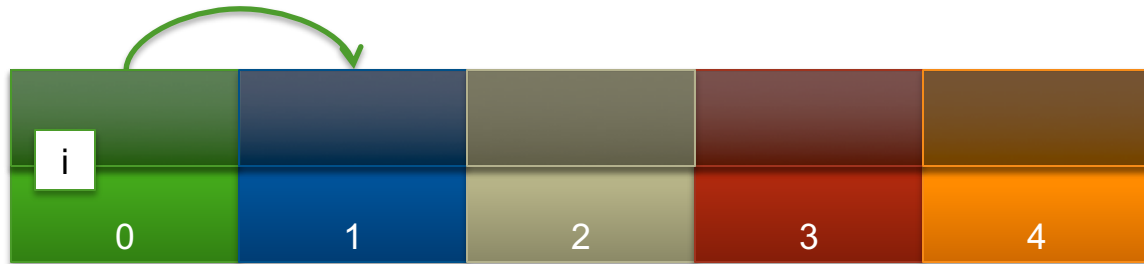
ANALYZE



Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {
```



Locales (think: “compute nodes”)

COMPUTE

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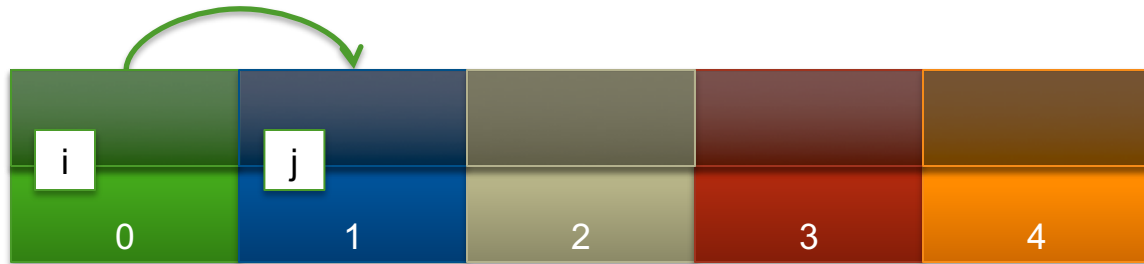
ANALYZE



Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
  var j: int;
```



Locales (think: “compute nodes”)

COMPUTE

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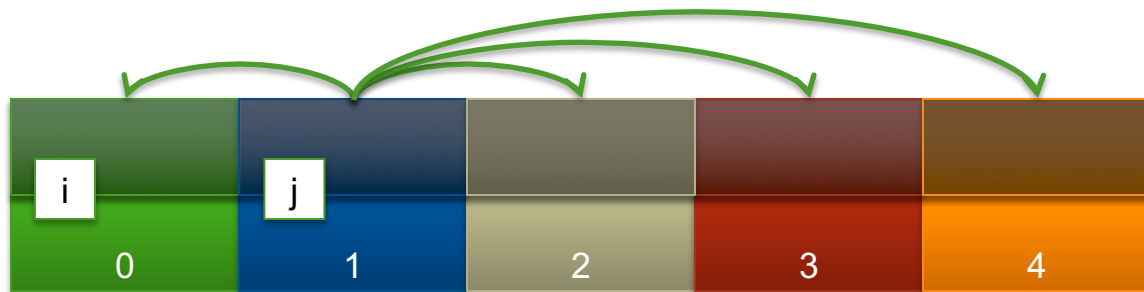
ANALYZE



Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
  var j: int;  
  coforall loc in Locales {  
    on loc {
```



Locales (think: “compute nodes”)

COMPUTE

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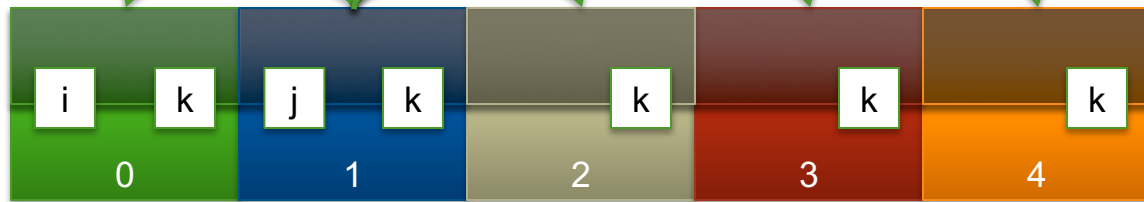
ANALYZE



Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
  var j: int;  
  coforall loc in Locales {  
    on loc {  
      var k: int;  
      ...  
    }  
  }  
}
```



Locales (think: “compute nodes”)

COMPUTE

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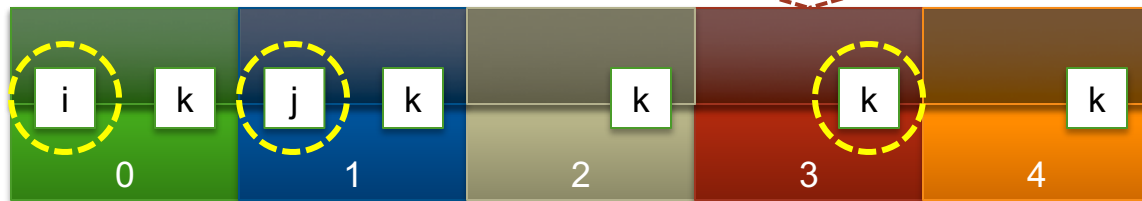
Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
  var j: int;  
  coforall loc in Locales {  
    on loc {  
      var k: int;  
      k = 2*i + j;  
    }  
  }  
}
```

OK to access i, j , and k
wherever they live

$k = 2*i + j;$



Locales (think: “compute nodes”)

COMPUTE

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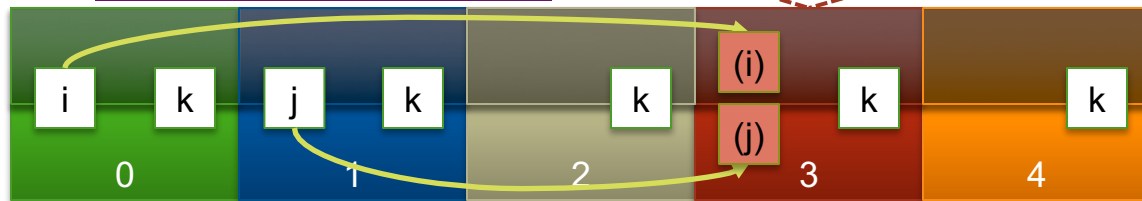
Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
  var j: int;  
  coforall loc in Locales {  
    on loc {  
      var k: int;  
      k = 2*i + j;  
    }  
  }  
}
```

here, i and j are remote, so
the compiler + runtime will
transfer their values

$k = 2*i + j;$



Locales (think: “compute nodes”)

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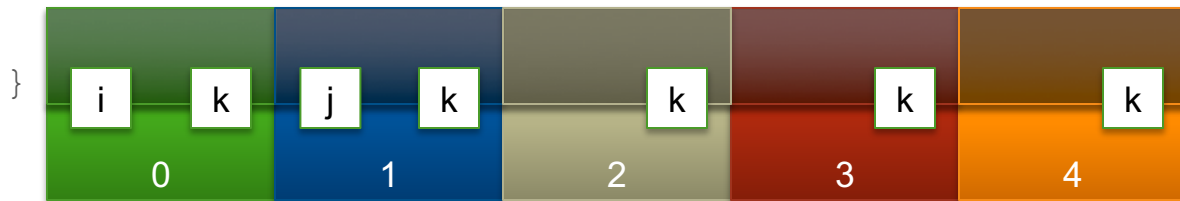
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Chapel: Locality queries



```
var i: int;  
on Locales[1] {  
  var j: int;  
  coforall loc in Locales {  
    on loc {  
      ...here...           // query the locale on which this task is running  
      ...j.locale...       // query the locale on which j is stored  
      ...here.physicalMemory(...) ... // query system characteristics  
      ...here.runningTasks() ...      // query runtime characteristics  
    }  
  }  
}
```



Locales (think: “compute nodes”)

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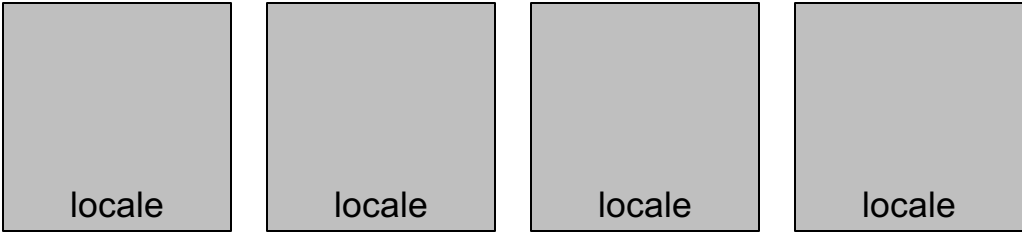
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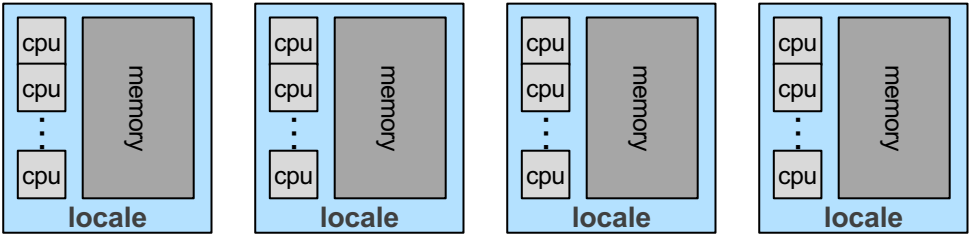


Classic Locales

- Historically, Chapel's locales were black boxes
 - Intra-node concerns handled by compiler, runtime, OS



- This was sufficient when compute nodes were simple

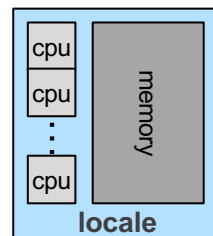
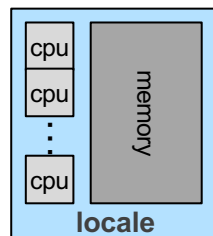
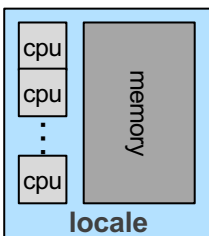
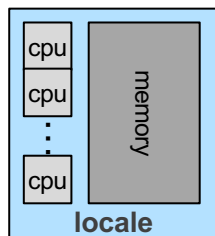
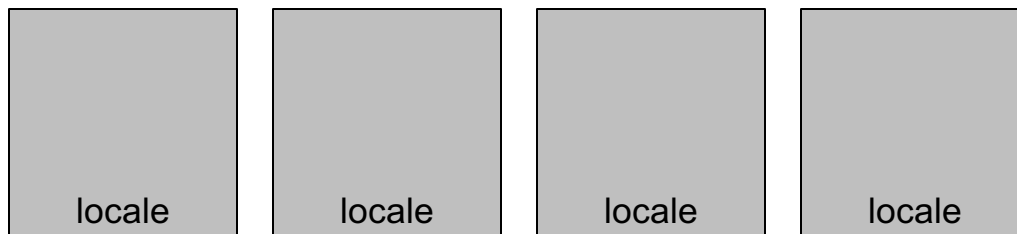


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Classic Locales



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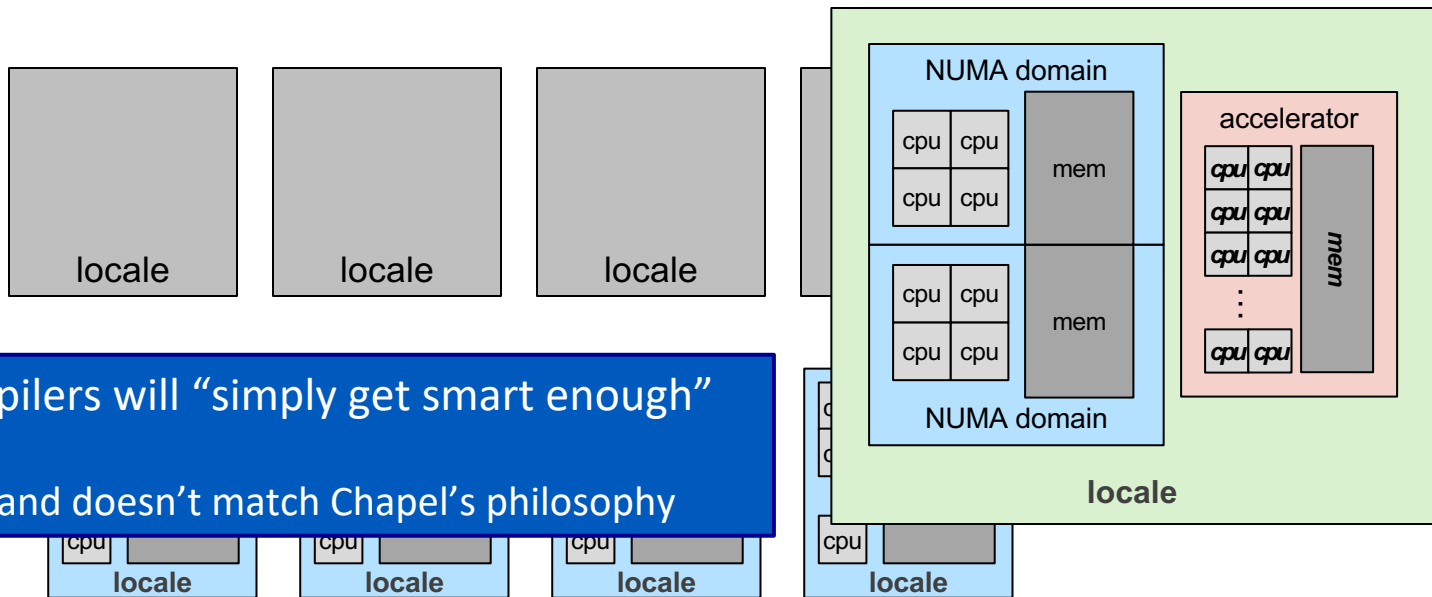
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Classic Locales



- **Classic model breaks down for more complex cases**
 - E.g. multiple flavors of memory or processors



Could hope compilers will “simply get smart enough”

...but seems naïve and doesn't match Chapel's philosophy



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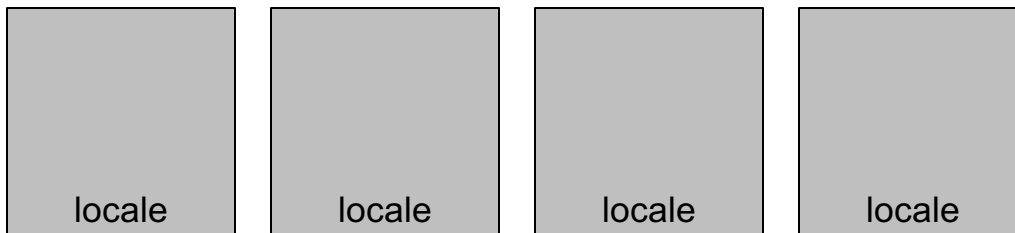
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Hierarchical Locales



- So, we made locales hierarchical



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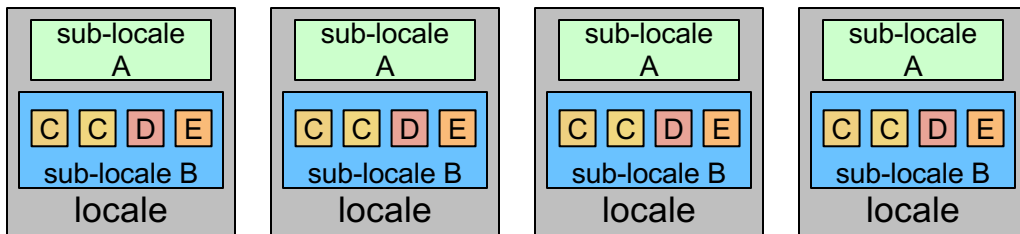
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Hierarchical Locales

- So, we made locales hierarchical

- Locales can now themselves contain locales
 - E.g., an accelerator sub-locale, a scratchpad memory sub-locale



- Target sub-locales with on-clauses, as before


```
on Locales[0].GPU do computationThatLikesGPUs();
```

 - Ideally, hide such logic in abstractions: domain maps, parallel iterators
- Introduced a new multiresolution type: *locale models*

Chapel's Locale Models



- **User-specified type representing locales**
- **Similar goals to domain maps:**
 - Support user implementation of key high-level abstractions
 - Make language future-proof (w.r.t. emerging architectures)



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Authoring a Locale Model



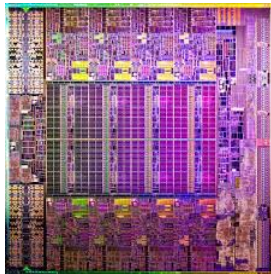
- **Creating a locale model:**

- Create a top-level locale object type
 - In turn, it can contain fields representing sub-locales
- Each locale / sub-locale type must meet a required interface:
 - **Memory:** How is it managed? (malloc, realloc, free)
 - **Tasking:** How do I launch and synchronize tasks?
 - **Communication:** How are data & control transferred between locales?
 - gets, puts, active messages
 - widening of pointers



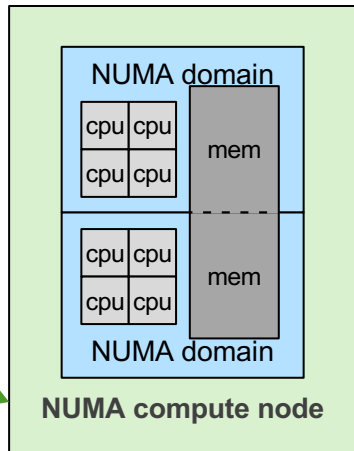
An Example: The numa Locale Model

physical



<http://www1.pcimag.com/media/images/337192-intel-xeon-e5-chip.jpg?thumb=y>

conceptual



`$CHPL_HOME/modules/.../numa/LocaleModel.chpl`

```
class NumaDomain : AbstractLocaleModel {
  const sid: chpl_sublocID_t;
}

// The node model
class LocaleModel : AbstractLocaleModel {
  const numSublocales: int;
  var childSpace: domain(1);
  var childLocales: [childSpace] NumaDomain;
}

// support for memory management
proc chpl_here_alloc(size:int, md:int(16)) { ... }

// support for "on" statements
proc chpl_executeOn
  (loc: chpl_localeID_t, // target locale
   fn: int,             // on-body func idx
   args: c_void_ptr,    // func args
   args_size: int(32)   // args size
  ) { ... }

// support for tasking stmts: begin, cobegin, coforall
proc chpl_taskListAddCoStmt
  (subloc_id: int,      // target subloc
   fn: int,            // body func idx
   args: c_void_ptr,   // func args
   ref tlist: _task_list, // task list
   tlist_node_id: int   // task list owner
  ) { ... }
```





Locale Models: Status

- **All Chapel compilations use a locale model**
 - Set via environment variable or compiler flag
- **Current locale models:**
 - **flat:** the default, has no sublocales (as in the classic model)
 - **numa:** supports a sub-locale per NUMA domain within the node
 - **knl:** for Intel® Xeon Phi™: numa w/ sublocale for HBM/MCDRAM
- **In practice...**
 - we use the 'flat' locale model almost exclusively





Performance: 'numa' vs. 'flat'

- **Using 'numa' leads to performance overheads**
 - Local arrays must be “chunked” between the numa sublocales
 - Indexing must do extra work to pick the right chunk
- **Though 'flat' has no sub-locales, it's also NUMA-aware**
 - First-touch heuristics used to map sub-arrays to NUMA domains
 - Yet array remains contiguous in memory \Rightarrow simple indexing
- **As a result, 'numa' rarely outperforms 'flat'**



Challenges: Static analysis & locality

- **Local vs. Remote Locales** (distributed or shared memory?)
 - In general, given:
`on x.locale ...`
can't statically tell whether locale shares memory with ``here`` or not
 - May result in overheads due to conservatism
 - “Assuming it's remote, I'll introduce wide pointers & communication, ...”





Challenges: Memory-only locales

- **What if a locale only represents memory? (say)**
 - Have interface run tasks on other locales with processors
 - Choice of policy is up to locale model author
 - Round-robin, dynamic load-balance, nearest, ...



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Wrapping Up



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Summary



- **Chapel's design uses a multiresolution philosophy**
 - High-level for productivity
 - Low-level for control
 - User-extensible for flexibility, future-proof design
- **Locale models support mapping to new architectures**
 - Provide bridge from compiler to system resources
 - Enables targeting of heterogeneous resources
- **Chapel performance can match C+MPI+OpenMP**
 - With improvements in readability, writability, code size



Possible Discussion Topics

- Error-handling / Exceptions




Chapel Resources




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Home
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chapel-lang.org
chapel_info@cray.com



The Chapel Parallel Programming Language

What is Chapel?

Chapel is a modern programming language that is...

- **parallel**: contains first-class concepts for concurrent and parallel computation
- **productive**: designed with programmability and performance in mind
- **portable**: runs on laptops, clusters, the cloud, and HPC systems
- **scalable**: supports locality-oriented features for distributed memory systems
- **open-source**: hosted on [GitHub](#), permissively [licensed](#)

New to Chapel?

As an introduction to Chapel, you may want to...

- read a [blog article](#) or [book chapter](#)
- watch [an overview talk](#) or browse its [slides](#)
- [download](#) the release
- browse [sample programs](#)
- view [other resources](#) to learn how to trivially write distributed programs like this:

```
use CyclicDist;           // use the Cyclic distribution library
config const n = 100;     // use ./a.out --n=<val> to override this default

forall i in {1..n} mapped Cyclic(startId=i) do
  writeln("Hello from iteration ", i, " of ", n, " running on node ", here.id);
```

What's Hot?

- **Chapel 1.16** is now available—[download](#) a copy today!
- The **CHIUV 2018** [call for participation](#) is now available!
- A recent [Cray blog post](#) reports on highlights from CHIUV 2017.
- Chapel is now one of the supported languages on [Try It Online!](#)
- Watch talks from [ACCU 2017](#), [CHIUV 2017](#), and [ATPESC 2016](#) on [YouTube](#).
- [Browse slides](#) from **PADAL**, **EAGE**, **EMBRACE**, **ACCU**, and other recent talks.
- See also: [What's New?](#)





How to Track Chapel

<http://facebook.com/ChapelLanguage>

<http://twitter.com/ChapelLanguage>

[https://www.youtube.com/channel/UCHmm27bYjhknK5mU7ZzPGsQ/
chapel-announce@lists.sourceforge.net](https://www.youtube.com/channel/UCHmm27bYjhknK5mU7ZzPGsQ/chapel-announce@lists.sourceforge.net)



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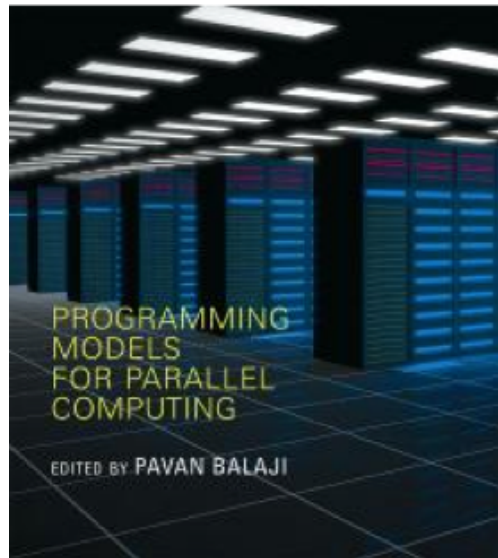
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Suggested Reading (healthy attention spans)

Chapel chapter from [*Programming Models for Parallel Computing*](#)

- a detailed overview of Chapel's history, motivating themes, features
- published by MIT Press, November 2015
- edited by Pavan Balaji (Argonne)
- chapter is now also available [online](#)



Other Chapel papers/publications available at <https://chapel-lang.org/papers.html>



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Suggested Reading (short attention spans)



[CHIUV 2017: Surveying the Chapel Landscape](#), [Cray Blog](#), July 2017.

- *a run-down of recent events*

[Chapel: Productive Parallel Programming](#), [Cray Blog](#), May 2013.

- *a short-and-sweet introduction to Chapel*

[Six Ways to Say “Hello” in Chapel](#) (parts [1](#), [2](#), [3](#)), [Cray Blog](#), Sep-Oct 2015.

- *a series of articles illustrating the basics of parallelism and locality in Chapel*

[Why Chapel?](#) (parts [1](#), [2](#), [3](#)), [Cray Blog](#), Jun-Oct 2014.

- *a series of articles answering common questions about why we are pursuing Chapel in spite of the inherent challenges*

[\[Ten\] Myths About Scalable Programming Languages](#), [IEEE TCSC Blog](#)

([index available on chapel-lang.org “blog posts” page](#)), Apr-Nov 2012.

- *a series of technical opinion pieces designed to argue against standard reasons given for not developing high-level parallel languages*



Chapel StackOverflow and GitHub Issues

The screenshot shows the Stack Overflow website with search results for the term 'chapel'. The top navigation bar includes links for Questions, Jobs, Documentation, Tags, and Users. The search bar contains the text '[chapel]'. Below the navigation bar, there's a section for 'Tagged Questions' with filters for 'info', 'newest', 'frequent', 'votes', and 'active'. The first question is 'Can one generate a grid of the Locales where a Distribution is mapped?' with 2 votes and 22 views. The second question is 'Is "[<var> in <distributed variable>]" equivalent to "forall"?' with 3 votes and 24 views. The third question is 'Get Non-primitive Variables from within a Cobegin - Chapel' with 2 votes and 45 views. The fourth question is 'Is there a default String conversion method in Chapel?' with 3 votes and 1 view.

The screenshot shows the GitHub repository page for 'chapel-lang / chapel'. The repository has 45 watchers, 455 stars, and 145 forks. The 'Issues' tab is selected, showing 292 issues. The issues are filtered by 'is:issue is:open'. The list of issues includes:

- 292 Open, 77 Closed
- Implement "bounded-coforall" optimization for remote coforalls (area: Compiler, type: Performance, #6357 opened 13 hours ago by ronawho)
- Consider using processor atomics for remote coforalls EndCount (area: Compiler, type: Performance, #6356 opened 13 hours ago by ronawho)
- make uninstall (area: BTR, type: Feature Request, #6353 opened 14 hours ago by mppf)
- make check doesn't work with ./configure (area: BTR, #6352 opened 16 hours ago by mppf)
- Passing variable via in intent to a forall loop seems to create an iteration-private variable, not a task-private one (area: Compiler, type: Bug, #6351 opened a day ago by cassella)
- Remove chpl_comm_make_progress (area: Runtime, easy, type: Design, #6349 opened a day ago by sungeunchol)
- Runtime error after make on Linux Mint (area: BTR, user issue, #6348 opened a day ago by danindiana)



Where to..



Submit bug reports:

GitHub issues for chapel-lang/chapel: public bug forum

chapel_bugs@cray.com: for reporting non-public bugs

Ask User-Oriented Questions:

StackOverflow: when appropriate / other users might care

#chapel-users (irc.freenode.net): user-oriented IRC channel

chapel-users@lists.sourceforge.net: user discussions

Discuss Chapel development

chapel-developers@lists.sourceforge.net: developer discussions

#chapel-developers (irc.freenode.net): developer-oriented IRC channel

Discuss Chapel's use in education

chapel-education@lists.sourceforge.net: educator discussions

Directly contact Chapel team at Cray: chapel_info@cray.com



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Questions?



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